## KHRONOS

# The OpenCL C Specification 

Version: 2.0
Document Revision: 33

Khronos OpenCL Working Group

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## 6. The OpenCL C Programming Language

This section describes the OpenCL C programming language used to create kernels that are executed on OpenCL device(s). The OpenCL C programming language (also referred to as OpenCL C) is based on the ISO/IEC 9899:1999 C language specification (a.k.a. C99 specification) with specific extensions and restrictions. Please refer to the ISO/IEC 9899:1999 specification for a detailed description of the language grammar. This section describes modifications and restrictions to ISO/IEC 9899:1999 supported in OpenCL C.

### 6.1 Supported Data Types

The following data types are supported.

### 6.1.1 Built-in Scalar Data Types

Table 6.1 describes the list of built-in scalar data types.

| Type | Description |
| :--- | :--- |
| bool $^{\mathbf{1}}$ | A conditional data type which is either true or false. The value true <br> expands to the integer constant 1 and the value false expands to the <br> integer constant 0. |
| char | A signed two's complement 8-bit integer. |
| unsigned char, <br> uchar | An unsigned 8-bit integer. |
| short | A signed two's complement 16-bit integer. |
| unsigned short, <br> ushort | An unsigned 16-bit integer. |
| int | A signed two's complement 32-bit integer. |
| unsigned int, <br> uint | An unsigned 32-bit integer. |
| long | A signed two's complement 64-bit integer. |
| unsigned long, <br> ulong | An unsigned 64-bit integer. |
| float | A 32-bit floating-point. The float data type must conform to the <br> IEEE 754 single precision storage format. |
| double ${ }^{2}$ | A 64-bit floating-point. The double data type must conform to the <br> IEEE 754 double precision storage format. |
| half | A 16-bit floating-point. The half data type must conform to the |

[^0]|  | IEEE 754-2008 half precision storage format. |
| :--- | :--- |
| size_t | The unsigned integer type of the result of the sizeof operator. This <br> is a 32-bit unsigned integer if CL_DEVICE_ADDRESS_BITS <br> defined in table 4.3 is 32-bits and is a 64-bit unsigned integer if <br> CL_DEVICE_ADDRESS_BITS is 64-bits. |
| ptrdiff_t | A signed integer type that is the result of subtracting two pointers. <br> This is a 32-bit signed integer if CL_DEVICE_ADDRESS_BITS <br> defined in table 4.3 is 32-bits and is a 64-bit signed integer if <br> CL_DEVICE_ADDRESS_BITS is 64-bits. |
| intptr_t | A signed integer type with the property that any valid pointer to <br> void can be converted to this type, then converted back to pointer <br> to void, and the result will compare equal to the original pointer. <br> This is a 32-bit signed integer if CL_DEVICE_ADDRESS_BITS <br> defined in table 4.3 is 32-bits and is a 64-bit signed integer if <br> CL_DEVICE_ADDRESS_BITS is 64-bits. |
| uintptr_t | An unsigned integer type with the property that any valid pointer to <br> void can be converted to this type, then converted back to pointer <br> to void, and the result will compare equal to the original pointer. <br> This is a 32-bit signed integer if CL_DEVICE_ADDRESS_BITS <br> defined in table 4.3 is 32-bits and is a 64-bit signed integer if <br> CL_DEVICE_ADDRESS_BITS is 64-bits. |
| void | The void type comprises an empty set of values; it is an incomplete <br> type that cannot be completed. |

Table 6.1 Built-in Scalar Data Types
Most built-in scalar data types are also declared as appropriate types in the OpenCL API (and header files) that can be used by an application. The following table describes the built-in scalar data type in the OpenCL C programming language and the corresponding data type available to the application:

| Type in OpenCL Language | API type for application |
| :---: | :---: |
| bool | n/a |
| char | cl_char |
| unsigned char, <br> uchar | cl_uchar |
| short | cl_short |
| unsigned short, <br> ushort | cl_ushort |
| int | cl_int |
| unsigned int, <br> uint | cl_uint |
| long | cl_long |
| unsigned long, | cl_ulong |


| ulong |  |
| :---: | :---: |
| float | cl_float |
| double | cl_double |
| half | cl_half |
| size_t | n/a |
| ptrdiff_t | n/a |
| intptr_t | n/a |
| uintptr_t | n/a |
| void | void |

### 6.1.1.1 The half data type

The half data type must be IEEE 754-2008 compliant. half numbers have 1 sign bit, 5 exponent bits, and 10 mantissa bits. The interpretation of the sign, exponent and mantissa is analogous to IEEE 754 floating-point numbers. The exponent bias is 15 . The half data type must represent finite and normal numbers, denormalized numbers, infinities and NaN . Denormalized numbers for the half data type which may be generated when converting a float to a half using vstore_half and converting a half to a float using vload_half cannot be flushed to zero. Conversions from float to half correctly round the mantissa to 11 bits of precision. Conversions from half to $f l o a t$ are lossless; all half numbers are exactly representable as float values.

The half data type can only be used to declare a pointer to a buffer that contains half values. A few valid examples are given below:

```
void
bar (__global half *p)
{
}
__kernel void
foo (___global half *pg, __local half *pl)
{
    __global half *ptr;
    int offset;
    ptr = pg + offset;
    bar(ptr);
}
```

Below are some examples that are not valid usage of the half type:

```
half a;
half b[100];
```

```
half *p;
a = *p; \leftarrow not allowed. must use vload_half function
```

Loads from a pointer to a half and stores to a pointer to a half can be performed using the vload_half, vload_half $\boldsymbol{n}$, vloada_halfn and vstore_half, vstore_half $\boldsymbol{n}$, vstorea_halfn functions respectively as described in section 6.13.7. The load functions read scalar or vector half values from memory and convert them to a scalar or vector float value. The store functions take a scalar or vector float value as input, convert it to a half scalar or vector value (with appropriate rounding mode) and write the half scalar or vector value to memory.

### 6.1.2 Built-in Vector Data Types ${ }^{3}$

The char, unsigned char, short, unsigned short, integer, unsigned integer, long, unsigned long, float vector data types are supported. The vector data type is defined with the type name i.e. char, uchar, short, ushort, int, uint, float, long, ulong followed by a literal value $n$ that defines the number of elements in the vector. Supported values of $n$ are $2,3,4,8$, and 16 for all vector data types.

Table 6.2 describes the list of built-in vector data types.

| Type | Description |
| :--- | :--- |
| char $n$ | A vector of $n$ 8-bit signed two's complement integer values. |
| uchar $n$ | A vector of $n$ 8-bit unsigned integer values. |
| short $\boldsymbol{n}$ | A vector of $n$ 16-bit signed two's complement integer values. |
| ushort $\boldsymbol{n}$ | A vector of $n$ 16-bit unsigned integer values. |
| int $\boldsymbol{n}$ | A vector of $n$ 32-bit signed two's complement integer values. |
| uint $\boldsymbol{n}$ | A vector of $n$ 32-bit unsigned integer values. |
| long $\boldsymbol{n}$ | A vector of $n$ 64-bit signed two's complement integer values. |
| ulong $\boldsymbol{n}$ | A vector of $n$ 64-bit unsigned integer values. |
| float $\boldsymbol{n}$ | A vector of $n$ 32-bit floating-point values. |
| double $\boldsymbol{n}^{4}$ | A vector of $n$ 64-bit floating-point values. |

Table 6.2 Built-in Vector Data Types
The built-in vector data types are also declared as appropriate types in the OpenCL API (and header files) that can be used by an application. The following table describes the built-in vector

[^1]data type in the OpenCL C programming language and the corresponding data type available to the application:

| Type in OpenCL Language | API type for application |
| :---: | :---: |
| charn | cl_charn |
| ucharn | cl_ucharn |
| short $n$ | cl_short $n$ |
| ushortn | cl_ushort $n$ |
| int $n$ | cl_int $n$ |
| uint $n$ | cl_uint $n$ |
| long $n$ | cl_longn |
| ulong $n$ | cl_ulong $n$ |
| float $n$ | cl_float $n$ |
| double $n$ | cl_doublen |

### 6.1.3 Other Built-in Data Types

Table 6.3 describes the list of additional data types supported by OpenCL.

| Type | Description |
| :--- | :--- |
| image2d_t | A 2D image. Refer to section 6.13 .14 for a detailed <br> description of the built-in functions that use this type. |
| image3d_t | A 3D image. Refer to section 6.13.14 for a detailed <br> description of the built-in functions that use this type. |
| image2d_array_t | A 2D image array. Refer to section 6.13.14 for a detailed <br> description of the built-in functions that use this type. |
| image1d_t | A 1D image. Refer to section 6.13.14 for a detailed <br> description of the built-in functions that use this type. |
| image1d_buffer_t | A 1D image created from a buffer object. Refer to section <br> 6.13.14 for a detailed description of the built-in functions <br> that use this type. |
| image1d_array_t | A 1D image array. Refer to section 6.13.14 for a detailed <br> description of the built-in functions that use this type. |
| image2d_depth_t | A 2D depth image. Refer to section 6.13.14 for a detailed <br> description of the built-in functions that use this type. |
| image2d_array_depth_t | A 2D depth image array. Refer to section 6.13.14 for a <br> detailed description of the built-in functions that use this <br> type. |
| sampler_t | A sampler type. Refer to section 6.13.14 for a detailed <br> description the built-in functions that use of this type. |
| queue_t | A device command queue. This queue can only be used to <br> enqueue commands from kernels executing on the device. |


| ndrange_t | The N-dimensional range over which a kernel executes. |
| :--- | :--- |
| clk_event_t | A device side event that identifies a command enqueue to a <br> device command queue. |
| reserve_id_t | A reservation ID. This opaque type is used to identify the <br>  <br>  <br>  <br>  <br> reservation for reading and writing a pipe. Refer to section <br> 6.13 .16. |
| event_t | An event. This can be used to identify async copies from <br> global to local memory and vice-versa. Refer to section |
|  | 6.13 .10. |

Table 6.3 Other Built-in Data Types
NOTE: The image2d_t, image3d_t, image2d_array_t, image1d_t, image1d_buffer_t, image1d_array_t, image2d_depth_t, image2d_array_depth_t and sampler_t types are only defined if the device supports images i.e. CL_DEVICE_IMAGE_SUPPORT as described in table 4.3 is CL_TRUE.

The C99 derived types (arrays, structs, unions, functions, and pointers), constructed from the built-in data types described in sections 6.1.1, 6.1.2 and 6.1.3 are supported, with restrictions described in section 6.9.

The following tables describe the other built-in data types in OpenCL described in table 6.3 and the corresponding data type available to the application:

| Type in OpenCL C | API type for application |
| :---: | :---: |
| queue $t$ | cl_command queue |
| clk_event_t | cl_event |

### 6.1.4 Reserved Data Types

The data type names described in table 6.4 are reserved and cannot be used by applications as type names. The vector data type names defined in table 6.2, but where $n$ is any value other than $2,3,4,8$ and 16 , are also reserved.
Type $\quad$ Description

| booln $\boldsymbol{n}$ | A boolean vector. |
| :--- | :--- |
| half $\boldsymbol{n}$ | A 16-bit floating-point vector. |
| quad, quad $\boldsymbol{n}$ | A 128-bit floating-point scalar and vector. |
| complex half, <br> complex half $\boldsymbol{n}$ | A complex 16-bit floating-point scalar and <br> vector. |
| imaginary half, <br> imaginary half $\boldsymbol{n}$ | An imaginary 16-bit floating-point scalar and <br> vector. |
| complex float, <br> complex float $\boldsymbol{n}$ | A complex 32-bit floating-point scalar and <br> vector. |
| imaginary float, <br> imaginary float $\boldsymbol{n}$ | An imaginary 32-bit floating-point scalar and <br> vector. |
| complex double, <br> complex doublen, | A complex 64-bit floating-point scalar and <br> vector. |
| imaginary double, <br> imaginary doublen | An imaginary 64-bit floating-point scalar and <br> vector. |
| complex quad, <br> complex quad $\boldsymbol{n}$, | A complex 128-bit floating-point scalar and <br> vector. |
| imaginary quad, <br> imaginary quad $\boldsymbol{n}$ | An imaginary 128-bit floating-point scalar and <br> vector. |
| float $\boldsymbol{x} \boldsymbol{m}$ | An $n \times m$ matrix of single precision floating- <br> point values stored in column-major order. |
| doublenx $\boldsymbol{m}$ | An $n$ x $m$ matrix of double precision floating- <br> point values stored in column-major order. |
| long double <br> long double $\boldsymbol{n}$ | A floating-point scalar and vector type with at <br> least as much precision and range as a double <br> and no more precision and range than a quad. |
| long long, long longn | A 128-bit signed integer scalar and vector. |
| unsigned long long, <br> ulong long, ulong long $\boldsymbol{n}$ | 128-bit unsigned integer scalar and vector. |

Table 6.4 Reserved Data Types

### 6.1.5 Alignment of Types

A data item declared to be a data type in memory is always aligned to the size of the data type in bytes. For example, a float4 variable will be aligned to a 16-byte boundary, a char2 variable will be aligned to a 2-byte boundary.

For 3-component vector data types, the size of the data type is $4 *$ sizeof(component). This means that a 3 -component vector data type will be aligned to a $4 *$ sizeof(component) boundary. The vload3 and vstore 3 built-in functions can be used to read and write, respectively, 3 -component vector data types from an array of packed scalar data type.

A built-in data type that is not a power of two bytes in size must be aligned to the next larger
power of two. This rule applies to built-in types only, not structs or unions.
The OpenCL compiler is responsible for aligning data items to the appropriate alignment as required by the data type. For arguments to a __kernel function declared to be a pointer to a data type, the OpenCL compiler can assume that the pointee is always appropriately aligned as required by the data type. The behavior of an unaligned load or store is undefined, except for the vload $\boldsymbol{n}$, vload_half $\boldsymbol{n}$, vstoren, and vstore_half $\boldsymbol{n}$ functions defined in section 6.13.7. The vector load functions can read a vector from an address aligned to the element type of the vector. The vector store functions can write a vector to an address aligned to the element type of the vector.

### 6.1.6 Vector Literals

Vector literals can be used to create vectors from a list of scalars, vectors or a mixture thereof. A vector literal can be used either as a vector initializer or as a primary expression. A vector literal cannot be used as an L-value.

A vector literal is written as a parenthesized vector type followed by a parenthesized comma delimited list of parameters. A vector literal operates as an overloaded function. The forms of the function that are available is the set of possible argument lists for which all arguments have the same element type as the result vector, and the total number of elements is equal to the number of elements in the result vector. In addition, a form with a single scalar of the same type as the element type of the vector is available. For example, the following forms are available for float4:

```
(float4)( float, float, float, float )
(float4)( float2, float, float )
(float4)( float, float2, float )
(float4)( float, float, float2 )
(float4)( float2, float2 )
(float4)( float3, float )
(float4)( float, float3 )
(float4)( float )
```

Operands are evaluated by standard rules for function evaluation, except that implicit scalar widening shall not occur. The order in which the operands are evaluated is undefined. The operands are assigned to their respective positions in the result vector as they appear in memory order. That is, the first element of the first operand is assigned to result.x, the second element of the first operand (or the first element of the second operand if the first operand was a scalar) is assigned to result.y, etc. In the case of the form that has a single scalar operand, the operand is replicated across all lanes of the vector.

Examples:

```
float4 f = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
```

```
uint4 u = (uint4)(1); < u will be (1, 1, 1, 1).
float4 f = (float4)((float2)(1.0f, 2.0f),
    (float2)(3.0f, 4.0f));
float4 f = (float4)(1.0f, (float2)(2.0f, 3.0f), 4.0f);
float4 f = (float4)(1.0f, 2.0f); < error
```


### 6.1.7 Vector Components

The components of vector data types with $1 \ldots 4$ components can be addressed as <vector_data_type>.xyzw. Vector data types of type char2, uchar2, short2, ushort2, int2, uint2, long2, ulong2, and float2 can access.xy elements. Vector data types of type char3, uchar3, short3, ushort3, int3, uint3, long3, ulong3, and float3 can access .xyz elements. Vector data types of type char4, uchar4, short4, ushort4, int4, uint4, long4, ulong4, float4 can access .xyzw elements.

Accessing components beyond those declared for the vector type is an error so, for example:

```
float2 pos;
pos.x = 1.Of; // is legal
pos.z = 1.Of; // is illegal
float3 pos;
pos.z = 1.Of; // is legal
pos.w = 1.Of; // is illegal
```

The component selection syntax allows multiple components to be selected by appending their names after the period (.).

```
float4 c;
c.xyzw = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
c.z = 1.0f;
c.xy = (float2)(3.0f, 4.0f);
c.xyz = (float3)(3.0f, 4.0f, 5.0f);
```

The component selection syntax also allows components to be permuted or replicated.

```
float4 pos = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
float4 swiz= pos.wzyx; // swiz = (4.0f, 3.0f, 2.0f, 1.0f)
```

```
float4 dup = pos.xxyy; // dup = (1.0f, 1.0f, 2.0f, 2.0f)
```

The component group notation can occur on the left hand side of an expression. To form an 1value, swizzling must be applied to an l-value of vector type, contain no duplicate components, and it results in an l-value of scalar or vector type, depending on number of components specified. Each component must be a supported scalar or vector type.

```
float4 pos = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
pos.xw = (float2)(5.0f, 6.0f);// pos = (5.0f, 2.0f, 3.0f, 6.0f)
pos.wx = (float2)(7.0f, 8.0f);// pos = (8.0f, 2.0f, 3.0f, 7.0f)
pos.xyz = (float3)(3.0f, 5.0f, 9.0f); // pos = (3.0f, 5.0f, 9.0f, 4.0f)
pos.xx = (float2)(3.0f, 4.0f);// illegal - 'x' used twice
// illegal - mismatch between float2 and float4
pos.xy = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
float4 a, b, c, d;
float16 x;
x = (float16)(a, b, c, d);
x = (float16)(a.xxxx, b.xyz, c.xyz, d.xyz, a.yzw);
// illegal - component a.xxxxxxx is not a valid vector type
x = (float16)(a.xxxxxxx, b.xyz, c.xyz, d.xyz);
```

Elements of vector data types can also be accessed using a numeric index to refer to the appropriate element in the vector. The numeric indices that can be used are given in the table below:

| Vector Components | Numeric indices that can be used |
| :---: | :---: |
| 2-component | 0, 1 |
| 3-component | 0, 1, 2 |
| 4-component | 0, 1, 2, 3 |
| 8-component | 0, 1, 2, 3, 4, 5, 6, 7 |
| 16-component | $\begin{aligned} & 0,1,2,3,4,5,6,7, \\ & 8,9, a, A, b, B, C, C, d, D, e, E, \\ & f, F \end{aligned}$ |

Table 6.5 Numeric indices for built-in vector data types

The numeric indices must be preceded by the letter $s$ or $S$.
In the following example

```
float8 f;
```

f.s 0 refers to the $1^{\text {st }}$ element of the $f l o a t 8$ variable $f$ and $f . s 7$ refers to the $8^{\text {th }}$ element of the float 8 variable $f$.

In the following example

```
float16 x;
```

x.sa (or x.sA) refers to the $11^{\text {th }}$ element of the float 16 variable $x$ and $x . s f$ (or $\mathrm{x} . \mathrm{sF}$ ) refers to the $16^{\text {th }}$ element of the float16 variable x .

The numeric indices used to refer to an appropriate element in the vector cannot be intermixed with . XYZw notation used to access elements of a $1 . .4$ component vector.

For example

```
float4 f, a;
a = f.x12w; // illegal use of numeric indices with .xyzw
a.xyzw = f.s0123; // valid
```

Vector data types can use the .lo (or .even) and .hi (or .odd) suffixes to get smaller vector types or to combine smaller vector types to a larger vector type. Multiple levels of .lo (or . even) and . hi (or . odd) suffixes can be used until they refer to a scalar term.

The. lo suffix refers to the lower half of a given vector. The . hi suffix refers to the upper half of a given vector.

The . even suffix refers to the even elements of a vector. The . odd suffix refers to the odd elements of a vector.

Some examples to help illustrate this are given below:

```
float4 vf;
float2 low = vf.lo; // returns vf.xy
float2 high = vf.hi; // returns vf.zw
float2 even = vf.even; // returns vf.xz
float2 odd = vf.odd; // returns vf.yw
```

The suffixes .lo (or .even) and .hi (or .odd) for a 3-component vector type operate as if the 3 -component vector type is a 4-component vector type with the value in the w component undefined.

## Some examples are given below:

```
float8 vf;
float4 odd = vf.odd;
float4 even = vf.even;
float2 high = vf.even.hi;
float2 low = vf.odd.lo;
// interleave L+R stereo stream
float4 left, right;
float8 interleaved;
interleaved.even = left;
interleaved.odd = right;
// deinterleave
left = interleaved.even;
right = interleaved.odd;
// transpose a 4x4 matrix
void transpose( float4 m[4] )
{
    // read matrix into a float16 vector
    float16 x = (float16)( m[0], m[1], m[2], m[3] );
    float16 t;
    //transpose
    t.even = x.lo;
    t.odd = x.hi;
    x.even = t.lo;
    x.odd = t.hi;
    //write back
    m[0] = x.lo.lo; // { m[0][0], m[1][0], m[2][0], m[3][0] }
    m[1] = x.lo.hi; // { m[0][1], m[1][1], m[2][1], m[3][1] }
    m[2] = x.hi.lo; // { m[0][2], m[1][2], m[2][2], m[3][2] }
    m[3] = x.hi.hi; // { m[0][3], m[1][3], m[2][3], m[3][3] }
}
float3 vf = (float3)(1.0f, 2.0f, 3.0f);
float2 low = vf.lo; // (1.0f, 2.0f);
float2 high = vf.hi; // (3.0f, undefined);
```

It is an error to take the address of a vector element and will result in a compilation error. For example:

```
float8 vf;
float *f = &vf.x; // is illegal
float2 *f2 = &vf.s07; // is illegal
float4 *odd = &vf.odd; // is illegal
float4 *even = &vf.even; // is illegal
float2 *high = &vf.even.hi; // is illegal
float2 *low = &vf.odd.lo; // is illegal
```


### 6.1.8 Aliasing Rules

OpenCL C programs shall comply with the C99 type-based aliasing rules (defined in section 6.5, item 7 of the C99 specification). The OpenCL C built-in vector data types are considered aggregate ${ }^{5}$ types for the purpose of applying these aliasing rules.

### 6.1.9 Keywords

The following names are reserved for use as keywords in OpenCL C and shall not be used otherwise.

Names reserved as keywords by C99.
4 OpenCL C data types defined in tables 6.2, 6.3 and 6.4.
4 Address space qualifiers: __global, global, __local, local, __constant, constant, __private, and private. __generic and generic are reserved for future use.

* Function qualifiers: __kernel and kernel.
* Access qualifiers: __read_only, read_only, __write_only, write_only, __read_write and read_write.
uniform, pipe.

[^2]
### 6.2 Conversions and Type Casting

### 6.2.1 Implicit Conversions

Implicit conversions between scalar built-in types defined in table 6.1 (except void and half $f^{6}$ ) are supported. When an implicit conversion is done, it is not just a re-interpretation of the expression's value but a conversion of that value to an equivalent value in the new type. For example, the integer value 5 will be converted to the floating-point value 5.0 .

Implicit conversions from a scalar type to a vector type are allowed. In this case, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector. The scalar type is then widened to the vector.

Implicit conversions between built-in vector data types are disallowed.
Implicit conversions for pointer types follow the rules described in the C 99 specification.

### 6.2.2 Explicit Casts

Standard typecasts for built-in scalar data types defined in table 6.1 will perform appropriate conversion (except void and half $\mathrm{f}^{7}$ ). In the example below:

```
float f = 1.0f;
int i = (int)f;
```

f stores $0 \times 3 \mathrm{~F} 800000$ and i stores $0 \times 1$ which is the floating-point value 1.0 f in f converted to an integer value.

Explicit casts between vector types are not legal. The examples below will generate a compilation error.

```
int4 i;
uint4 u = (uint4) i; < not allowed
float4 f;
int4 i = (int4) f; < not allowed
float4 f;
int8 i = (int8) f; & not allowed
```

Scalar to vector conversions may be performed by casting the scalar to the desired vector data

[^3]type. Type casting will also perform appropriate arithmetic conversion. The round to zero rounding mode will be used for conversions to built-in integer vector types. The default rounding mode will be used for conversions to floating-point vector types. When casting a bool to a vector integer data type, the vector components will be set to -1 (i.e. all bits set) if the bool value is true and 0 otherwise.

Below are some correct examples of explicit casts.

```
float f = 1.0f;
float4 va = (float4)f;
// va is a float4 vector with elements (f, f, f, f).
uchar u = 0xFF;
float4 vb = (float4)u;
// vb is a float4 vector with elements((float)u, (float)u,
// (float)u, (float)u).
float f = 2.0f;
int2 vc = (int2)f;
// vc is an int2 vector with elements ((int)f, (int)f).
uchar4 vtrue = (uchar4)true;
// vtrue is a uchar4 vector with elements (Oxff, Oxff,
// 0xff, Oxff).
```


### 6.2.3 Explicit Conversions

Explicit conversions may be performed using the

```
convert_destType(sourceType)
```

suite of functions. These provide a full set of type conversions between supported types (see sections 6.1.1, 6.1.2 and 6.1.3) except for the following types: bool, half, size_t, ptrdiff_t, intptr_t, uintptr_t, and void.

The number of elements in the source and destination vectors must match.
In the example below:

```
uchar4 u;
int4 c = convert_int4(u);
```

convert_int 4 converts a uchar 4 vector $u$ to an int 4 vector $c$.

```
    float f;
    int i = convert_int(f);
convert_int converts a float scalar f to an int scalar i.
```

The behavior of the conversion may be modified by one or two optional modifiers that specify saturation for out-of-range inputs and rounding behavior.

The full form of the scalar convert function is:

```
destType convert_destType<_sat><_roundingMode> (sourceType)
```

The full form of the vector convert function is:

```
destTypen convert_destTypen<_sat><_roundingMode> (sourceTypen)
```


### 6.2.3.1 Data Types

Conversions are available for the following scalar types: char, uchar, short, ushort, int, uint, long, ulong, float, and built-in vector types derived therefrom. The operand and result type must have the same number of elements. The operand and result type may be the same type in which case the conversion has no effect on the type or value of an expression.

Conversions between integer types follow the conversion rules specified in sections 6.3.1.1 and 6.3.1.3 of the C99 specification except for out-of-range behavior and saturated conversions which are described in section 6.2.3.3 below.

### 6.2.3.2 Rounding Modes

Conversions to and from floating-point type shall conform to IEEE-754 rounding rules. Conversions may have an optional rounding mode modifier described in table 6.6.

| Modifier | Rounding Mode Description |
| :--- | :--- |
| $-r t e$ | Round to nearest even |
| $-r t z$ | Round toward zero |
| $-r t p$ | Round toward positive infinity |
| rtn | Round toward negative infinity |
| no modifier specified | Use the default rounding mode for this destination type, <br> rtz for conversion to integers or the default rounding <br> mode for conversion to floating-point types. |

Table 6.6 Rounding Modes

By default, conversions to integer type use the _rtz (round toward zero) rounding mode and conversions to floating-point type ${ }^{8}$ use the default rounding mode. The only default floatingpoint rounding mode supported is round to nearest even i.e the default rounding mode will be _rte for floating-point types.

### 6.2.3.3 Out-of-Range Behavior and Saturated Conversions

When the conversion operand is either greater than the greatest representable destination value or less than the least representable destination value, it is said to be out-of-range. The result of out-of-range conversion is determined by the conversion rules specified by the C99 specification in section 6.3. When converting from a floating-point type to integer type, the behavior is implementation-defined.

Conversions to integer type may opt to convert using the optional saturated mode by appending the _sat modifier to the conversion function name. When in saturated mode, values that are outside the representable range shall clamp to the nearest representable value in the destination format. ( NaN should be converted to 0 ).

Conversions to floating-point type shall conform to IEEE-754 rounding rules. The _sat modifier may not be used for conversions to floating-point formats.

### 6.2.3.4 Explicit Conversion Examples

Example 1:

```
short4 s;
// negative values clamped to 0
ushort4 u = convert_ushort4_sat( s );
// values > CHAR_MAX converted to CHAR_MAX
// values < CHAR_MIN converted to CHAR_MIN
char4 c = convert_char4_sat( s );
```


## Example 2:

```
float4 f;
// values implementation defined for
// f > INT MAX, f < INT MIN or NaN
int4 i = convert_int4(` f );
```

[^4]```
// values > INT_MAX clamp to INT_MAX, values < INT_MIN clamp
// to INT_MIN. NaN should produce 0.
// The _rtz rounding mode is used to produce the integer values.
int4 i}2= convert_int4_sat( f );
// similar to convert_int4, except that floating-point values
// are rounded to the nearest integer instead of truncated
int4 i3 = convert_int4_rte( f );
// similar to convert_int4_sat, except that floating-point values
// are rounded to the nearest integer instead of truncated
int4 i4 = convert_int4_sat_rte( f );
```

Example 3:

```
int4 i;
```

// convert ints to floats using the default rounding mode.
float4 $\mathbf{f}=$ convert_float4 ( i ) ;
// convert ints to floats. integer values that cannot
// be exactly represented as floats should round up to the
// next representable float.
float4 $\mathbf{f}=$ convert_float4_rtp ( i ) ;

### 6.2.4 Reinterpreting Data As Another Type

It is frequently necessary to reinterpret bits in a data type as another data type in OpenCL. This is typically required when direct access to the bits in a floating-point type is needed, for example to mask off the sign bit or make use of the result of a vector relational operator (see section 6.3.d) on floating-point data ${ }^{9}$. Several methods to achieve this (non-) conversion are frequently practiced in C, including pointer aliasing, unions and memcpy. Of these, only memcpy is strictly correct in C99. Since OpenCL does not provide memcpy, other methods are needed.

### 6.2.4.1 Reinterpreting Types Using Unions

The OpenCL language extends the union to allow the program to access a member of a union object using a member of a different type. The relevant bytes of the representation of the object

[^5]are treated as an object of the type used for the access. If the type used for access is larger than the representation of the object, then the value of the additional bytes is undefined.

Examples:

```
union{ float f; uint u; double d 10;} u;
u.u = 1; // u.f contains 2**-149. u.d is undefined --
    // depending on endianness the low or high half
    // of d is unknown
u.f = 1.0f; // u.u contains 0x3f800000, u.d contains an
    // undefined value -- depending on endianness
    // the low or high half of d is unknown
u.d = 1.0; // u.u contains 0x3ff00000 (big endian) or 0
    // (little endian). u.f contains either 0x1.ep0f
    // (big endian) or 0.0f (little endian)
```


### 6.2.4.2 Reinterpreting Types Using as_type() and as_typen()

All data types described in tables 6.1 and 6.2 (except bool, half ${ }^{11}$ and void) may be also reinterpreted as another data type of the same size using the as_type() operator for scalar data types and the as_typen () operator ${ }^{12}$ for vector data types. When the operand and result type contain the same number of elements, the bits in the operand shall be returned directly without modification as the new type. The usual type promotion for function arguments shall not be performed.

For example, as_float $(0 \times 3 f 800000)$ returns $1.0 f$, which is the value that the bit pattern $0 \times 3 £ 800000$ has if viewed as an IEEE-754 single precision value.

When the operand and result type contain a different number of elements, the result shall be implementation-defined except if the operand is a 4-component vector and the result is a 3component vector. In this case, the bits in the operand shall be returned directly without

[^6]modification as the new type. That is, a conforming implementation shall explicitly define a behavior, but two conforming implementations need not have the same behavior when the number of elements in the result and operand types does not match. The implementation may define the result to contain all, some or none of the original bits in whatever order it chooses. It is an error to use as_type() or as_typen() operator to reinterpret data to a type of a different number of bytes.

## Examples:

```
float f = 1.0f;
uint u = as_uint(f); // Legal. Contains: 0x3f800000
float4 f = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
// Legal. Contains:
// (int4)(0x3f800000, 0x40000000, 0x40400000, 0x40800000)
int4 i = as_int4(f);
float4 f, g;
int4 is_less = f < g;
// Legal. f[i] = f[i] < g[i] ? f[i] : 0.0f
f = as_float4(as_int4(f) & is_less);
int i;
// Legal. Result is implementation-defined.
short2 j = as_short2(i);
int4 i;
// Legal. Result is implementation-defined.
short8 j = as_short8(i);
float4 f;
// Error. Result and operand have different sizes
double4 g = as_double4 }\mp@subsup{}{}{13}(f)
float4 f;
// Legal. g.xyz will have same values as f.xyz. g.w is undefined
float3 g = as_float3(f);
```


### 6.2.5 Pointer Casting

Pointers to old and new types may be cast back and forth to each other. Casting a pointer to a new type represents an unchecked assertion that the address is correctly aligned. The developer will also need to know the endianness of the OpenCL device and the endianness of the data to determine how the scalar and vector data elements are stored in memory.

[^7]
### 6.2.6 Usual Arithmetic Conversions

Many operators that expect operands of arithmetic type cause conversions and yield result types in a similar way. The purpose is to determine a common real type for the operands and result. For the specified operands, each operand is converted, without change of type domain, to a type whose corresponding real type is the common real type. For this purpose, all vector types shall be considered to have higher conversion ranks than scalars. Unless explicitly stated otherwise, the common real type is also the corresponding real type of the result, whose type domain is the type domain of the operands if they are the same, and complex otherwise. This pattern is called the usual arithmetic conversions. If the operands are of more than one vector type, then an error shall occur. Implicit conversions between vector types are not permitted, per section 6.2.1.

Otherwise, if there is only a single vector type, and all other operands are scalar types, the scalar types are converted to the type of the vector element, then widened into a new vector containing the same number of elements as the vector, by duplication of the scalar value across the width of the new vector. An error shall occur if any scalar operand has greater rank than the type of the vector element. For this purpose, the rank order defined as follows:

1. The rank of a floating-point type is greater than the rank of another floating-point type, if the first floating-point type can exactly represent all numeric values in the second floating-point type. (For this purpose, the encoding of the floating-point value is used, rather than the subset of the encoding usable by the device.)
2. The rank of any floating-point type is greater than the rank of any integer type.
3. The rank of an integer type is greater than the rank of an integer type with less precision.
4. The rank of an unsigned integer type is greater than the rank of a signed integer type with the same precision. ${ }^{14}$
5. The rank of the bool type is less than the rank of any other type.
6. The rank of an enumerated type shall equal the rank of the compatible integer type.
7. For all types, $\mathrm{T} 1, \mathrm{~T} 2$ and T 3 , if T 1 has greater rank than T 2 , and T 2 has greater rank than T 3 , then T 1 has greater rank than T 3 .

Otherwise, if all operands are scalar, the usual arithmetic conversions apply, per section 6.3.1.8 of the C99 standard.

NOTE: Both the standard orderings in sections 6.3.1.8 and 6.3.1.1 of C99 were examined and rejected. Had we used integer conversion rank here, int $4+$ OU would have been legal and had int 4 return type. Had we used standard C99 usual arithmetic conversion rules for scalars, then the standard integer promotion would have been performed on vector integer element types and short8 + char would either have return type of int8 or be illegal.

[^8]
### 6.3 Operators

a. The arithmetic operators add $(+)$, subtract $(-)$, multiply $(*)$ and divide ( $/$ ) operate on built-in integer and floating-point scalar, and vector data types. The remainder (\%) operates on builtin integer scalar and integer vector data types. All arithmetic operators return result of the same built-in type (integer or floating-point) as the type of the operands, after operand type conversion. After conversion, the following cases are valid:

The two operands are scalars. In this case, the operation is applied, resulting in a scalar.
\# One operand is a scalar, and the other is a vector. In this case, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in the same size vector.

The two operands are vectors of the same type. In this case, the operation is done component-wise resulting in the same size vector.

All other cases of implicit conversions are illegal. Division on integer types which results in a value that lies outside of the range bounded by the maximum and minimum representable values of the integer type will not cause an exception but will result in an unspecified value. A divide by zero with integer types does not cause an exception but will result in an unspecified value. Division by zero for floating-point types will result in $\pm$ infinity or NaN as prescribed by the IEEE-754 standard. Use the built-in functions dot and cross to get, respectively, the vector dot product and the vector cross product.
b. The arithmetic unary operators (+ and -) operate on built-in scalar and vector types.
c. The arithmetic post- and pre-increment and decrement operators ( $-\mathbf{-}$ and $+\boldsymbol{+}$ ) operate on built-in scalar and vector types except the built-in scalar and vector float types ${ }^{15}$. All unary operators work component-wise on their operands. These result with the same type they operated on. For post- and pre-increment and decrement, the expression must be one that could be assigned to (an l-value). Pre-increment and pre-decrement add or subtract 1 to the contents of the expression they operate on, and the value of the pre-increment or predecrement expression is the resulting value of that modification. Post-increment and postdecrement expressions add or subtract 1 to the contents of the expression they operate on, but the resulting expression has the expression's value before the post-increment or postdecrement was executed.

[^9]d. The relational operators ${ }^{16}$ greater than ( $>$ ), less than ( $<$ ), greater than or equal ( $>=$ ), and less than or equal $(<=)$ operate on scalar and vector types. All relational operators result in an integer type. After operand type conversion, the following cases are valid:

The two operands are scalars. In this case, the operation is applied, resulting in an int scalar.
\# One operand is a scalar, and the other is a vector. In this case, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in the same size vector.
\$ The two operands are vectors of the same type. In this case, the operation is done component-wise resulting in the same size vector.

All other cases of implicit conversions are illegal.
The result is a scalar signed integer of type int if the source operands are scalar and a vector signed integer type of the same size as the source operands if the source operands are vector types. Vector source operands of type charn and ucharn return a charn result; vector source operands of type short $n$ and ushort $n$ return a short $n$ result; vector source operands of type int $n$, uint $n$ and float $n$ return an int $n$ result; vector source operands of type longn, ulong $n$ and doublen return a longn result. For scalar types, the relational operators shall return 0 if the specified relation is false and 1 if the specified relation is true. For vector types, the relational operators shall return 0 if the specified relation is false and -1 (i.e. all bits set) if the specified relation is true. The relational operators always return 0 if either argument is not a number ( NaN ).
e. The equality operators ${ }^{17}$ equal ( $==$ ), and not equal ( $!=$ ) operate on built-in scalar and vector types. All equality operators result in an integer type. After operand type conversion, the following cases are valid:

The two operands are scalars. In this case, the operation is applied, resulting in a scalar.

One operand is a scalar, and the other is a vector. In this case, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting

[^10]in the same size vector.
The two operands are vectors of the same type. In this case, the operation is done component-wise resulting in the same size vector.

All other cases of implicit conversions are illegal.
The result is a scalar signed integer of type int if the source operands are scalar and a vector signed integer type of the same size as the source operands if the source operands are vector types. Vector source operands of type charn and ucharn return a charn result; vector source operands of type short $n$ and ushort $n$ return a short $n$ result; vector source operands of type int $n$, uint $n$ and float $n$ return an int $n$ result; vector source operands of type longn, ulongn and doublen return a long $n$ result.

For scalar types, the equality operators return 0 if the specified relation is false and return 1 if the specified relation is true. For vector types, the equality operators shall return 0 if the specified relation is false and -1 (i.e. all bits set) if the specified relation is true. The equality operator equal (==) returns 0 if one or both arguments are not a number ( NaN ). The equality operator not equal (!=) returns 1 (for scalar source operands) or -1 (for vector source operands) if one or both arguments are not a number ( NaN ).
f. The bitwise operators and (\&), or (I), exclusive or ( $\wedge$ ), not ( $\sim$ ) operate on all scalar and vector built-in types except the built-in scalar and vector float types. For vector built-in types, the operators are applied component-wise. If one operand is a scalar and the other is a vector, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in the same size vector.
g. The logical operators and (\&\&), or (II) operate on all scalar and vector built-in types. For scalar built-in types only, and $(\& \&)$ will only evaluate the right hand operand if the left hand operand compares unequal to 0 . For scalar built-in types only, or (II) will only evaluate the right hand operand if the left hand operand compares equal to 0 . For built-in vector types, both operands are evaluated and the operators are applied component-wise. If one operand is a scalar and the other is a vector, the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in the same size vector.

The logical operator exclusive or ( ${ }^{\wedge} \wedge$ ) is reserved.
The result is a scalar signed integer of type int if the source operands are scalar and a vector signed integer type of the same size as the source operands if the source operands are vector types. Vector source operands of type charn and ucharn return a charn result; vector source operands of type short $n$ and ushort $n$ return a short $n$ result; vector source
operands of type int $n$, uint $n$ and float $n$ return an int $n$ result; vector source operands of type longn, ulong $n$ and doublen return a longn result.

For scalar types, the logical operators shall return 0 if the result of the operation is false and 1 if the result is true. For vector types, the logical operators shall return 0 if the result of the operation is false and -1 (i.e. all bits set) if the result is true.
h. The logical unary operator not (!) operates on all scalar and vector built-in types. For built-in vector types, the operators are applied component-wise.

The result is a scalar signed integer of type int if the source operands are scalar and a vector signed integer type of the same size as the source operands if the source operands are vector types. Vector source operands of type charn and ucharn return a charn result; vector source operands of type short $n$ and ushort $n$ return a short $n$ result; vector source operands of type int $n$, uint $n$ and float $n$ return an int $n$ result; vector source operands of type longn, ulong $n$ and double $n$ return a long $n$ result.

For scalar types, the result of the logical unary operator is 0 if the value of its operand compares unequal to 0 , and 1 if the value of its operand compares equal to 0 . For vector types, the unary operator shall return a 0 if the value of its operand compares unequal to 0 , and -1 (i.e. all bits set) if the value of its operand compares equal to 0 .
i. The ternary selection operator (? : ) operates on three expressions (exp1 $\boldsymbol{e x p} 2: \exp 3)$. This operator evaluates the first expression $\exp 1$, which can be a scalar or vector result except float. If all three expressions are scalar values, the C99 rules for ternary operator are followed. If the result is a vector value, then this is equivalent to calling $\operatorname{select}(\exp 3, \exp 2$, exp1). The select function is described in table 6.14. The second and third expressions can be any type, as long their types match, or there is a conversion in section 6.2.1 Implicit Conversions that can be applied to one of the expressions to make their types match, or one is a vector and the other is a scalar and the scalar may be subject to the usual arithmetic conversion to the element type used by the vector operand and widened to the same type as the vector type. This resulting matching type is the type of the entire expression.
j. The operators right-shift ( $\gg$ ), left-shift ( $\ll$ ) operate on all scalar and vector built-in types except the built-in scalar and vector float types. For built-in vector types, the operators are applied component-wise. For the right-shift ( $\gg$ ), left-shift ( $\ll$ ) operators, the rightmost operand must be a scalar if the first operand is a scalar, and the rightmost operand can be a vector or scalar if the first operand is a vector.

The result of E1 << E2 is E1 left-shifted by $\log _{2}(\mathrm{~N})$ least significant bits in E2 viewed as an unsigned integer value, where $N$ is the number of bits used to represent the data type of E1 after integer promotion ${ }^{18}$, if E1 is a scalar, or the number of bits used to represent the type of E1 elements, if E1 is a vector. The vacated bits are filled with zeros.

[^11]The result of E1 >> E2 is E1 right-shifted by $\log _{2}(\mathrm{~N})$ least significant bits in E2 viewed as an unsigned integer value, where $N$ is the number of bits used to represent the data type of E1 after integer promotion, if E1 is a scalar, or the number of bits used to represent the type of E1 elements, if E1 is a vector. If E1 has an unsigned type or if E1 has a signed type and a nonnegative value, the vacated bits are filled with zeros. If E1 has a signed type and a negative value, the vacated bits are filled with ones.
k. The sizeof operator yields the size (in bytes) of its operand, including any padding bytes (refer to section 6.1.5) needed for alignment, which may be an expression or the parenthesized name of a type. The size is determined from the type of the operand. The result is of type size_t. If the type of the operand is a variable length array ${ }^{19}$ type, the operand is evaluated; otherwise, the operand is not evaluated and the result is an integer constant.

When applied to an operand that has type char, uchar, the result is 1 . When applied to an operand that has type short, ushort, or half the result is 2 . When applied to an operand that has type int, uint or float, the result is 4. When applied to an operand that has type long, ulong or double, the result is 8 . When applied to an operand that is a vector type, the result ${ }^{20}$ is number of components * size of each scalar component. When applied to an operand that has array type, the result is the total number of bytes in the array. When applied to an operand that has structure or union type, the result is the total number of bytes in such an object, including internal and trailing padding. The sizeof operator shall not be applied to an expression that has function type or an incomplete type, to the parenthesized name of such a type, or to an expression that designates a bit-field struct member ${ }^{21}$.

The behavior of applying the sizeof operator to the bool, image2d_t, image3d_t, image2d_array_t,image2d_depth_t, image2d_array_depth_t, image1d_t, image1d_buffer_t or image1d_array_t, sampler_t, clk_event_t, queue_t and event_t types is implementation-defined.

1. The comma (,) operator operates on expressions by returning the type and value of the rightmost expression in a comma separated list of expressions. All expressions are evaluated, in order, from left to right.
m . The unary $\left({ }^{*}\right)$ operator denotes indirection. If the operand points to an object, the result is an lvalue designating the object. If the operand has type "pointer to type", the result has type "type". If an invalid value has been assigned to the pointer, the behavior of the unary *

[^12]operator is undefined ${ }^{22}$.
n. The unary $(\boldsymbol{\&})$ operator returns the address of its operand. If the operand has type "type", the result has type "pointer to type". If the operand is the result of a unary * operator, neither that operator nor the $\boldsymbol{\&}$ operator is evaluated and the result is as if both were omitted, except that the constraints on the operators still apply and the result is not an lvalue. Similarly, if the operand is the result of a [] operator, neither the $\boldsymbol{\&}$ operator nor the unary * that is implied by the [] is evaluated and the result is as if the $\boldsymbol{\&}$ operator were removed and the [] operator were changed to a + operator. Otherwise, the result is a pointer to the object designated by its operand ${ }^{23}$.
o. Assignments of values to variable names are done with the assignment operator ( = ), like
lvalue = expression

The assignment operator stores the value of expression into lvalue. The expression and lvalue must have the same type, or the expression must have a type in table 6.1, in which case an implicit conversion will be done on the expression before the assignment is done.

If expression is a scalar type and lvalue is a vector type, the scalar is converted to the element type used by the vector operand. The scalar type is then widened to a vector that has the same number of components as the vector operand. The operation is done component-wise resulting in the same size vector.

Any other desired type-conversions must be specified explicitly. L-values must be writable. Variables that are built-in types, entire structures or arrays, structure fields, l-values with the field selector (. ) applied to select components or swizzles without repeated fields, l-values within parentheses, and l-values dereferenced with the array subscript operator ( [ ] ) are all lvalues. Other binary or unary expressions, function names, swizzles with repeated fields, and constants cannot be l-values. The ternary operator (?:) is also not allowed as an l-value.

The order of evaluation of the operands is unspecified. If an attempt is made to modify the result of an assignment operator or to access it after the next sequence point, the behavior is undefined. Other assignment operators are the assignments add into (+=), subtract from $(-=)$, multiply into $\left(^{*}=\right)$, divide into $(/=)$, modulus into ( $\%=$ ), left shift by ( $\ll=$ ), right shift by (>>=), and into ( $\&=$ ), inclusive or into ( $\mid=$ ), and exclusive or into ( $\wedge=$ ).

The expression
lvalue op= expression

[^13]is equivalent to

```
lvalue = lvalue op expression
```

and the l-value and expression must satisfy the requirements for both operator $o p$ and assignment (=).

Note: Except for the sizeof operator, the half data type cannot be used with any of the operators described in this section.

### 6.4 Vector Operations

Vector operations are component-wise. Usually, when an operator operates on a vector, it is operating independently on each component of the vector, in a component-wise fashion.

For example,

```
float4 v, u;
float f;
v = u + f;
```

will be equivalent to

```
v.x = u.x + f;
v.y = u.y + f;
v.z = u.z + f;
v.w = u.w + f;
```

And

```
float4 v, u, w;
w = v + u;
```

will be equivalent to

$$
\begin{aligned}
& \mathrm{w} \cdot \mathrm{x}=\mathrm{v} \cdot \mathrm{x}+\mathrm{u} \cdot \mathrm{x} ; \\
& \mathrm{w} \cdot \mathrm{y}=\mathrm{v} \cdot \mathrm{y}+\mathrm{u} \cdot \mathrm{y} ; \\
& \mathrm{w} \cdot \mathrm{z}=\mathrm{v} \cdot \mathrm{z}+\mathrm{u} \cdot \mathrm{z} ; \\
& \mathrm{w} \cdot \mathrm{w}=\mathrm{v} \cdot \mathrm{w}+\mathrm{u} \cdot \mathrm{w} ;
\end{aligned}
$$

and likewise for most operators and all integer and floating-point vector types.

### 6.5 Address Space Qualifiers

OpenCL implements the following disjoint named address spaces: __global, __local, constant and __private. The address space qualifier may be used in variable declarations to specify the region of memory that is used to allocate the object. The C syntax for type qualifiers is extended in OpenCL to include an address space name as a valid type qualifier. If the type of an object is qualified by an address space name, the object is allocated in the specified address name; otherwise, the object is allocated in the generic address space.

The address space names without the _ prefix i.e. global, local, constant and private may be substituted for the corresponding address space names with the __prefix.

The address space name for arguments to a function in a program, or local variables of a function is __private. All function arguments shall be in the $\qquad$ private address space. The address space for a variable at program scope, a static or extern variable inside a function can either be __global or __constant, but defaults to __global if not specified.

Examples:

```
// declares a pointer p in the private address space that
// points to an object in address space global
global int *p;
void foo (...)
{
    // declares an array of 4 floats in the private address space
    float x[4];
}
```

OpenCL 2.0 adds support for an unnamed generic address space. Pointers that are declared without pointing to a named address space point to the generic address space. Before referring to the region pointed to, the pointer must be associated with a named address space. Functions written with pointer arguments and return values which do not declare an address space are defined to point to the generic address space.
kernel function arguments declared to be a pointer or an array of a type must point to one of the named address spaces __global, __local or __constant.

The named address spaces are a subset of the generic address space except for the constant address space.

A pointer to address space A can only be assigned to a pointer to the same address space A or a pointer to the generic address space. Casting a pointer to address space A to a pointer to address space $B$ is illegal if $A$ and $B$ are named address spaces and $A$ is not the same as $B$.

Examples:

```
private int f() { ... } // should generate an error
local int *f() { ... } // allowed
local int * private f() { ... }; // should generate an error.
```

The $\qquad$ constant, __local, __private, global, constant, local, and private names are reserved for use as address space qualifiers and shall not be used otherwise. The $\qquad$ generic and generic names are reserved for future use.

NOTE: The size of pointers to different address spaces may differ. It is not correct to assume that, for example, sizeof(__global int *) always equals sizeof(__local int *).

### 6.5.1 __global (or global)

The __global or global address space name is used to refer to memory objects (buffer or image objects) allocated from the global memory pool.

A buffer memory object can be declared as a pointer to a scalar, vector or user-defined struct. This allows the kernel to read and/or write any location in the buffer.

The actual size of the array memory object is determined when the memory object is allocated via appropriate API calls in the host code.

Some examples are:

```
global float4 *color; // An array of float4 elements
typedef struct {
    floata[3];
    int b[2];
} foo_t;
global foo_t *my_info; // An array of foo_t elements.
```

As image objects are always allocated from the global address space, the __global or global qualifier should not be specified for image types. The elements of an image object cannot be directly accessed. Built-in functions to read from and write to an image object are provided.

Variables defined at program scope and static variables inside a function can also be declared in the global address space. They can be defined with any valid OpenCL C data type except for those in table 6.3. In particular, such program scope variables may be of any user-defined type, or a pointer to a user-defined type. In the presence of shared virtual memory, these pointers or pointer members should work as expected as long as they are shared virtual memory pointers
and the referenced storage has been mapped appropriately. These variables in the global address space have the same lifetime as the program, and their values persist between calls to any of the kernels in the program. These variables are not shared across devices. They have distinct storage.

Program scope and static variables in the global address space may be initialized, but only with constant expressions.

## Examples:

```
global int foo; // OK.
int foo; // OK. Declared in the global address space
global uchar buf[512]; // OK.
global int baz = 12; // OK. Initialization is allowed
static global int bat; // OK. Internal linkage
static int foo; // OK. Declared in the global address space
static global int foo; // OK.
int *foo; // OK. foo is allocated in global address space.
    // pointer to foo in generic address space
void func(...)
{
    int *foo; // OK. foo is allocated in private address space.
    // foo points to a location in generic address space.
}
global int * global ptr; // OK.
int * global ptr; // OK.
constant int *global ptr=&baz; // error since baz is in global address
    // space.
global int * constant ptr = &baz; // OK
// Pointers work. Also, initialization to a constant known at
// program load time
global int *global baz_ptr = &baz;
global image2d_t im; // Error. Invalid type for program scope
    // variables
global event_t ev; // Error. Invalid type for program scope variables
global int *bad_ptr; // Error. No implicit address space
```

The const qualifier can also be used with the __global qualifier to specify a read-only buffer memory object.

### 6.5.2 __local (or local)

The __local or local address space name is used to describe variables that need to be allocated in local memory and are shared by all work-items of a work-group. Pointers to the local address space are allowed as arguments to functions (including kernel functions). Variables declared in the __local address space inside a kernel function must occur at kernel function scope.

Some examples of variables allocated in the __local address space inside a kernel function are:

```
kernel void my_func(...)
{
    local float a; // A single float allocated
            // in local address space
    local float b[10];// An array of 10 floats
                            // allocated in local address space.
    if (...)
    {
        // example of variable in __local address space but not
        // declared at __kernel function scope.
        local float c; < not allowed.
    }
}
```

Variables allocated in the $\qquad$ local address space inside a kernel function cannot be initialized.

```
kernel void my_func(...)
{
    local float a = 1; < not allowed
    local float b;
    b = 1; < allowed
}
```

NOTE: Variables allocated in the _local address space inside a kernel function are allocated for each work-group executing the kernel and exist only for the lifetime of the work-group executing the kernel.

### 6.5.3 __constant (or constant)

The $\qquad$ constant or constant address space name is used to describe variables allocated in global memory and which are accessed inside a kernel(s) as read-only variables. These readonly variables can be accessed by all (global) work-items of the kernel during its execution.

Pointers to the $\qquad$ constant address space are allowed as arguments to functions (including kernel functions) and for variables declared inside functions.

All string literal storage shall be in the $\qquad$ constant address space.

NOTE: Each argument to a kernel that is a pointer to the $\qquad$ constant address space is counted separately towards the maximum number of such arguments, defined as CL_DEVICE_MAX_CONSTANT_ARGS in table 4.3.

Variables in the program scope can be declared in the $\qquad$ constant address space. Variables in the outermost scope of kernel functions can be declared in the __ constant address space. These variables are required to be initialized and the values used to initialize these variables must be a compile time constant. Writing to such a variable results in a compile-time error.

Implementations are not required to aggregate these declarations into the fewest number of constant arguments. This behavior is implementation defined.

Thus portable code must conservatively assume that each variable declared inside a function or in program scope allocated in the __constant address space counts as a separate constant argument.

### 6.5.4 __private (or private)

Variables inside a kernel function not declared with an address space qualifier, all variables inside non-kernel functions, and all function arguments are in the __private or private address space. Variables declared as pointers are considered to point to the __private address space if an address space qualifier is not specified.

### 6.5.5 The generic address space

The following rules apply when using pointers that point to the generic address space:

- A pointer that points to the global, local or private address space can be implicitly converted to a pointer to the unnamed generic address space but not vice-versa.
- Pointer casts can be used to cast a pointer that points to the global, local or private space to the unnamed generic address space and vice-versa.
- A pointer that points to the constant address space cannot be cast or implicitly converted to the generic address space.

A few examples follow.

This is the canonical example. In this example, function foo is declared with an argument that is a pointer with no address space qualifier.

```
void foo(int *a)
{
    *a = *a + 2;
}
kernel void kl(local int *a)
{
    foo(a);
}
kernel void k2(global int *a)
{
    foo(a);
}
```

In the example below, var is in the unnamed generic address space which gets mapped to the global or local address space depending on the result of the conditional expression.

```
kernel void bar(global int *g, local int *l)
{
    int *var;
    if (is_even(get_global_id(0))
            var = g;
        else
        var = l;
        *var = 42;
}
```

The example below is an example with one unnamed generic address space pointer with multiple named address space assignments.

```
int *ptr;
global int g;
ptr = &g; // legal
local int l;
ptr = &l; // legal
```

```
private int p;
ptr = &p; // legal
constant int c;
ptr = &c; // illegal
```

The example below is an example with one unnamed generic address space pointer being assigned to point to several named address spaces.

```
global int * gp;
local int *lp;
private int *pp;
int *p;
p = gp; // legal
p = lp; // legal
p = pp; // legal
// it is illegal to convert from a generic pointer
// to an explicit address space pointer without a cast:
gp = p; // compile-time error
lp = p; // compile-time error
pp = p; // compile-time error
```


### 6.5.6 Changes to ISO/IEC 9899:1999

This section details the modifications to ISO/IEC 9899:1999 needed to incorporate the functionality of named address space and the generic address space:

## Clause 6.2.5 - Types, replace paragraph 26 with the following paragraphs:

If type $T$ is qualified by the address space qualifier for address space $A$, then " $T$ is in $A$ ". If type $T$ is in address space $A$, a pointer to $T$ is also a "pointer into $A$ " and the referenced address space of the pointer is $A$.

A pointer to void in any address space shall have the same representation and alignment requirements as a pointer to a character type in the same address space. Similarly, pointers to differently access-qualified versions of compatible types shall have the same representation and alignment requirements. All pointers to structure types in the same address space shall have the same representation and alignment requirements as each other. All pointers to union types in the same address space shall have the same representation and alignment requirements as each other.

Clause 6.3.2.3 - Pointers, replace the first two paragraphs with the following paragraphs:
If a pointer into one address space is converted to a pointer into another address space, then
unless the original pointer is a null pointer or the location referred to by the original pointer is within the second address space, the behavior is undefined. (For the original pointer to refer to a location within the second address space, the two address spaces must overlap).

A pointer to void in any address space may be converted to or from a pointer to any incomplete or object type. A pointer to any incomplete or object type in some address space may be converted to a pointer to void in an enclosing address space and back again; the result shall compare equal to the original pointer.

For any qualifier q , a pointer to a non-q-qualified type may be converted to a pointer to the qqualified version of the type (but with the same address-space qualifier or the generic address space); the values stored in the original and converted pointers shall compare equal.

Clause 6.3.2.3 - Pointers, replace the last sentence of paragraph 4 with:
Conversion of a null pointer to another pointer type yields a null pointer of that type. Any two null pointers whose referenced address spaces overlap shall compare equal.

Clause 6.5.2.2 - Function calls, change the second bullet of paragraph 6 to:
both types are pointers to qualified or unqualified versions of a character type or void in the same address space or one type is a pointer in a named address space and the other is a pointer in the generic address space.

## Clause 6.5.6 - Additive operators, add another constraint paragraph:

For subtraction, if the two operands are pointers into different address spaces, the address spaces must overlap.

## Clause 6.5.8 - Relational operators, add another constraint paragraph:

If the two operands are pointers into different address spaces, the address spaces must overlap.

## Clause 6.5.8 - Relational operators, add a new paragraph between existing paragraphs 3 and 4 :

If the two operands are pointers into different address spaces, one of the address spaces encloses the other. The pointer into the enclosed address space is first converted to a pointer to the same reference type except with any address-space qualifier removed and any address-space qualifier of the other pointer's reference type added. (After this conversion, both pointers are pointers into the same address space).

Examples:

```
kernel void test1()
{
```

```
    global int arr[5] = { 0, 1, 2, 3, 4 };
    int *p = &arr[1];
    global int *q = &arr[3];
    // q implicitly converted to the generic address space
    // since the generic address space encloses the global
    // address space
    if (q >= p)
    printf("true\n");
    // q implicitly converted to the generic address space
    // since the generic address space encloses the global
    // address space
    if (p <= q)
    printf("true\n");
}
```


## Clause 6.5.9 - Equality operators, add another constraint paragraph:

If the two operands are pointers into different address spaces, the address spaces must overlap.

## Clause 6.5.9 - Equality operators, replace paragraph 5 with:

Otherwise, at least one operand is a pointer. If one operand is a pointer and the other is a null pointer constant, the null pointer constant is converted to the type of the pointer. If both operands are pointers, each of the following conversions is performed as applicable:

* If the two operands are pointers into different address spaces, one of the address spaces encloses the other. The pointer into the enclosed address space is first converted to a pointer to the same reference type except with any address-space qualifier removed and any address-space qualifier of the other pointer's reference type added. (After this conversion, both pointers are pointers into the same address space).
\$ Then, if one operand is a pointer to an object or incomplete type and the other is a pointer to a qualified or unqualified version of void, the former is converted to the type of the latter.


## Examples:

```
int *ptr = NULL;
local int lval = SOME_VAL;
local int *lptr = &lval;
global int gval = SOME_OTHER_VAL;
global int *gptr = &gval;
ptr = lptr;
```

```
if (ptr == gptr) // legal
{
}
if (ptr == lptr) // legal
{
}
if (lptr == gptr) // illegal, compiler error
{
}
```

Consider the following example:

```
bool callee(int *p1, int *p2)
{
    if (p1 == p2)
        return true;
    return false;
}
void caller()
{
    global int *gptr = Oxdeadbeef;
    private int *pptr = Oxdeadbeef;
    // behavior of callee is undefined
    bool b = callee(gptr, pptr);
}
```

The behavior of callee is undefined as gptr and pptr are in different address spaces. The example above would have the same undefined behavior if the equality operator is replaced with a relational operator.

## Examples:

```
int *ptr = NULL;
local int *lptr = NULL;
global int *gptr = NULL;
if (ptr == NULL) // legal
{
}
```

```
if (ptr == lptr) // legal
{
}
if (lptr == gptr) // compile-time error
{
}
ptr = lptr; // legal
intptr l = (intptr_t)lptr;
if (l == 0) // legal
{
    ...
}
if (l == NULL) // legal
{
    ...
}
```


## Clause 6.5.9 - Equality operators, replace first sentence of paragraph 6 with:

Two pointers compare equal if and only if both are null pointers with overlapping address spaces.

Clause 6.5.15 - Conditional operator, add another constraint paragraph:
If the second and third operands are pointers into different address spaces, the address spaces must overlap.

Examples:

```
kernel void testl()
{
    global int arr[5] = { 0, 1, 2, 3, 4 };
    int *p = &arr[1];
    global int *q = &arr[3];
    local int *r = NULL;
    int *val = NULL;
    // legal. 2nd and 3rd operands are in address spaces
```

```
    // that overlap
    val = (q >= p) ? q : p;
    // compiler error. 2nd and 3rd operands are in disjoint
    // address spaces
    val = (q >= p) ? q : r;
}
```

Clause 6.5.16.1 - Simple assignment, change the third and fourth bullets of paragraph 1 to:

- both operands are pointers to qualified or unqualified versions of compatible types, the referenced address space of the left encloses the referenced address space of the right, and the type pointed to by the left has all the qualifiers of the type pointed to by the right.
- one operand is a pointer to an object or incomplete type and the other is a pointer to a qualified or unqualified version of void, the referenced address space of the left encloses the referenced address space of the right, and the type pointed to by the left has all the qualifiers of the type pointed to by the right.

Examples:

```
kernel void f()
{
    int *ptr;
    local int *lptr;
    global int *gptr;
    local int val = 55;
    ptr = &val; // legal: implicit cast to generic, then assign
    lptr = ptr; // illegal: no implicit cast from
        // generic to local
    lptr = gptr; // illegal: no implicit cast from
        // global to local
    ptr = gptr; // legal: implicit cast from global to generic,
        // then assign
}
```


## Clause 6.7.2.1 - Structure and union specifiers, add a new constraint paragraph:

Within a structure or union specifier, the type of a member shall not be qualified by an address space qualifier.

## Clause 6.7.3 - Type qualifiers, add three new constraint paragraphs:

No type shall be qualified by qualifiers for two or more different address spaces.

### 6.6 Access Qualifiers

Image objects specified as arguments to a kernel can be declared to be read-only, write-only or read-write. A kernel cannot read from and write to the same image object. The __read_only (or read_only) and __write_only (or write_only) qualifiers must be used with image object arguments to declare if the image object is being read or written by a kernel or function. The __read_write (or read_write) qualifier must be used with image object arguments of kernels and of user-defined functions to declare if the image object is being both read and written by the kernel. If no qualifier is provided, read_only is assumed.

In the following example

```
kernel void
foo (read_only image2d_t imageA,
    writ\overline{e_only image2\overline{d_t imageB)}}\mathbf{~}\mathrm{ ( }
{
}
```

imageA is a read-only 2D image object, and imageB is a write-only 2D image object.
The sampler-less read image and write image built-ins can be used with image declared with the read_write (or read_write) qualifier. Calls to built-ins that read from an image using a sampler for images declared with the __read_write (or read_write) qualifier will be a compilation error.

Pipe objects specified as arguments to a kernel also use these access qualifiers. Refer to section 6.13.16 for a detailed description on how these access qualifiers can be used with pipes.

The $\qquad$ read_only, __ write_only, __read_write, read_only, write_only and read_write names are reserved for use as access qualifiers and shall not be used otherwise.

### 6.7 Function Qualifiers

### 6.7.1 __kernel (or kernel)

The kernel (or kernel) qualifier declares a function to be a kernel that can be executed by an application on an OpenCL device(s). The following rules apply to functions that are declared with this qualifier:

It can be executed on the device only

* It can be called by the host
* It is just a regular function call if a __ kernel function is called by another kernel function.


## NOTE:

Kernel functions with variables declared inside the function with the local or local qualifier can be called by the host using appropriate APIs such as clEnqueueNDRangeKernel.

The $\qquad$ kernel and kernel names are reserved for use as functions qualifiers and shall not be used otherwise.

### 6.7.2 Optional Attribute Qualifiers

The $\qquad$ kernel qualifier can be used with the keyword $\qquad$ attribute $\qquad$ to declare additional information about the kernel function as described below.

The optional $\qquad$ attribute $\qquad$ ( (vec_type_hint (<type>)) ) ${ }^{24}$ is a hint to the compiler and is intended to be a representation of the computational width of the $\qquad$ kernel, and should serve as the basis for calculating processor bandwidth utilization when the compiler is looking to autovectorize the code. In the $\qquad$ attribute $\qquad$ ((vec_type_hint(<type>))) qualifier <type> is one of the built-in vector types listed in table 6.2 or the constituent scalar element types. If vec_type_hint (<type>) is not specified, the kernel is assumed to have the __attribute__((vec_type_hint(int))) qualifier.

For example, where the developer specified a width of $f l o a t 4$, the compiler should assume that the computation usually uses up to 4 lanes of a float vector, and would decide to merge work-items or possibly even separate one work-item into many threads to better match the hardware capabilities. A conforming implementation is not required to autovectorize code, but

[^14]shall support the hint. A compiler may autovectorize, even if no hint is provided. If an implementation merges N work-items into one thread, it is responsible for correctly handling cases where the number of global or local work-items in any dimension modulo N is not zero.

Examples:

```
// autovectorize assuming float4 as the
// basic computation width
__kernel __attribute__((vec_type_hint(float4)))
void foo( __global float4 *p ) { .... }
// autovectorize assuming double as the
// basic computation width
__kernel ___attribute__((vec_type_hint(double)))
void foo( __global float4 *\overline{p}) { .... }
// autovectorize assuming int (default)
// as the basic computation width
    kernel
void foo( __global float4 *p ) { .... }
```

If for example, a __kernel function is declared with $\qquad$ attribute $\qquad$ ( vec_type_hint (float 4)) ) (meaning that most operations in the __kernel function are explicitly vectorized using float 4) and the kernel is running using Intel® Advanced Vector Instructions (Intel® AVX) which implements a 8 -float-wide vector unit, the autovectorizer might choose to merge two work-items to one thread, running a second work-item in the high half of the 256-bit AVX register.

As another example, a Power4 machine has two scalar double precision floating-point units with an 6 -cycle deep pipe. An autovectorizer for the Power4 machine might choose to interleave six kernels declared with the___attribute__(( vec_type_hint (double2))) qualifier into one hardware thread, to ensure that there is always 12 -way parallelism available to saturate the FPUs. It might also choose to merge 4 or 8 work-items (or some other number) if it concludes that these are better choices, due to resource utilization concerns or some preference for divisibility by 2 .

The optional __attribute__((work_group_size_hint (X, Y, Z)) ) is a hint to the compiler and is intended to specify the work-group size that may be used i.e. value most likely to be specified by the local_work_size argument to clEnqueueNDRangeKernel. For example the __attribute__((work_group_size_hint(1, 1, 1))) is a hint to the compiler that the kernel will most likely be executed with a work-group size of 1 .

The optional __attribute__((reqd_work_group_size (X, Y, Z))) is the workgroup size that must be used as the local_work_size argument to clEnqueueNDRangeKernel. This allows the compiler to optimize the generated code appropriately for this kernel.

If Z is one, the work_dim argument to clEnqueueNDRangeKernel can be 2 or 3 . If Y and Z are one, the work_dim argument to clEnqueueNDRangeKernel can be 1,2 or 3 .

The optional___attribute__( (nosvm)) qualifier can be used with a pointer variable to inform the compiler that the pointer does not refer to a shared virtual memory region.
Note: $\qquad$ attribute $\qquad$ ( (nosvm)) is depractated and the compiler can ignore it.

### 6.8 Storage-Class Specifiers

The typedef, extern and static storage-class specifiers are supported. The auto and register storage-class specifiers are not supported.

The extern storage-class specifier can only be used for functions (kernel and non-kernel functions) and global variables declared in program scope or variables declared inside a function (kernel and non-kernel functions). The static storage-class specifier can only be used for non-kernel functions, global variables declared in program scope and variables inside a function declared in the global or constant address space.

Examples:

```
extern constant float4 noise_table[256];
static constant float4 color_table[256];
extern kernel void my_foo(image2d_t img);
extern void my_bar(global float *a);
kernel void my_func(image2d_t img, global float *a)
{
        extern constant float4 a;
        static constant float4 b = (float4)(1.0f); // OK.
        static float c; // Error: No implicit address space
        global int hurl; // Error: Must be static
        my foo(img);
        ...
        my_bar(a);
        ...
        while (1)
        {
            static global int inside; // OK.
        }
        ...
}
```


### 6.9 Restrictions ${ }^{25}$

a. The use of pointers is somewhat restricted. The following rules apply:

Arguments to kernel functions declared in a program that are pointers must be declared with the $\qquad$ global, $\qquad$ constant or $\qquad$ local qualifier.

A pointer declared with the __constant qualifier can only be assigned to a pointer declared with the $\qquad$ constant qualifier respectively.

Pointers to functions are not allowed.

* Arguments to kernel functions in a program cannot be declared as a pointer to a peinter(s). Variables inside a function or arguments to non kernel functions in a program can be declared as a pointer to a pointer(s).
b. An image type (image2d_t, image3d_t, image2d_array_t, image1d_t, image1d_buffer_t or image1d_array_t) can only be used as the type of a function argument. An image function argument cannot be modified. Elements of an image can only be accessed using built-in functions described in section 6.13.14.

An image type cannot be used to declare a variable, a structure or union field, an array of images, a pointer to an image, or the return type of a function. An image type cannot be used with the __global, __private, __local and __constant address space qualifiers. The image 3d_t type cannet be used with the__write_only access qualifier unless the cl_khr_3d_image_writes extension is enabled. An image type cannot be used with the __read_write access qualifer which is reserved for fature use.

The sampler type (sampler_t) can only be used as the type of a function argument or a variable declared in the program scope or the outermost scope of a kernel function. The behavior of a sampler variable declared in a non-outermost scope of a kernel function is implementation-defined. A sampler argument or variable cannot be modified.

The sampler type cannot be used to declare a structure or union field, an array of samplers, a pointer to a sampler, or the return type of a function. The sampler type cannot be used with the __ local and __global address space qualifiers.
c. Bit-field struct members are currently not supported.
d. Variable length arrays and structures with flexible (or unsized) arrays are not supported.

[^15]e. Variadic macros and functions with the exception of printf and enqueue_kernel are not supported.
f. Unless defined in the OpenCL specification, the library functions, macros, types, and constants defined in the C99 standard headers assert.h, ctype.h, complex.h, errno.h, fenv.h, float.h, inttypes.h, limits.h, locale.h, setjmp.h, signal.h, stdarg.h, stdio.h, stdlib.h, string.h, tgmath.h, time.h, wchar.h and wctype.h are not available and cannot be included by a program.
g. The auto and register storage-class specifiers are not supported.
h. Predefined identifiers are not supported.
i. Recursion is not supported.
j. The return type of a kernel function must be void.
k. Arguments to kernel functions in a program cannot be declared with the built-in scalar types bool, size_t, ptrdiff_t, intptr_t, and uintptr_t or a struct and/or union that contain fields declared to be one of these built-in scalar types. The size in bytes of these types are implementation-defined and in addition can also be different for the OpenCL device and the host processor making it difficult to allocate buffer objects to be passed as arguments to a kernel declared as pointer to these types.

1. half is not supported as half can be used as a storage format ${ }^{26}$ only and is not a data type on which floating-point arithmetic can be performed.
m . Whether or not irreducible control flow is illegal is implementation defined.
n. Built-in types that are less than 32 bits in size i.e. char, uchar, char2, uchar2, short, ushort, and half have the following restriction:
\# Writes to a pointer (or arrays) of type char, uchar, char2, uchar2, short, ushort, and half or to elements of a struct that are of type char, uchar, char2, uchar2, short and ushort are not supported. Refer to section 9.9 for additional information.

The kernel example below shows what memory operations are not supported on built in types less than 32 bits in size.


[^16]
o. The type qualifiers const, restrict and volatile as defined by the C99 specification are supported. These qualifiers cannot be used with image2d_t, image3d_t, image2d_array_t, image2d_depth_t, image2d_array_depth_t, image1d_t, image1d_buffer_t and image1d_array_t types. Types other than pointer types shall not use the restrict qualifier.
p. The event type (event_t) cannot be used as the type of a kernel function argument. The event type cannot be used to declare a program scope variable. The event type cannot be used to declare a structure or union field. The event type cannot be used with the $\qquad$ local, $\qquad$ constant and $\qquad$ global address space qualifiers.
q. The clk_event_t, ndrange_t and reserve_id_t types cannot be used as arguments to kernel functions that get enqueued from the host. The clk_event_t and reserve_id_t types cannot be declared in program scope.
r. The values returned by applying the sizeof operator to the queue_t, clk_event_t, ndrange_t and reserve_id_t types is implementation-defined.
s. Kernels enqueued by the host must continue to have their arguments that are a pointer to a type declared to point to a named address space.
t. A function in an OpenCL program cannot be called main.
u. Implicit function declaration is not supported.

### 6.10 Preprocessor Directives and Macros

The preprocessing directives defined by the C99 specification are supported.
The \# pragma directive is described as:
\# pragma pp-tokens opt new-line
A \# pragma directive where the preprocessing token OPENCL (used instead of STDC) does not immediately follow pragma in the directive (prior to any macro replacement) causes the implementation to behave in an implementation-defined manner. The behavior might cause translation to fail or cause the translator or the resulting program to behave in a non-conforming manner. Any such pragma that is not recognized by the implementation is ignored. If the preprocessing token OPENCL does immediately follow pragma in the directive (prior to any macro replacement), then no macro replacement is performed on the directive, and the directive shall have one of the following forms whose meanings are described elsewhere:

```
#pragma OPENCL FP_CONTRACT on-off-switch
    on-off-switch: one of ON OFF DEFAULT
#pragma OPENCL EXTENSION extensionname : behavior
#pragma OPENCL EXTENSION all : behavior
```

The following predefined macro names are available.
__FILE__The presumed name of the current source file (a character string literal).
__LINE __ The presumed line number (within the current source file) of the current source line (an integer constant).
__OPENCL_VERSION__ substitutes an integer reflecting the version number of the OpenCL supported by the OpenCL device. The version of OpenCL described in this document will have
$\qquad$ OPENCL_VERSION $\qquad$ substitute the integer 200.

CL_VERSION_1_0 substitutes the integer 100 reflecting the OpenCL 1.0 version.
CL_VERSION_1_1 substitutes the integer 110 reflecting the OpenCL 1.1 version.
CL_VERSION_1_2 substitutes the integer 120 reflecting the OpenCL 1.2 version.

CL_VERSION_2_0 substitutes the integer 200 reflecting the OpenCL 2.0 version.
__OPENCL_C_VERSION__ substitutes an integer reflecting the OpenCL C version
specified by the -cl-std build option (specified in section 5.8.4.5) to clBuildProgram or clCompileProgram. If the -cl-std build option is not specified, the highest OpenCL C 1.x language version supported by each device is used as the version of OpenCL C when compiling the program for each device. The version of OpenCL C described in this document will have __OPENCL_C_VERSION__ substitute the integer 200 if -cl-std=CL2.0 is specified.
__ENDIAN_LITTLE $\qquad$ is used to determine if the OpenCL device is a little endian architecture $\overline{\text { or a }}$ a big endian architecture (an integer constant of 1 if device is little endian and is undefined otherwise). Also refer to CL_DEVICE_ENDIAN_LITTLE specified in table 4.3.

```
__kernel_exec(X, typen) (and kernel_exec(X, typen)) is defined as
    __kernel __attribute__((work_group_size_hint(X, 1, 1))) \
        __attribute__((vec_type_hint(typen)))
```

        IMAGE_SUPPORT__ is used to determine if the OpenCL device supports images. This is an integer constant of 1 if images are supported and is undefined otherwise. Also refer to CL_DEVICE_IMAGE_SUPPORT specified in table 4.3.
    __FAST_RELAXED_MATH__ is used to determine if the -cl-fast-relaxed-math optimization $\overline{\text { option }}$ is specified in build options given to clBuildProgram or clCompileProgram. This is an integer constant of 1 if the -cl-fast-relaxed-math build option is specified and is undefined otherwise.

The NULL macro expands to a null pointer constant. An integer constant expression with the value 0 , or such an expression cast to type void * is called a null pointer constant.

The macro names defined by the C 99 specification but not currently supported by OpenCL are reserved for future use.

The predefined identifier $\qquad$ func $\qquad$ is available.

### 6.11 Attribute Qualifiers

This section describes the syntax with which $\qquad$ attribute $\qquad$ may be used, and the constructs to which attribute specifiers bind.

An attribute specifier is of the form $\qquad$ attribute $\qquad$ ((attribute-list)).

An attribute list is defined as:

```
attribute-list:
    attributeopt
    attribute-list , attributeopt
attribute:
    attribute-token attribute-argument-clauseopt
attribute-token:
    identifier
attribute-argument-clause:
    ( attribute-argument-list )
attribute-argument-list:
    attribute-argument
    attribute-argument-list, attribute-argument
attribute-argument:
    assignment-expression
```

This syntax is taken directly from GCC but unlike GCC, which allows attributes to be applied only to functions, types, and variables, OpenCL attributes can be associated with:

4 types;

* functions;
* variables;

4 blocks; and
\# control-flow statements.

In general, the rules for how an attribute binds, for a given context, are non-trivial and the reader is pointed to GCC's documentation and Maurer and Wong's paper [See 16. and 17. in section 11 - References] for the details.

### 6.11.1 Specifying Attributes of Types

The keyword ___attribute__ allows you to specify special attributes of enum, struct and union types when you define such types. This keyword is followed by an attribute specification inside double parentheses. Two attributes are currently defined for types: aligned, and packed.

You may specify type attributes in an enum, struct or union type declaration or definition, or for other types in a typedef declaration.

For an enum, struct or union type, you may specify attributes either between the enum, struct or union tag and the name of the type, or just past the closing curly brace of the definition. The former syntax is preferred.
aligned (alignment)
This attribute specifies a minimum alignment (in bytes) for variables of the specified type. For example, the declarations:

```
struct S { short f[3]; } __attribute__ ((aligned (8)));
typedef int more_aligned_int __attribute__ ((aligned (8)));
```

force the compiler to insure (as far as it can) that each variable whose type is struct $S$ or more_aligned_int will be allocated and aligned at least on a 8 -byte boundary.

Note that the alignment of any given struct or union type is required by the ISO C standard to be at least a perfect multiple of the lowest common multiple of the alignments of all of the members of the struct or union in question and must also be a power of two. This means that you can effectively adjust the alignment of a struct or union type by attaching an aligned attribute to any one of the members of such a type, but the notation illustrated in the example above is a more obvious, intuitive, and readable way to request the compiler to adjust the alignment of an entire struct or union type.

As in the preceding example, you can explicitly specify the alignment (in bytes) that you wish the compiler to use for a given struct or union type. Alternatively, you can leave out the alignment factor and just ask the compiler to align a type to the maximum useful alignment for the target machine you are compiling for. For example, you could write:
struct $S$ \{ short f[3]; \} ___attribute__ ((aligned));
Whenever you leave out the alignment factor in an aligned attribute specification, the compiler automatically sets the alignment for the type to the largest alignment which is ever used for any data type on the target machine you are compiling for. In the example
above, the size of each short is 2 bytes, and therefore the size of the entire struct $S$ type is 6 bytes. The smallest power of two which is greater than or equal to that is 8 , so the compiler sets the alignment for the entire struct $S$ type to 8 bytes.

Note that the effectiveness of aligned attributes may be limited by inherent limitations of the OpenCL device and compiler. For some devices, the OpenCL compiler may only be able to arrange for variables to be aligned up to a certain maximum alignment. If the OpenCL compiler is only able to align variables up to a maximum of 8 byte alignment, then specifying aligned (16) in an __attribute__ will still only provide you with 8 byte alignment. See your platform-specific documentation for further information.

The aligned attribute can only increase the alignment; but you can decrease it by specifying packed as well. See below.

```
packed
```

This attribute, attached to struct or union type definition, specifies that each member of the structure or union is placed to minimize the memory required. When attached to an enum definition, it indicates that the smallest integral type should be used.

Specifying this attribute for struct and union types is equivalent to specifying the packed attribute on each of the structure or union members.

In the following example struct my_packed_struct's members are packed closely together, but the internal layout of its s member is not packed. To do that, struct my_unpacked_struct would need to be packed, too.

```
struct my_unpacked_struct
{
    char c;
    int i;
};
struct __attribute__ ((packed)) my_packed_struct
{
    char c;
    int i;
    struct my_unpacked_struct s;
};
```

You may only specify this attribute on the definition of a enum, struct or union, not on a typedef which does not also define the enumerated type, structure or union.

### 6.11.2 Specifying Attributes of Functions

Refer to section 6.7 for the function attribute qualifiers currently supported.

### 6.11.3 Specifying Attributes of Variables

The keyword
$\qquad$ attribute $\qquad$ allows you to specify special attributes of variables or structure fields. This keyword is followed by an attribute specification inside double parentheses. The following attribute qualifiers are currently defined:

```
aligned (alignment)
```

This attribute specifies a minimum alignment for the variable or structure field, measured in bytes. For example, the declaration:

```
int x ___attribute__ ((aligned (16))) = 0;
```

causes the compiler to allocate the global variable x on a 16-byte boundary. The alignment value specified must be a power of two.

You can also specify the alignment of structure fields. For example, to create doubleword aligned int pair, you could write:

```
struct foo { int x[2] ___attribute__ ((aligned (8))); };
```

This is an alternative to creating a union with a double member that forces the union to be double-word aligned.

As in the preceding examples, you can explicitly specify the alignment (in bytes) that you wish the compiler to use for a given variable or structure field. Alternatively, you can leave out the alignment factor and just ask the compiler to align a variable or field to the maximum useful alignment for the target machine you are compiling for. For example, you could write:

```
short array[3] __attribute__ ((aligned));
```

Whenever you leave out the alignment factor in an aligned attribute specification, the OpenCL compiler automatically sets the alignment for the declared variable or field to the largest alignment which is ever used for any data type on the target device you are compiling for.

When used on a struct, or struct member, the aligned attribute can only increase
the alignment; in order to decrease it, the packed attribute must be specified as well. When used as part of a typedef, the aligned attribute can both increase and decrease alignment, and specifying the packed attribute will generate a warning.

Note that the effectiveness of aligned attributes may be limited by inherent limitations of the OpenCL device and compiler. For some devices, the OpenCL compiler may only be able to arrange for variables to be aligned up to a certain maximum alignment. If the OpenCL compiler is only able to align variables up to a maximum of 8 byte alignment, then specifying aligned (16) in an $\qquad$ attribute $\qquad$ will still only provide you with 8 byte alignment. See your platform-specific documentation for further information.
packed
The packed attribute specifies that a variable or structure field should have the smallest possible alignment-one byte for a variable, unless you specify a larger value with the aligned attribute.

Here is a structure in which the field $x$ is packed, so that it immediately follows $a$ :

```
struct foo
{
    char a;
    int x[2] __attribute__ ((packed));
};
```

An attribute list placed at the beginning of a user-defined type applies to the variable of that type and not the type, while attributes following the type body apply to the type. For example:

```
    /* a has alignment of 128 */
    __attribute__((aligned(128))) struct A {int i;} a;
    /* b has alignment of 16 */
    __attribute__((aligned(16))) struct B {double d;}
        __attribute__((aligned(32))) b ;
struct A a1; /* a1 has alignment of 4 */
struct B b1; /* b1 has alignment of 32 */
endian (endiantype)
```

The endian attribute determines the byte ordering of a variable. endiantype can be set to host indicating the variable uses the endianness of the host processor or can be set to device indicating the variable uses the endianness of the device on which the kernel will be executed. The default is device.

For example:

```
global float4 *p ___attribute__ ((endian(host)));
```

specifies that data stored in memory pointed to by p will be in the host endian format.
The endian attribute can only be applied to pointer types that are in the global or constant address space. The endian attribute cannot be used for variables that are not a pointer type. The endian attribute value for both pointers must be the same when one pointer is assigned to another.

### 6.11.4 Specifying Attributes of Blocks and Control-FlowStatements

For basic blocks and control-flow-statements the attribute is placed before the structure in question, for example:

```
__attribute__((attr1)) {...}
for __attribute__((attr2)) (...) ___attribute__((attr3)) {...}
```

Here attr 1 applies to the block in braces and attr2 and attr3 apply to the loop's control construct and body, respectively.

No attribute qualifiers for blocks and control-flow-statements are currently defined.

### 6.11.5 Specifying Attribute For Unrolling Loops

The $\qquad$ attribute $\qquad$ ( (opencl_unroll_hint)) and __attribute__( (opencl_unroll_hint(n))) attribute qualifiers can be used to specify that a loop (for, while and do loops) can be unrolled. This attribute qualifier can be used to specify full unrolling or partial unrolling by a specified amount. This is a compiler hint and the compiler may ignore this directive.
n is the loop unrolling factor and must be a positive integral compile time constant expression. An unroll factor of 1 disables unrolling. If n is not specified, the compiler determines the unrolling factor for the loop.

NOTE: The $\qquad$ attribute $\qquad$ ((opencl_unroll_hint(n))) attribute qualifier must appear immediately before the loop to be affected.

Examples:

```
__attribute__((opencl_unroll_hint(2)))
while (*s != 0)
    * p++ = *s++;
```

The tells the compiler to unroll the above while loop by a factor of 2 .

```
__attribute__((opencl_unroll_hint))
for (int i=\\overline{0; i<2; i+')}
{
}
```

In the example above, the compiler will determine how much to unroll the loop.

```
__attribute__((opencl_unroll_hint(1)))
for (int i=0; i<32; i++)
{
}
```

The above is an example where the loop should not be unrolled.
Below are some examples of invalid usage of

```
__attribute__((opencl_unroll_hint(n))).
__attribute__((opencl_unroll_hint(-1)))
while (...)
{
    ...
}
```

The above example is an invalid usage of the loop unroll factor as the loop unroll factor is negative.


The above example is invalid because the unroll attribute qualifier is used on a non-loop construct

```
kernel void
my_kernel( ... )
```

```
        int x;
    attribute__((opencl_unroll_hint(x))
    for (int i=0; i<x; i++)
    {
    }
}
```

The above example is invalid because the loop unroll factor is not a compile-time constant expression.

### 6.11.6 Extending Attribute Qualifiers

The attribute syntax can be extended for standard language extensions and vendor specific extensions. Any extensions should follow the naming conventions outlined in the introduction to section 9 in the OpenCL 1.2 Extension Specification.

Attributes are intended as useful hints to the compiler. It is our intention that a particular implementation of OpenCL be free to ignore all attributes and the resulting executable binary will produce the same result. This does not preclude an implementation from making use of the additional information provided by attributes and performing optimizations or other transformations as it sees fit. In this case it is the programmer's responsibility to guarantee that the information provided is in some sense correct.

### 6.12 Blocks

OpenCL C 2.0 adds support for the clang block syntax ${ }^{27}$. Like function types, the Block type is a pair consisting of a result value type and a list of parameter types very similar to a function type. Blocks are intended to be used much like functions with the key distinction being that in addition to executable code they also contain various variable bindings to automatic (stack) or global memory.

### 6.12.1 Declaring and Using a Block

You use the ${ }^{\wedge}$ operator to declare a Block variable and to indicate the beginning of a Block literal. The body of the Block itself is contained within \{\}, as shown in this example (as usual with C , ; indicates the end of the statement):

The example is explained in the following illustration:


Notice that the Block is able to make use of variables from the same scope in which it was defined.

If you declare a Block as a variable, you can then use it just as you would a function:

```
int multiplier = 7;
```

[^17]```
int (^myBlock)(int) = ^(int num) {
    return num * multiplier;
};
printf("%d\n", myBlock(3));
// prints 21
```


### 6.12.2 Declaring a Block Reference

Block variables hold references to Blocks. You declare them using syntax similar to that you use to declare a pointer to a function, except that you use ${ }^{\wedge}$ instead of *. The Block type fully interoperates with the rest of the C type system. The following are valid Block variable declarations:

```
void (^blockReturningVoidWithVoidArgument) (void);
int (^blockReturningIntWithIntAndCharArguments)(int,
char);
```

A Block that takes no arguments must specify void in the argument list. A Block reference may not be dereferenced via the pointer dereference operation *, and thus a Block's size may not be computed at compile time.

Blocks are designed to be fully type safe by giving the compiler a full set of metadata to use to validate use of Blocks, parameters passed to blocks, and assignment of the return value.

You can also create types for Blocks-doing so is generally considered to be best practice when you use a block with a given signature in multiple places:

```
typedef float (^MyBlockType)(float, float);
MyBlockType myFirstBlock = // ...;
MyBlockType mySecondBlock = // ...;
```


### 6.12.3 Block Literal Expressions

A Block literal expression produces a reference to a Block. It is introduced by the use of the ${ }^{\wedge}$ token as a unary operator.

```
Block_literal_expression ::= ^ block_decl compound_statement_body
block_decl ::=
block_decl ::= parameter_list
block_decl ::= type_expression
```

where type expression is extended to allow ${ }^{\wedge}$ as a Block reference where $*$ is allowed as a function reference.

The following Block literal:

```
^ void (void) { printf("hello world\n"); }
```

produces a reference to a Block with no arguments with no return value.
The return type is optional and is inferred from the return statements. If the return statements return a value, they all must return a value of the same type. If there is no value returned the inferred type of the Block is void; otherwise it is the type of the return statement value. If the return type is omitted and the argument list is (void ), the ( void ) argument list may also be omitted.

So:

```
^ ( void ) { printf("hello world\n"); }
```

and:

```
^ { printf("hello world\n"); }
```

are exactly equivalent constructs for the same expression.
The compound statement body establishes a new lexical scope within that of its parent. Variables used within the scope of the compound statement are bound to the Block in the normal manner with the exception of those in automatic (stack) storage. Thus one may access functions and global variables as one would expect, as well as static local variables.

Local automatic (stack) variables referenced within the compound statement of a Block are imported and captured by the Block as const copies. The capture (binding) is performed at the time of the Block literal expression evaluation.

The compiler is not required to capture a variable if it can prove that no references to the variable will actually be evaluated.

The lifetime of variables declared in a Block is that of a function..
Block literal expressions may occur within Block literal expressions (nested) and all variables captured by any nested blocks are implicitly also captured in the scopes of their enclosing Blocks.

A Block literal expression may be used as the initialization value for Block variables at global or local static scope.

You can also declare a Block as a global literal in program scope.

```
int GlobalInt = 0;
int (^getGlobalInt)(void) = ^{ return GlobalInt; };
```


### 6.12.4 Control Flow

The compound statement of a Block is treated much like a function body with respect to control flow in that continue, break and goto do not escape the Block.

### 6.12.5 Restrictions

The following Blocks features are currently not supported in OpenCL C.
4 The __block storage type.

* The Block_copy() and Block_release() functions that copy and release Blocks.
* Blocks with variadic arguments.

4 Arrays of Blocks.

* Blocks as structures and union members.

Block literals are assumed to allocate memory at the point of definition and to be destroyed at the end of the same scope. To support these behaviors, additional restrictions ${ }^{28}$ in addition to the above feature restrictions are:

* Block variables must be defined and used in a way that allows them to be statically determinable at build or "link to executable" time. In particular:
- Block variables assigned in one scope must be used only with the same or any nested scope.
- The "extern" storage-class specified cannot be used with program scope block variables.
- Block variable declarations are implicitly qualified with const. Therefore all block variables must be initialized at declaration time and may not be reassigned.
- A block cannot be the return value of a function.

The unary operators (* and \&) cannot be used with a Block.
Pointers to Blocks are not allowed.
Blocks cannot be used as expressions of the ternary selection operator (?:).

[^18]A Block cannot reference or capture another Block variable declared in the outer scope (example 4).

Some examples that describe legal and illegal issue of Blocks in OpenCL C are described below.

## Example 1:

```
void foo(int *x, int (^bar)(int, int))
{
    *x = bar(*x, *x);
}
kernel
void k(global int *x, global int *z)
{
    if (some expression)
        *x = foo(x,
            ^int(int x, int y){return x+y+*z;}); // legal
    else
        *x = foo(x,
            ^int(int x, int y){return (x*y)-*z;}); // legal
}
```


## Example 2:

```
kernel
void k(global int *x, global int *z)
{
    int ^(tmp)(int, int);
    if (some expression)
    {
        tmp = ^int(int x, int y){return x+y+*z;}); // illegal
    }
    *x = foo(x, tmp);
}
```


## Example 3:

```
int GlobalInt = 0;
int (^getGlobalInt)(void) = ^{ return GlobalInt; }; // legal
int (^getAnotherGlobalInt)(void); // illegal
extern int (^getExternGlobalInt)(void); // illegal
void foo()
{
    getGlobalInt = ^{ return 0; }; // illegal - cannot assign to
```


## Example 4:

```
int ^(block1)(int, int) = ^int(int x, int y){return x+y+*z;});
int foo(int x, int y) { return blockl(x, y); }
kernel
void k(global int *x, global int *z)
{
    int ^(block2)(int, int) =
            ^int(int x, int y){return foo(x, y);}); // illegal
    int ^(block3)(int, int) =
            ^int(int x, int y){return blockl(x, y);}); // illegal
    int ^(block4)(int, int) =
            ^int(int x, int y){return block2(x, y);}); // illegal
}
```


### 6.13 Built-in Functions

The OpenCL C programming language provides a rich set of built-in functions for scalar and vector operations. Many of these functions are similar to the function names provided in common C libraries but they support scalar and vector argument types. Applications should use the built-in functions wherever possible instead of writing their own version.

User defined OpenCL C functions, behave per C standard rules for functions (C99, TC2, Section 6.9.1). On entry to the function, the size of each variably modified parameter is evaluated and the value of each argument expression is converted to the type of the corresponding parameter as per usual arithmetic conversion rules described in section 6.2.6. Built-in functions described in this section behave similarly, except that in order to avoid ambiguity between multiple forms of the same built-in function, implicit scalar widening shall not occur. Note that some built-in functions described in this section do have forms that operate on mixed scalar and vector types, however.

### 6.13.1 Work-Item Functions

Table 6.7 describes the list of built-in work-item functions that can be used to query the number of dimensions, the global and local work size specified to cIEnqueueNDRangeKernel, and the global and local identifier of each work-item when this kernel is being executed on a device.

| Function | Description |
| :--- | :--- |
| uint get_work_dim () | Returns the number of dimensions in use. This is the <br> value given to the work_dim argument specified in <br> clEnqueueNDRangeKernel. |
| size_t get_global_size (uint dimindx) | Returns the number of global work-items specified for <br> dimension identified by dimindx. This value is given by <br> the global_work_size argument to <br> clEnqueueNDRangeKernel. Valid values of dimindx <br> are 0 to get_work_dim() - 1. For other values of <br> dimindx, get_global_size() returns 1. |
| size_t get_global_id (uint dimindx) | Returns the unique global work-item ID value for <br> dimension identified by dimindx. The global work-item <br> ID specifies the work-item ID based on the number of <br> global work-items specified to execute the kernel. Valid <br> values of dimindx are 0 to get_work_dim() - 1. For <br> other values of dimindx, get_global_id() returns 0. |
| size_t get_local_size (uint dimindx) | Returns the number of local work-items specified in <br> dimension identified by dimindx. This value is at most <br> the value given by the local_work_size argument to <br> clEnqueueNDRangeKernel if local_work_size is not <br> NULL; otherwise the OpenCL implementation chooses |


|  | an appropriate local_work_size value which is returned by this function. If the kernel is executed with a nonuniform work-group size ${ }^{29}$, calls to this built-in from some work-groups may return different values than calls to this built-in from other work-groups. <br> Valid values of dimindx are 0 to get_work_dim()-1. <br> For other values of dimindx, get_local_size() returns 1. |
| :---: | :---: |
| size_t get_enqueued_local_size ( uint dimindx) | Returns the same value as that returned by get_local_size $(\operatorname{dimindx} x$ ) if the kernel is executed with a uniform work-group size. <br> If the kernel is executed with a non-uniform work-group size, returns the number of local work-items in each of the work-groups that make up the uniform region of the global range in the dimension identified by dimindx. If the local_work_size argument to <br> cIEnqueueNDRangeKernel is not NULL, this value will match the value specified in local_work_size $[$ dimindx]. If local_work_size is NULL, this value will match the local size that the implementation determined would be most efficient at implementing the uniform region of the global range. <br> Valid values of dimind $x$ are 0 to get_work_dim( $)-1$. For other values of dimindx, get_enqueued_local_size() returns 1. |
| size_t get_local_id (uint dimindx) | Returns the unique local work-item ID i.e. a work-item within a specific work-group for dimension identified by dimindx. Valid values of dimindx are 0 to get_work_dim ()$-1$. For other values of dimindx, get_local_id() returns 0 . |
| size_t get_num_groups (uint dimindx) | Returns the number of work-groups that will execute a kernel for dimension identified by dimindx. Valid values of dimind $x$ are 0 to get_work_dim() - 1 . For other values of dimindx, get_num_groups () returns 1. |
| size_t get_group_id (uint dimindx) | get_group_id returns the work-group ID which is a number from 0 .. get_num_groups $(\operatorname{dimind} d x)-1$. Valid values of dimindx are 0 to get_work_dim()-1. For other values, get_group_id() returns 0 . |
| size_t get_global_offset (uint dimindx) | get_global_offset returns the offset values specified in global_work_offset argument to <br> clEnqueueNDRangeKernel. Valid values of dimindx are 0 to get_work_dim ()$-1$. For other values, |

[^19]|  | get_global_offset() returns 0. |
| :---: | :---: |
| size_t get_global_linear_id () | Returns the work-items 1-dimensional global ID. For 1D work-groups, it is computed as get_global_id(0) - get_global_offset(0). <br> For 2D work-groups, it is computed as (get_global_id(1) - get_global_offset(1)) * get_global_size(0) + (get_global_id(0) get_global_offset(0)). <br> For 3D work-groups, it is computed as ((get_global_id(2) - get_global_offset(2)) * get_global_size(1) * get_global_size(0)) + ((get_global_id(1) - get_global_offset(1)) * get_global_size (0)) + (get_global_id(0) - get_global_offset(0)). |
| size_t get_local_linear_id () | Returns the work-items 1-dimensional local ID. <br> For 1D work-groups, it is the same value as get_local_id(0). <br> For 2D work-groups, it is computed as get_local_id(1) * get_local_size (0) + get_local_id(0). <br> For 3D work-groups, it is computed as (get_local_id(2) * get_local_size(1) * get_local_size(0)) + (get_local_id(1) * get_local_size(0)) + get_local_id(0). |

Table 6.7 Work-Item Functions Table

### 6.13.2 Math Functions

The list of built-in math functions is described in table 6.8. The built-in math functions are categorized into the following:

4 A list of built-in functions that have scalar or vector argument versions, and,
4 A list of built-in functions that only take scalar float arguments.
The vector versions of the math functions operate component-wise. The description is percomponent.

The built-in math functions are not affected by the prevailing rounding mode in the calling environment, and always return the same value as they would if called with the round to nearest even rounding mode.

Table 6.8 describes the list of built-in math functions that can take scalar or vector arguments. We use the generic type name gentype to indicate that the function can take float, float2, float3, float4, float8, float16, double, double2, double3, double4, double8 or double16 as the type for the arguments. We use the generic type name gentypef to indicate that the function can take float, float2, float3, float4, float8, or float16 as the type for the arguments. We use the generic type name gentyped to indicate that the function can take double, double2, double3, double4, double8 or double16 as the type for the arguments. For any specific use of a function, the actual type has to be the same for all arguments and the return type, unless otherwise specified.

| Function | Description |
| :--- | :--- |
| gentype acos (gentype) | Arc cosine function. Returns an angle in radians. |
| gentype acosh (gentype) | Inverse hyperbolic cosine. Returns an angle in <br> radians. |
| gentype acospi (gentype $x$ ) | Compute acos $(x) / \pi$. |
| gentype asin (gentype) | Arc sine function. Returns an angle in radians. |
| gentype asinh (gentype) | Inverse hyperbolic sine. Returns an angle in <br> radians. |
| gentype asinpi (gentype $x$ ) | Compute asin $(x) / \pi$. |
| gentype atan (gentype $y \_$over_ $x$ ) | Arc tangent function. Returns an angle in radians. |
| gentype atan2 (gentype $y$, gentype $x$ ) | Arc tangent of $y / x$. Returns an angle in radians. |
| gentype atanh (gentype) | Hyperbolic arc tangent. Returns an angle in radians. |
| gentype atanpi (gentype $x$ ) | Compute atan $(x) / \pi$. |
| gentype atan2pi (gentype $y$, gentype $x$ ) | Compute atan2 $(y, x) / \pi$. |
| gentype cbrt (gentype) | Compute cube-root. |
| gentype ceil (gentype) | Round to integral value using the round to positive <br> infinity rounding mode. |


| gentype copysign (gentype $x$, gentype $y$ ) | Returns $x$ with its sign changed to match the sign of $y$. |
| :---: | :---: |
| gentype cos (gentype $x$ ) | Compute cosine, where $x$ is an angle in radians. |
| gentype cosh (gentype $x$ ) | Compute hyperbolic consine, where $x$ is an angle in radians. |
| gentype cospi (gentype $x$ ) | Compute $\cos (\pi x)$. |
| gentype erfc (gentype) | Complementary error function. |
| gentype erf (gentype) | Error function encountered in integrating the normal distribution. |
| gentype exp (gentype $x$ ) | Compute the base- e exponential of $x$. |
| gentype exp2 (gentype) | Exponential base 2 function. |
| gentype exp10 (gentype) | Exponential base 10 function. |
| gentype expm1 (gentype $x$ ) | Compute $e^{x}$ - 1.0. |
| gentype fabs (gentype) | Compute absolute value of a floating-point number. |
| gentype fdim (gentype $x$, gentype $y$ ) | $x-y$ if $x>y,+0$ if x is less than or equal to y . |
| gentype floor (gentype) | Round to integral value using the round to negative infinity rounding mode. |
|  | Returns the correctly rounded floating-point representation of the sum of $c$ with the infinitely precise product of $a$ and $b$. Rounding of intermediate products shall not occur. Edge case behavior is per the IEEE 754-2008 standard. |
| gentype fmax (gentype $x$, gentype $y$ ) <br> gentypef fmax (gentypef $x$, float $y$ ) <br> gentyped fmax (gentyped $x$, double $y$ ) | Returns $y$ if $x<y$, otherwise it returns $x$. If one argument is a $\mathrm{NaN}, \mathbf{f m a x}()$ returns the other argument. If both arguments are NaNs, fmax() returns a NaN . |
| gentype fmin ${ }^{30}$ (gentype $x$, gentype $y$ ) gentypef $\mathbf{f m i n}$ (gentypef $x$, float $y$ ) gentyped fmin (gentyped $x$, double $y$ ) | Returns $y$ if $y<x$, otherwise it returns $x$. If one argument is a $\mathrm{NaN}, \mathbf{f m i n}()$ returns the other argument. If both arguments are $\mathrm{NaNs}, \mathbf{f m i n}()$ returns a NaN . |
| gentype fmod (gentype $x$, gentype $y$ ) | Modulus. Returns $x-y$ * trunc ( $x / y$ ). |
| gentype fract (gentype $x$, gentype ${ }^{*}$ iptr) ${ }^{31}$ | Returns fmin( $x$ - floor ( $x$ ), 0x1.fffffep-1f ). floor( x ) is returned in iptr. |
| float $n$ frexp (float $n x$, int $n$ *exp) float frexp (float $x$, int *exp) | Extract mantissa and exponent from $x$. For each component the mantissa returned is a float with magnitude in the interval $[1 / 2,1$ ) or 0 . Each component of $x$ equals mantissa returned $* 2^{e x p}$. |
| double $n$ frexp (double $n x, \operatorname{int} n * \exp )$ double frexp (double $x$, int *exp) | Extract mantissa and exponent from $x$. For each component the mantissa returned is a float with |

[^20]|  | magnitude in the interval $[1 / 2,1$ ) or 0 . Each component of $x$ equals mantissa returned $* 2^{\text {exp }}$. |
| :---: | :---: |
| gentype hypot (gentype $x$, gentype $y$ ) | Compute the value of the square root of $x^{2}+y^{2}$ without undue overflow or underflow. |
| int $n$ ilogb (float $n x$ ) int ilogb (float $x$ ) <br> int $n$ ilogb (double $n x$ ) int ilogb (double $x$ ) | Return the exponent as an integer value. |
| float $n \operatorname{ldexp}$ (float $n x$, int $n k$ ) float $n$ ldexp (float $n x$, int $k$ ) float ldexp (float $x$, int $k$ ) <br> doublen ldexp (doublen $x$, int $n k$ ) doublen Idexp (doublen $x$, int $k$ ) double Idexp (double $x$, int $k$ ) | Multiply $x$ by 2 to the power $k$. |
| ```gentype lgamma (gentype }x\mathrm{ ) floatn lgamma_r (floatn x, intn *signp) float lgamma_r (float x, int *signp) doublen Igamma_r (doublen x, intn *signp) double lgamma_r (double }x\mathrm{ , int *signp)``` | Log gamma function. Returns the natural logarithm of the absolute value of the gamma function. The sign of the gamma function is returned in the signp argument of Igamma_r. |
| gentype $\log$ (gentype) | Compute natural logarithm. |
| gentype log2 (gentype) | Compute a base 2 logarithm. |
| gentype $\log 10$ (gentype) | Compute a base 10 logarithm. |
| gentype log1p (gentype $x$ ) | Compute $\log _{\mathrm{e}}(1.0+x)$. |
| gentype logb (gentype $x$ ) | Compute the exponent of $x$, which is the integral part of $\log _{r}\|x\|$. |
| gentype mad (gentype $a$, gentype $b$, gentype $c$ ) | mad approximates $a * b+c$. Whether or how the product of $a * b$ is rounded and how supernormal or subnormal intermediate products are handled is not defined. mad is intended to be used where speed is preferred over accuracy ${ }^{32}$. |
| gentype maxmag (gentype $x$, gentype $y$ ) | Returns $x$ if $\|x\|>\|y\|, y$ if $\|y\|>\|x\|$, otherwise $\operatorname{fmax}(x, y)$. |
| gentype minmag (gentype $x$, gentype $y$ ) | Returns $x$ if $\|x\|<\|y\|, y$ if $\|y\|<\|x\|$, otherwise $\operatorname{fmin}(x, y)$. |
|  | Decompose a floating-point number. The modf function breaks the argument $x$ into integral and fractional parts, each of which has the same sign as the argument. It stores the integral part in the object |

[^21]|  | pointed to by iptr. |
| :---: | :---: |
| floatn nan (uintn nancode) float nan (uint nancode) <br> doublen nan (ulongn nancode) double nan (ulong nancode) | Returns a quiet NaN. The nancode may be placed in the significand of the resulting NaN . |
| gentype nextafter (gentype $x$, gentype $y$ ) | Computes the next representable single-precision floating-point value following $x$ in the direction of $y$. Thus, if $y$ is less than $x$, nextafter() returns the largest representable floating-point number less than $x$. |
| gentype pow (gentype $x$, gentype $y$ ) | Compute $x$ to the power $y$. |
| float $n$ pown (float $n x$, int $n y$ ) float pown (float $x$, int $y$ ) <br> doublen pown (doublen $x$, int $n y$ ) double pown (double $x$, int $y$ ) | Compute $x$ to the power $y$, where $y$ is an integer. |
| gentype powr (gentype $x$, gentype $y$ ) | Compute $x$ to the power $y$, where $x$ is $>=0$. |
| gentype remainder (gentype $x$, gentype $y$ ) | Compute the value $r$ such that $r=x-n^{*} y$, where $n$ is the integer nearest the exact value of $x / y$. If there are two integers closest to $x / y, n$ shall be the even one. If $r$ is zero, it is given the same sign as $x$. |
| float $n$ remquo (float $n x$, float $n y$, int $n * q u o$ ) float remquo (float $x$, float $y$, int *quo) | The remquo function computes the value r such that $r=x-k^{*} y$, where $k$ is the integer nearest the exact value of $x / y$. If there are two integers closest to $x / y, k$ shall be the even one. If $r$ is zero, it is given the same sign as $x$. This is the same value that is returned by the remainder function. remquo also calculates the lower seven bits of the integral quotient $x / y$, and gives that value the same sign as $x / y$. It stores this signed value in the object pointed to by quo. |
| double $n$ remquo (double $n x$, double $y$, int $n$ *quo) double remquo (double $x$, double $y$, int *quo) | The remquo function computes the value r such that $r=x-k^{*} y$, where $k$ is the integer nearest the exact value of $x / y$. If there are two integers closest to $x / y, k$ shall be the even one. If $r$ is zero, it is given the same sign as $x$. This is the same value that is returned by the remainder function. remquo also calculates the lower seven bits of the integral quotient $x / y$, and gives that value the same sign as $x / y$. It stores this signed value in the object pointed to by quo. |
| gentype rint (gentype) | Round to integral value (using round to nearest even rounding mode) in floating-point format. Refer to section 7.1 for description of rounding modes. |


| float $n$ rootn (float $n x$, int $n y$ ) <br> float rootn (float $x$, int $y$ ) <br> doublen rootn (doublen $x$, int $n$ ) <br> doublen rootn (double $x$, int $y$ ) | Compute $x$ to the power 1/y. |
| :--- | :--- |
| gentype round (gentype $x$ ) | Return the integral value nearest to $x$ rounding <br> halfway cases away from zero, regardless of the <br> current rounding direction. |
| gentype rsqrt (gentype) | Compute inverse square root. |
| gentype sin (gentype $x$ ) | Compute sine, where $x$ is an angle in radians. |
| gentype sincos (gentype $x$, |  |
| gentype *cosval) | Compute sine and cosine of x. The computed sine <br> is the return value and computed cosine is returned <br> in cosval, where $x$ is an angle in radians |
| gentype sinh (gentype $x$ ) | Compute hyperbolic sine, where $x$ is an angle in <br> radians |
| gentype sinpi (gentype $x$ ) | Compute sin $(\pi x)$. |
| gentype sqrt (gentype) | Compute square root. |
| gentype tan (gentype $x$ ) | Compute tangent, where $x$ is an angle in radians. <br> Compute hyperbolic tangent, where $x$ is an angle in <br> radians. |
| gentype tanh (gentype $x$ ) | Compute tan $(\pi x$ ). |
| gentype tanpi (gentype $x$ ) | Compute the gamma function. |
| gentype tgamma (gentype) | Round to integral value using the round to zero <br> rounding mode. |
| gentype trunc (gentype) |  |

Table 6.8 Scalar and Vector Argument Built-in Math Function Table

Table 6.9 describes the following functions:

* A subset of functions from table 6.8 that are defined with the half_ prefix. These functions are implemented with a minimum of 10-bits of accuracy i.e. an ULP value <= 8192 ulp.

4 A subset of functions from table 6.8 that are defined with the native_ prefix. These functions may map to one or more native device instructions and will typically have better performance compared to the corresponding functions (without the native prefix) described in table 6.8. The accuracy (and in some cases the input range(s)) of these functions is implementation-defined.
half_ and native_functions for following basic operations: divide and reciprocal.
We use the generic type name gentype to indicate that the functions in table 6.9 can take float, float2, float3, float4, float8 or float16 as the type for the arguments.

| Function | Description |
| :---: | :---: |
| gentype half_cos (gentype $x$ ) | Compute cosine. $x$ is an angle in radians and it must be in the range $-2^{16} \ldots+2^{16}$. |
| gentype half_divide (gentype $x$, gentype $y$ ) | Compute $x / y$. |
| gentype half_exp (gentype $x$ ) | Compute the base- e exponential of $x$. |
| gentype half_exp2 (gentype $x$ ) | Compute the base- 2 exponential of $x$. |
| gentype half_exp 10 (gentype $x$ ) | Compute the base- 10 exponential of $x$. |
| gentype half_log (gentype $x$ ) | Compute natural logarithm. |
| gentype half_log2 (gentype $x$ ) | Compute a base 2 logarithm. |
| gentype half_log10 (gentype $x$ ) | Compute a base 10 logarithm. |
| gentype half_powr (gentype $x$, gentype $y$ ) | Compute $x$ to the power $y$, where $x$ is $>=0$. |
| gentype half_recip (gentype $x$ ) | Compute reciprocal. |
| gentype half_rsqrt (gentype $x$ ) | Compute inverse square root. |
| gentype half_sin (gentype $x$ ) | Compute sine. $x$ is an angle in radians and it must be in the range $-2^{16} \ldots+2^{16}$. |
| gentype half_sqrt (gentype $x$ ) | Compute square root. |
| gentype half_tan (gentype $x$ ) | Compute tangent. $x$ is an angle in radians and it must be in the range $-2^{16} \ldots+2^{16}$. |
| gentype native_cos (gentype $x$ ) | Compute cosine over an implementation-defined range, where $x$ is an angle in radians. The maximum error is implementation-defined. |
| gentype native_divide (gentype $x$, gentype $y$ ) | Compute $x / y$ over an implementation-defined range. The maximum error is implementation-defined. |
| gentype native_exp (gentype $x$ ) | Compute the base- e exponential of $x$ over an implementation-defined range. The maximum error is implementation-defined. |
| gentype native_exp2 (gentype $x$ ) | Compute the base- 2 exponential of $x$ over an implementation-defined range. The maximum error is implementation-defined. |
| gentype native_exp10 (gentype $x$ ) | Compute the base- 10 exponential of $x$ over an implementation-defined range. The maximum error is implementation-defined. |
| gentype native_log (gentype $x$ ) | Compute natural logarithm over an implementationdefined range. The maximum error is implementationdefined. |
| gentype native_log2 (gentype $x$ ) | Compute a base 2 logarithm over an implementationdefined range. The maximum error is implementationdefined. |
| gentype native_log10 (gentype $x$ ) | Compute a base 10 logarithm over an implementationdefined range. The maximum error is implementationdefined. |
| gentype native_powr (gentype $x$, | Compute $x$ to the power $y$, where $x$ is $>=0$. The range of |


| gentype $y$ ) | $x$ and $y$ are implementation-defined. The maximum error <br> is implementation-defined. |
| :--- | :--- |
| gentype native_recip (gentype $x$ ) | Compute reciprocal over an implementation-defined <br> range. The maximum error is implementation-defined. |
| gentype native_rsqrt (gentype $x$ ) | Compute inverse square root over an implementation- <br> defined range. The maximum error is implementation- <br> defined. |
| gentype native_sin (gentype $x$ ) | Compute sine over an implementation-defined range, <br> where $x$ is an angle in radians. The maximum error is <br> implementation-defined. |
| gentype native_sqrt (gentype $x$ ) | Compute square root over an implementation-defined <br> range. The maximum error is implementation-defined. |
| gentype native_tan (gentype $x$ ) | Compute tangent over an implementation-defined range, <br> where $x$ is an angle in radians. The maximum error is <br> implementation-defined. |

## Table 6.9 Scalar and Vector Argument Built-in half__ and native__ Math Functions

Support for denormal values is optional for half_functions. The half_ functions may return any result allowed by section 7.5.3, even when -cl-denorms-are-zero (see section 5.8.4.2) is not in force. Support for denormal values is implementation-defined for native_functions.

The following symbolic constants are available. Their values are of type float and are accurate within the precision of a single precision floating-point number.

| Constant Name | Description |
| :--- | :--- |
| MAXFLOAT | Value of maximum non-infinite single-precision floating-point <br> number. |
| HUGE_VALF | A positive float constant expression. HUGE_VALF evaluates <br> to +infinity. Used as an error value returned by the built-in <br> math functions. |
| INFINITY | A constant expression of type float representing positive or <br> unsigned infinity. |
| NAN | A constant expression of type float representing a quiet NaN. |

If double precision is supported by the device, the following symbolic constant will also be available:

| Constant Name | Description |
| :--- | :--- |
| HUGE_VAL | A positive double constant expression. HUGE_VAL evaluates <br> to +infinity. Used as an error value returned by the built-in <br> math functions. |

### 6.13.2.1 Floating-point macros and pragmas

The FP_CONTRACT pragma can be used to allow (if the state is on) or disallow (if the state is $\circ f f$ ) the implementation to contract expressions. Each pragma can occur either outside external declarations or preceding all explicit declarations and statements inside a compound statement. When outside external declarations, the pragma takes effect from its occurrence until another FP_CONTRACT pragma is encountered, or until the end of the translation unit. When inside a compound statement, the pragma takes effect from its occurrence until another FP_CONTRACT pragma is encountered (including within a nested compound statement), or until the end of the compound statement; at the end of a compound statement the state for the pragma is restored to its condition just before the compound statement. If this pragma is used in any other context, the behavior is undefined.

The pragma definition to set $\mathbf{F P}$ _CONTRACT is:

```
#pragma OPENCL FP_CONTRACT on-off-switch
on-off-switch is one of:
    ON, OFF or DEFAULT.
    The DEFAULT value is ON.
```

The FP_FAST_FMAF macro indicates whether the fma function is fast compared with direct code for single precision floating-point. If defined, the FP_FAST_FMAF macro shall indicate that the fma function generally executes about as fast as, or faster than, a multiply and an add of float operands.

The macro names given in the following list must use the values specified. These constant expressions are suitable for use in \#if preprocessing directives.

```
#define FLT_DIG 6
#define FLT_MANT_DIG 24
#define FLT_MAX_10_EXP +38
#define FLT_MAX_EXP +128
#define FLT MIN 10 EXP -37
#define FLT_MIN_EXP -125
#define FLT_RADIX 2
#define FLT_MAX 0x1.fffffep127f
#define FLT_MIN 0x1.0p-126f
#define FLT_EPSILON 0x1.0p-23f
```

The following table describes the built-in macro names given above in the OpenCL C programming language and the corresponding macro names available to the application.

| Macro in OpenCL Language | Macro for application |
| :---: | :---: |
| FLT_DIG | CL_FLT_DIG |
| FLT_MANT_DIG | CL_FLT_MANT_DIG |


| FLT_MAX_10_EXP | CL_FLT_MAX_10_EXP |
| :---: | :---: |
| FLT_MAX_EXP | CL_FLT_MAX_EXP |
| FLT_MIN_10_EXP | CL_FLT_MIN_10_EXP $^{\text {FLT_MI }}$ |
| FLT_MIN_EXP | CL_FLT_MIN_EXP |
| FLT_RADIX | CL_FLT_RADIX |
| FLT_MAX | CL_FLT_MAX |
| FLT_MIN | CL_FLT_MIN |
| FLT_EPSILSON | CL_FLT_EPSILON |

The following macros shall expand to integer constant expressions whose values are returned by $\operatorname{ilogb}(x)$ if $x$ is zero or NaN, respectively. The value of $\mathbf{F P}$ _ILOGB0 shall be either $\{$ INT_MIN $\}$ or - \{INT_MAX\}. The value of FP_ILOGBNAN shall be either \{INT_MAX\} or \{INT_MIN\}.

The following constants are also available. They are of type float and are accurate within the precision of the float type.

| Constant | Description |
| :--- | :--- |
| M_E_F | Value of e |
| M_LOG2E_F | Value of $\log _{2} \mathrm{e}$ |
| M_LOG10E_F | Value of $\log _{10} \mathrm{e}$ |
| M_LN2_F | Value of $\log _{\mathrm{e}} 2$ |
| M_LN10_F | Value of $\log _{e} 10$ |
| M_PI_F | Value of $\pi$ |
| M_PI_2_F | Value of $\pi / 2$ |
| M_PI_4_F | Value of $\pi / 4$ |
| M_1_PI_F | Value of $1 / \pi$ |
| M_2_PI_F | Value of $2 / \pi$ |
| M_2_SQRTPI_F | Value of $2 / \sqrt{\pi}$ |
| M_SQRT2_F | Value of $\sqrt{ } 2$ |
| M_SQRT1_2_F | Value of $1 / \sqrt{2}$ |

If double precision is supported by the device, the following macros and constants are also available:

The FP_FAST_FMA macro indicates whether the fma() family of functions are fast compared with direct code for double precision floating-point. If defined, the FP_FAST_FMA macro shall indicate that the fma() function generally executes about as fast as, or faster than, a multiply and an add of double operands

The macro names given in the following list must use the values specified. These constant expressions are suitable for use in \#if preprocessing directives.

```
#define DBL_DIG 15
#define DBL_MANT_DIG 53
```

```
#define DBL_MAX_10_EXP +308
#define DBL_MAX_EXP +1024
#define DBL_MIN_10_EXP -307
#define DBL_MIN_EXP - -1021
#define DBL_MAX 0x1.fffffffffffffp1023
#define DBL_MIN 0x1.0p-1022
#define DBL_EPSILON 0x1.0p-52
```

The following table describes the built-in macro names given above in the OpenCL C programming language and the corresponding macro names available to the application.

| Macro in OpenCL Language | Macro for application |
| :---: | :---: |
| DBL_DIG | CL_DBL_DIG |
| DBL_MANT_DIG | CL_DBL_MANT_DIG |
| DBL_MAX_10_EXP | CL_DBL_MAX_10_EXP |
| DBL_MAX_EXP | CL_DBL_MAX_EXP |
| DBL_MIN_10_EXP | CL_DBL_MIN_10_EXP |
| DBL_MIN_EXP | CL_DBL_MIN_EXP |
| DBL_MAX | CL_DBL_MAX |
| DBL_MIN | CL_DBL_MIN |
| DBL_EPSILSON | CL_DBL_EPSILON |

The following constants are also available. They are of type double and are accurate within the precision of the double type.

| Constant | Description |
| :--- | :--- |
| M_E | Value of e |
| M_LOG2E | Value of $\log _{2} \mathrm{e}$ |
| M_LOG10E | Value of $\log _{10} \mathrm{e}$ |
| M_LN2 | Value of $\log _{e} 2$ |
| M_LN10 | Value of $\log _{e} 10$ |
| M_PI | Value of $\pi$ |
| M_PI_2 | Value of $\pi / 2$ |
| M_PI_4 | Value of $\pi / 4$ |
| M_1_PI | Value of $1 / \pi$ |
| M_2_PI | Value of $2 / \pi$ |
| M_2_SQRTPI | Value of $2 / \sqrt{\pi}$ |
| M_SQRT2 | Value of $\sqrt{ } 2$ |
| M_SQRT1_2 | Value of $1 / \sqrt{ } 2$ |

### 6.13.3 Integer Functions

Table 6.10 describes the built-in integer functions that take scalar or vector arguments. The vector versions of the integer functions operate component-wise. The description is percomponent.

We use the generic type name gentype to indicate that the function can take char, char $\{2|3| 4|8| 16\}$, uchar, uchar $\{2|3| 4|8| 16\}$, short, short $\{2|3| 4|8| 16\}$, ushort, $u s h o r t\{2|3| 4|8| 16\}$, int, int $\{2|3| 4|8| 16\}$, uint, $u i n t\{2|3| 4|8| 16\}$, long, $\operatorname{long}\{2|3| 4|8| 16\}$ ulong, or ulong $\{2|3| 4|8| 16\}$ as the type for the arguments. We use the generic type name ugentype to refer to unsigned versions of gentype. For example, if gentype is char4, ugentype is uchar4. We also use the generic type name sgentype to indicate that the function can take a scalar data type i.e. char, uchar, short, ushort, int, uint, long, or ulong as the type for the arguments. For built-in integer functions that take gentype and sgentype arguments, the gentype argument must be a vector or scalar version of the sgentype argument. For example, if sgentype is uchar, gentype must be uchar or uchar $\{2|3| 4|8| 16\}$. For vector versions, sgentype is implicitly widened to gentype as described in section 6.3.a.

For any specific use of a function, the actual type has to be the same for all arguments and the return type unless otherwise specified.

| Function | Description |
| :--- | :--- |
| ugentype abs (gentype $x$ ) | Returns $\|\mathrm{x}\|$. |
| ugentype abs_diff (gentype $x$, gentype $y$ ) | Returns $\|\mathrm{x}-\mathrm{y}\|$ without modulo overflow. |
| gentype add_sat (gentype $x$, gentype $y$ ) | Returns $x+y$ and saturates the result. |
| gentype hadd (gentype $x$, gentype $y$ ) | Returns $(x+y) \gg 1$. The intermediate sum does <br> not modulo overflow. |
| gentype rhadd (gentype $x$, gentype $y)^{33}$ | Returns $(x+y+1) \gg 1$. The intermediate sum <br> does not modulo overflow. |
| gentype clamp (gentype $x$, <br> gentype minval, <br> gentype maxval) | Returns min(max( $x$, minval), maxval). <br> gentype clamp (gentype $x$, <br> sgentype minval, <br> sgentype maxval) |
| Results are undefined if minval > maxval. |  |
| gentype clz (gentype $x$ ) | Returns the number of leading 0-bits in $x$, starting <br> at the most significant bit position. If $x$ is 0, <br> returns the size in bits of the type of $x$ or |

[^22]|  | component type of $x$, if $x$ is a vector. |
| :---: | :---: |
| gentype ctz (gentype $x$ ) | Returns the count of trailing 0 -bits in $x$. If $x$ is 0 , returns the size in bits of the type of $x$ or component type of $x$, if $x$ is a vector. |
| gentype mad_hi (gentype $a$, gentype $b$, gentype $c$ ) | Returns mul_hi $(a, b)+c$. |
| gentype mad_sat (gentype $a$, gentype $b$, gentype $c$ ) | Returns $a * b+c$ and saturates the result. |
| gentype max (gentype $x$, gentype $y$ ) <br> gentype max (gentype $x$, sgentype $y$ ) | Returns $y$ if $x<y$, otherwise it returns $x$. |
| gentype min (gentype $x$, gentype $y$ ) <br> gentype min (gentype $x$, sgentype $y$ ) | Returns $y$ if $y<x$, otherwise it returns $x$. |
| gentype mul_hi (gentype $x$, gentype $y$ ) | Computes $x * y$ and returns the high half of the product of $x$ and $y$. |
| gentype rotate (gentype $v$, gentype $i$ ) | For each element in $v$, the bits are shifted left by the number of bits given by the corresponding element in $i$ (subject to usual shift modulo rules described in section 6.3). Bits shifted off the left side of the element are shifted back in from the right. |
| gentype sub_sat (gentype $x$, gentype $y$ ) | Returns $x-y$ and saturates the result. |
| short upsample (char hi, uchar lo) ushort upsample (uchar hi, uchar lo) shortn upsample (charn hi, ucharn lo) ushortn upsample (ucharn hi, ucharn lo) <br> int upsample (short hi, ushort lo) uint upsample (ushort hi, ushort lo) int $n$ upsample (shortn hi, ushortn lo) uintn upsample (ushortn hi, ushortn lo) <br> long upsample (int hi, uint lo) ulong upsample (uint hi, uint lo) longn upsample (intn hi, uint $n l o$ ) ulongn upsample (uintn hi, uintn lo) | $\begin{aligned} & \operatorname{result}[\mathrm{i}]=((\text { short }) h i[\mathrm{i}] \ll 8) \mid l o[\mathrm{i}] \\ & \operatorname{result}[\mathrm{i}]=((\text { ushort }) h i[\mathrm{i}] \ll 8) \mid l o[\mathrm{i}] \\ & \operatorname{result}[\mathrm{i}]=((\mathrm{int}) h i[\mathrm{i}] \ll 16) \mid l o[\mathrm{i}] \\ & \operatorname{result}[\mathrm{i}]=((\text { uint }) h i[\mathrm{i}] \ll 16) \mid l o[\mathrm{i}] \\ & \operatorname{result}[\mathrm{i}]=((\operatorname{long}) h i[\mathrm{i}] \ll 32) \mid l o[\mathrm{i}] \\ & \operatorname{result}[\mathrm{i}]=((\text { ulong }) h i[\mathrm{i}] \ll 32) \mid l o[\mathrm{i}] \end{aligned}$ |
| gentype popcount (gentype $x$ ) | Returns the number of non-zero bits in $x$. |

Table 6.10 Scalar and Vector Integer Argument Built-in Functions
Table 6.11 describes fast integer functions that can be used for optimizing performance of kernels. We use the generic type name gentype to indicate that the function can take int, int2, int3, int4, int8, int16, uint, uint2, uint3, uint4, uint8 or uint16 as the type for the arguments.

| Function | Description |
| :--- | :--- |
| gentype mad24 (gentype $x$, <br> gentype $y$, gentype z) | Multipy two 24-bit integer values $x$ and $y$ and add <br> the 32-bit integer result to the 32-bit integer $z$. <br> Refer to definition of mul24 to see how the 24-bit <br> integer multiplication is performed. |
| gentype mul24 (gentype $x$, gentype $y$ ) | Multiply two 24-bit integer values $x$ and $y . x$ and $y$ <br> are 32-bit integers but only the low 24-bits are used <br> to perform the multiplication. mul24 should only |
|  | be used when values in $x$ and $y$ are in the range [- |
|  | $\left.2^{23,}, 2^{23}-1\right]$ if $x$ and $y$ are signed integers and in the |
| range [0, 24-1] if $x$ and $y$ are unsigned integers. If |  |
|  | $x$ and $y$ are not in this range, the multiplication |
| result is implementation-defined. |  |

## Table 6.11 Fast Integer Built-in Functions

The macro names given in the following list must use the values specified. The values shall all be constant expressions suitable for use in \#if preprocessing directives.

```
#define CHAR BIT 8
#define CHAR_MAX SCHAR_MAX
#define CHAR_MIN SCHAR_MIN
#define INT_MAX 2147483647
#define INT_MIN (-2147483647 - 1)
#define LONG_MAX 0x7fffffffffffffffL
#define LONG_MIN (-0x7fffffffffffffffl - 1)
#define SCHA\overline{R}MMAX 127
#define SCHAR_MIN (-127 - 1)
#define SHRT_M
#define SHRT_MIN (-32767 - 1)
#define UCHAR MAX 255
#define USHRT_MAX 65535
#define UINT_\overline{MAX Oxffffffff}
#define ULONG_MAX 0xffffffffffffffffuL
```

The following table describes the built-in macro names given above in the OpenCL C programming language and the corresponding macro names available to the application.

| Macro in OpenCL Language | Macro for application |
| :---: | :---: |
| CHAR_BIT | CL_CHAR_BIT |
| CHAR_MAX | CL_CHAR_MAX |
| CHAR_MIN | CL_CHAR_MIN |
| INT_MAX | CL_INT_MAX |
| INT_MIN | CL_INT_MIN |
| LONG_MAX | CL_LONG_MAX |


| LONG_MIN | CL_LONG_MIN |
| :---: | :---: |
| SCHAR_MAX | CL_SCHAR_MAX |
| SCHAR_MIN | CL_SCHAR_MIN |
| SHRT_MAX | CL_SHRT_MAX |
| SHRT_MIN | CL_SHRT_MIN |
| UCHAR_MAX | CL_UCHAR_MAX |
| USHRT_MAX | CL_USHRT_MAX |
| UINT_MAX | CL_UINT_MAX |
| ULONG_MAX | CL_ULONG_MAX |

### 6.13.4 Common Functions ${ }^{34}$

Table 6.12 describes the list of built-in common functions. These all operate component-wise. The description is per-component. We use the generic type name gentype to indicate that the function can take float, float2, float3, float4, float8, float16, double, double2, double3, double4, double8 or double16 as the type for the arguments. We use the generic type name gentypef to indicate that the function can take float, float2, float3, float4, float8, or float16 as the type for the arguments. We use the generic type name gentyped to indicate that the function can take double, double2, double3, double4, double8 or double16 as the type for the arguments.

The built-in common functions are implemented using the round to nearest even rounding mode.

| Function | Description |
| :--- | :--- |
| gentype clamp (gentype $x$, <br> gentype minval, <br> gentype maxval) | Returns fmin(fmax(x, minval), maxval). |
| gentypef clamp (gentypef $x$, |  |
| float minval, |  |
| float maxval) |  |$\quad$| Results are undefined if minval $>$ maxval. |
| :--- |

[^23]| ```gentypef mix (gentypef x, gentypef y, float a) gentyped mix (gentyped x, gentyped y, double a)``` | $a$ must be a value in the range $0.0 \ldots 1.0$. If $a$ is not in the range $0.0 \ldots 1.0$, the return values are undefined. |
| :---: | :---: |
| gentype radians (gentype degrees) | Converts degrees to radians, i.e. ( $\pi / 180$ ) * degrees. |
| gentype step (gentype edge, gentype $x$ ) gentypef step (float edge, gentypef $x$ ) gentyped step (double edge, gentyped $x$ ) | Returns 0.0 if $x<e d g e$, otherwise it returns 1.0. |
| gentype smoothstep (gentype edge0,gentype edgel,gentype $x)$ gentypef smoothstep (float edge0,float edgel, <br> gentypef $x)$ <br> gentyped smoothstep (double edge 0, <br> double edgel, <br> gentyped $x)$ | Returns 0.0 if $x<=e d g e 0$ and 1.0 if $x>=e d g e 1$ and performs smooth Hermite interpolation between 0 and 1 when edge $0<x<$ edge 1 . This is useful in cases where you would want a threshold function with a smooth transition. <br> This is equivalent to: gentype t ; $\mathrm{t}=$ clamp $((\mathrm{x}-$ edge 0$) /($ edge $1-$ edge 0$), 0,1) ;$ return $\mathrm{t} * \mathrm{t} *(3-2 * \mathrm{t})$; <br> Results are undefined if edge $0>=e d g e 1$ or if $x$, edge0 or edgel is a NaN . |
| gentype sign (gentype $x$ ) | Returns 1.0 if $x>0,-0.0$ if $x=-0.0,+0.0$ if $x=$ +0.0 , or -1.0 if $x<0$. Returns 0.0 if $x$ is a NaN. |

Table 6.12 Scalar and Vector Argument Built-in Common Function Table

### 6.13.5 Geometric Functions ${ }^{35}$

Table 6.13 describes the list of built-in geometric functions. These all operate component-wise. The description is per-component. floatn is float, float2, float 3 , or float 4 and doublen is double, double2, double3, or double4. The built-in geometric functions are implemented using the round to nearest even rounding mode.

| Function | Description |
| :---: | :---: |
| float4 cross (float4 p0, float4 pl) <br> float 3 cross (float $3 p 0$, float $3 p 1$ ) <br> double 4 cross (double 4 p0, double 4 p1) <br> double 3 cross (double $3 p 0$, double $3 p 1$ ) | Returns the cross product of $p 0 . x y z$ and p1.xyz. The $w$ component of float 4 result returned will be 0.0 . |
| float dot (float $n p 0$, float $n p 1$ ) <br> double dot (doublen $p 0$, doublen $p 1$ ) | Compute dot product. |
| float distance (float $n p 0$, float $n p 1$ ) double distance (doublen $p 0$, doublen $p 1$ ) | Returns the distance between $p 0$ and $p 1$. This is calculated as length $(p 0-p 1)$. |
| float length (float $n p$ ) <br> double length (doublen $p$ ) | Return the length of vector $p$, i.e., $\sqrt{p \cdot x^{2}+p \cdot y^{2}+\ldots}$ |
| float $n$ normalize (float $n p$ ) $\qquad$ | Returns a vector in the same direction as $p$ but with a length of 1. |
| float fast_distance (float $n p 0$, floatn $p$ 1) | Returns fast_length $(p 0-p 1)$. |
| float fast_length (float $n p$ ) | Returns the length of vector $p$ computed as: half_sqrt $\left(p . x^{2}+p . y^{2}+\ldots.\right)$ |
| float $n$ fast_normalize (float $n p$ ) | Returns a vector in the same direction as $p$ but with a length of 1 . fast_normalize is computed as: $p^{*} \text { half_rsqrt }\left(p \cdot x^{2}+p \cdot y^{2}+\ldots\right)$ <br> The result shall be within 8192 ulps error from the infinitely precise result of |

[^24]|  | ```if \((\boldsymbol{\operatorname { l l l }}(p==0.0 \mathrm{f}))\) result \(=p\); else result \(=p / \mathbf{s q r t}\left(p \cdot x^{2}+p \cdot y^{2}+\ldots\right) ;\)``` <br> with the following exceptions: <br> 1) If the sum of squares is greater than FLT_MAX then the value of the floating-point values in the result vector are undefined. <br> 2) If the sum of squares is less than FLT_MIN then the implementation may return back $p$. <br> 3) If the device is in "denorms are flushed to zero" mode, individual operand elements with magnitude less than $\mathbf{~ s q r t ( F L T \_ M I N ) ~ m a y ~ b e ~ f l u s h e d ~ t o ~ z e r o ~}$ before proceeding with the calculation. |
| :---: | :---: |

Table 6.13 Scalar and Vector Argument Built-in Geometric Function Table

### 6.13.6 Relational Functions

The relational and equality operators ( $\langle,\langle=,>,>=,!=,==$ ) can be used with scalar and vector built-in types and produce a scalar or vector signed integer result respectively as described in section 6.3.

The functions ${ }^{36}$ described in table 6.14 can be used with built-in scalar or vector types as arguments and return a scalar or vector integer result. The argument type gentype refers to the following built-in types: char, charn, uchar, ucharn, short, shortn, ushort, ushortn, int, intn, uint, uintn, long, longn, ulong, ulongn, float, floatn, double, and doublen. The argument type igentype refers to the built-in signed integer types i.e. char, charn, short, shortn, int, int $n$, long and longn. The argument type ugentype refers to the built-in unsigned integer types i.e. uchar, ucharn, ushort, ushortn, uint, uint $n$, ulong and ulongn. $n$ is $2,3,4$, 8 , or 16 .

The functions isequal, isnotequal, isgreater, isgreaterequal, isless, islessequal, islessgreater, isfinite, isinf, isnan, isnormal, isordered, isunordered and signbit described in table 6.14 shall return a 0 if the specified relation is false and a 1 if the specified relation is true for scalar argument types. These functions shall return a 0 if the specified relation is false and a -1 (i.e. all bits set) if the specified relation is true for vector argument types.

The relational functions isequal, isgreater, isgreaterequal, isless, islessequal, and islessgreater always return 0 if either argument is not a number ( NaN ). isnotequal returns 1 if one or both arguments are not a number $(\mathrm{NaN})$ and the argument type is a scalar and returns -1 if one or both arguments are not a number ( NaN ) and the argument type is a vector.

| Function | Description |
| :--- | :--- |
| int isequal (float $x$, float $y$ ) <br> int $n$ isequal (float $n x$, float $n y$ ) | Returns the component-wise compare of $x==y$. |
| int isequal (double $x$, double $y$ ) |  |
| long $n$ isequal (double $x$, double $n y$ ) |  |$\quad$.

[^25]| int $n$ isgreater (float $n x$, float $n y$ ) <br> int isgreater (double $x$, double $y$ ) <br> long $n$ isgreater (double $n x$, doublen $y$ ) |  |
| :---: | :---: |
| int isgreaterequal (float $x$, float $y$ )int $n$ isgreaterequal (float $n x$, float $n y$ ) int isgreaterequal (double $x, ~ 子$double $y$ ) <br> long $n$ isgreaterequal (double $x$, <br> double $n y$ ) | Returns the component-wise compare of $x>=y$. |
| int isless (float $x$, float $y$ ) <br> int $n$ isless (floatn $x$, floatn $y$ ) <br> int isless (double $x$, double $y$ ) <br> long $n$ isless (double $n x$, doublen $y$ ) | Returns the component-wise compare of $x<y$. |
| int islessequal (float $x$, float $y$ ) int $n$ islessequal (float $n x$, float $n y$ ) <br> int islessequal (double $x$, double $y$ ) <br> long $n$ islessequal (double $n x$, double $n y$ ) | Returns the component-wise compare of $x<=y$. |
| int islessgreater (float $x$, float $y$ ) intn islessgreater (float $n x$, floatn $y$ ) <br> int islessgreater (double $x$, double $y$ ) long $n$ islessgreater (double $x$, doublen $y$ ) | Returns the component-wise compare of $(x<y) \\|(x>y)$. |
| int isfinite (float) int $n$ isfinite (float $n$ ) <br> int isfinite (double) long $n$ isfinite (doublen) | Test for finite value. |
| int isinf (float) int $n \operatorname{isinf}$ (float $n$ ) <br> int isinf (double) long $n \operatorname{isinf}$ (doublen) | Test for infinity value (positive or negative) |
| int isnan (float) int $n$ isnan (float $n$ ) <br> int isnan (double) longn isnan (doublen) | Test for a NaN. |
| int isnormal (float) int $n$ isnormal (floatn) <br> int isnormal (double) | Test for a normal value. |


| long $n$ isnormal (double $n$ ) |  |
| :---: | :---: |
| int isordered (float $x$, float $y$ ) <br> int $n$ isordered (float $n x$, float $n y$ ) <br> int isordered (double $x$, double $y$ ) <br> long $n$ isordered (double $n x$, doublen $y$ ) | Test if arguments are ordered. isordered() takes arguments $x$ and $y$, and returns the result isequal $(x, x) \& \&$ isequal $(y, y)$. |
| int isunordered (float $x$, float $y$ ) int $n$ isunordered (float $n x$, float $n y$ ) <br> int isunordered (double $x$, double $y$ ) long $n$ isunordered (double $n x$, double $y$ ) | Test if arguments are unordered. isunordered() takes arguments $x$ and $y$, returning non-zero if $x$ or $y$ is NaN , and zero otherwise. |
| int signbit (float) int $n$ signbit (float $n$ ) <br> int signbit (double) long $n$ signbit (doublen) | Test for sign bit. The scalar version of the function returns a 1 if the sign bit in the float is set else returns 0 . The vector version of the function returns the following for each component in floatn: -1 (i.e all bits set) if the sign bit in the float is set else returns 0 . |
| int any (igentype $x$ ) | Returns 1 if the most significant bit in any component of $x$ is set; otherwise returns 0 . |
| int all (igentype $x$ ) | Returns 1 if the most significant bit in all components of $x$ is set; otherwise returns 0 . |
| gentype bitselect (gentype $a$, gentype $b$, gentype $c$ ) | Each bit of the result is the corresponding bit of $a$ if the corresponding bit of $c$ is 0 . Otherwise it is the corresponding bit of $b$. |
|  | For each component of a vector type, result $[i]=$ if MSB of $c[i]$ is set $? b[i]: a[i]$. <br> For a scalar type, result $=c ? b: a$. <br> igentype and ugentype must have the same number of elements and bits as gentype. <br> NOTE: The above definition means that the behavior of select and the ternary operator for vector and scalar types is dependent on different interpretations of the bit pattern of $c$. |

Table 6.14 Scalar and Vector Relational Functions

### 6.13.7 Vector Data Load and Store Functions

Table 6.15 describes the list of supported functions that allow you to read and write vector types from a pointer to memory. We use the generic type gentype to indicate the built-in data types char, uchar, short, ushort, int, uint, long, ulong, float or double. We use the generic type name gentypen to represent $n$-element vectors of gentype elements. We use the type name halfn to represent $n$-element vectors of half elements ${ }^{37}$. The suffix $n$ is also used in the function names (i.e. vload $n$, vstoren etc.), where $n=$ $2,3,4,8$ or 16 .

| Function | Description |
| :---: | :---: |
| ```gentypen vloadn (size_t offset, const gentype *p) gentypen vloadn (size_t offset, const constant gentype *p)``` | Return sizeof (gentypen) bytes of data read from address $(p+($ offset $* n))$. The address computed as $(p+($ offset $* n))$ must be 8 -bit aligned if gentype is char, uchar; 16-bit aligned if gentype is short, ushort; 32-bit aligned if gentype is int, uint, float; 64-bit aligned if gentype is long, ulong. |
| void vstoren (gentypen data, size_t offset, gentype *p) | Write sizeof (gentypen) bytes given by data to address $(p+($ offset $* n))$. The address computed as $(p+($ offset $* n))$ must be 8 -bit aligned if gentype is char, uchar; 16-bit aligned if gentype is short, ushort; 32-bit aligned if gentype is int, uint, float; 64-bit aligned if gentype is long, ulong. |
| ```float vload_half (size_t offset, const half *p) float vload_half (size_t offset, const constant half *p)``` | Read sizeof (half) bytes of data from address ( $p+$ offset). The data read is interpreted as a half value. The half value is converted to a float value and the float value is returned. The read address computed as ( $p+$ offset) must be 16 -bit aligned. |
| ```floatn vload_halfn (size_t offset, const half *p) floatn vload_half}n\mathrm{ (size_t offset, const constant half *p)``` | Read sizeof (halfn) bytes of data from address $(p+($ offset $* n))$. The data read is interpreted as a half $n$ value. The half $n$ value read is converted to a float $n$ value and the float $n$ value is returned. The read address computed as ( $\mathrm{p}+($ offset $* n$ ) ) must be 16 -bit aligned. |

[^26]| ```void vstore_half (float data, size_t offset, half *p) void vstore_half_rte (float data, size_t offset, half *p) void vstore_half_rtz (float data, size_t offset, half *p) void vstore_half_rtp (float data, size_t offset, half *p) void vstore_half_rtn (float data, size_t offset, half *p)``` | The float value given by data is first converted to a half value using the appropriate rounding mode. The half value is then written to address computed as $(p+$ offset). The address computed as ( $p+$ offset) must be 16-bit aligned. <br> vstore_half uses the default rounding mode. The default rounding mode is round to nearest even. |
| :---: | :---: |
| ```void vstore_half}\boldsymbol{n}\mathrm{ (float n data, size_t offset, half *p) void vstore_halfn_rte (floatn data, size_t offset, half *p) void vstore_halfn_rtz (floatn data, size_t offset, half *p) void vstore_halfn_rtp (floatn data, size_t offset, half *p) void vstore_halfn_rtn (floatn data, size_t offset, half *p)``` | The float $n$ value given by data is converted to a half $n$ value using the appropriate rounding mode. The half $n$ value is then written to address computed as ( $p+$ (offset *n)). The address computed as ( $p+$ (offset * $n$ )) must be 16 -bit aligned. <br> vstore_half $\boldsymbol{n}$ uses the default rounding mode. The default rounding mode is round to nearest even. |
| void vstore_half (double data, size_t offset, half *p) void vstore_half_rte (double data, size_t offset, half *p) void vstore_half_rtz (double data, size_t offset, half *p) void vstore_half_rtp (double data, size_t offset, half *p) void vstore_half_rtn (double data, size_t offset, half *p) | The double value given by data is first converted to a half value using the appropriate rounding mode. The half value is then written to address computed as $(p+$ offset). The address computed as ( $p+$ offset) must be16-bit aligned. <br> vstore_half use the default rounding mode. The default rounding mode is round to nearest even. |
| void vstore_half $\boldsymbol{n}$ (doublen data, size_t offset, half *p) void vstore_half $\boldsymbol{n}$ _rte (doublen data, size_t offset, half *p) void vstore_half $\boldsymbol{n}$ _rtz (doublen data, size_t offset, half *p) void vstore_halfn_rtp (doublen data, size_t offset, half *p) void vstore_halfn_rtn (doublen data, size_t offset, half *p) | The double $n$ value given by data is converted to a half $n$ value using the appropriate rounding mode. The half $n$ value is then written to address computed as $(p+(o f f s e t * n))$. The address computed as $(p+($ offset $* n))$ must be 16 -bit aligned. <br> vstore_half $\boldsymbol{n}$ uses the default rounding mode. The default rounding mode is round to nearest even. |
| float $n$ | For $\mathrm{n}=2,4,8$ and 16 read sizeof (halfn) |


| ```const half *p) floatn vloada_halfn (size_t offset, const constant half *p)``` | bytes of data from address ( $p+$ (offset * $n)$ ). The data read is interpreted as a half $n$ value. The half $n$ value read is converted to a float $n$ value and the float $n$ value is returned. <br> The address computed as $(p+($ offset $* n))$ must be aligned to sizeof (half $n$ ) bytes. <br> For $\mathrm{n}=3$, vloada_half3 reads a half3 from address $(p+($ offset $* 4))$ and returns a float3. The address computed as ( $p+$ $($ offset * 4)) must be aligned to sizeof (half) * 4 bytes. |
| :---: | :---: |
| void vstorea_halfn (float $n$ data, size_t offset, half *p) void vstorea_halfn_rte (floatn data, size_t offset, half *p) void vstorea_halfn_rtz (floatn data, size_t offset, half *p) void vstorea_halfn_rtp (floatn data, size_t offset, half *p) void vstorea_halfn_rtn (floatn data, size_t offset, half *p) | The float $n$ value given by data is converted to a half $n$ value using the appropriate rounding mode. <br> For $\mathrm{n}=2,4,8$ and 16 , the half $n$ value is written to the address computed as $(p+$ $($ offset $* n)$ ). The address computed as $(p+$ (offset $* n)$ ) must be aligned to sizeof (halfn) bytes. <br> For $\mathrm{n}=3$, the half 3 value is written to the address computed as $(p+($ offset $* 4))$. The address computed as ( $p+($ offset $* 4)$ ) must be aligned to sizeof (half) $* 4$ bytes. <br> vstorea_half $\boldsymbol{n}$ uses the default rounding mode. The default rounding mode is round to nearest even. |
| void vstorea_half $\boldsymbol{n}$ (double data, size_t offset, half *p) void vstorea_halfn_rte (doublen data, size_t offset, half *p) void vstorea_half $\boldsymbol{n} \_\boldsymbol{r t z}$ (doublen data, size_t offset, half *p) void vstorea_halfn_rtp (doublen data, size_t offset, half *p) void vstorea_halfn_rtn (doublen data, size_t offset, half *p) | The double $n$ value is converted to a half $n$ value using the appropriate rounding mode. <br> For $\mathrm{n}=2,4,8$ or 16 , the half $n$ value is written to the address computed as ( $p+$ $($ offset $* n)$ ). The address computed as $(p+$ $($ offset $* n)$ ) must be aligned to sizeof (halfn) bytes. <br> For $\mathrm{n}=3$, the half 3 value is written to the address computed as $(p+($ offset $* 4))$. The address computed as ( $p+($ offset $* 4)$ ) must be aligned to sizeof (half) $* 4$ bytes. |


|  | vstorea_half $\boldsymbol{n}$ uses the default rounding <br> mode. The default rounding mode is round <br> to nearest even. |
| :--- | :--- |

Table 6.15 Vector Data Load and Store Functions ${ }^{38}$
The results of vector data load and store functions are undefined if the address being read from or written to is not correctly aligned as described in table 6.15. The pointer argument p can be a pointer to global, local or private memory for store functions described in table 6.15. The pointer argument p can be a pointer to global, local, constant or private memory for load functions described in table 6.15.

NOTE: The vector data load and store functions variants that take pointer arguments which point to the generic address space are also supported.

[^27]In addition vloada_half3 reads $x, y, z$ components from address $(p+(o f f s e t * 4))$ into a 3-component vector and vstorea_half3 writes $x, y, z$ components from a 3-component vector to address $(p+($ offset $* 4)$ ).

### 6.13.8 Synchronization Functions

The OpenCL C programming language implements the following synchronization functions.

| Function | Description |
| :---: | :---: |
| ```void work_group_barrier }\mp@subsup{}{}{39}\mathrm{ ( cl_mem_fence_flags flags) void work_group_barrier ( cl_mem_fence_flags flags, memory_scope scope }\mp@subsup{}{}{40}\mathrm{ )``` | All work-items in a work-group executing the kernel on a processor must execute this function before any are allowed to continue execution beyond the work_group_barrier. This function must be encountered by all work-items in a work-group executing the kernel. These rules apply to NDranges implemented with uniform and non-uniform work-groups. <br> If work_group_barrier is inside a conditional statement, then all work-items must enter the conditional if any work-item enters the conditional statement and executes the work_group_barrier. <br> If work_group_barrier is inside a loop, all workitems must execute the work_group_barrier for each iteration of the loop before any are allowed to continue execution beyond the work_group_barrier. <br> The work_group_barrier function also supports a variant that specifies the memory scope. For the work_group_barrier variant that does not take a memory scope, the scope is memory_scope_work_group. <br> The scope argument specifies whether the memory accesses of work-items in the work-group to memory address space(s) identified by flags become visible to all work-items in the work-group, the device or all SVM devices. <br> The work_group_barrier function can also be used to specify which memory operations i.e. to global memory, local memory or images become visible to |

[^28]\(\left.\left.$$
\begin{array}{|l|l|}\hline & \begin{array}{l}\text { the appropriate memory scope identified by scope. } \\
\text { The flags argument specifies the memory address } \\
\text { spaces. This is a bitfield and can be set to 0 or a } \\
\text { combination of the following values ORed together. } \\
\text { When these flags are OR'ed together the } \\
\text { work_group_barrier acts as a combined barrier for } \\
\text { all address spaces specified by the flags ordering } \\
\text { memory accesses both within and across the } \\
\text { specified address spaces. }\end{array} \\
\\
\text { CLK_LOCAL_MEM_FENCE - The } \\
\text { work_group_barrier function will ensure that all } \\
\text { local memory accesses become visible to all work- } \\
\text { items in the work-group. Note that the value of } \\
\text { scope is ignored as the memory scope is always } \\
\text { memory_scope_work_group. }\end{array}
$$\right\} \begin{array}{l}CLK_GLOBAL_MEM_FENCE - The <br>
work_group_barrier function ensure that all global <br>
memory accesses become visible to the appropriate <br>

scope as given by scope.\end{array}\right\}\)| CLK_IMAGE_MEM_FENCE - The |
| :--- |
| work_group_barrier function will ensure that all |
| image memory accesses become visible to the |
| appropriate scope as given by scope. The value of |
| scope must be memory_scope_work_group. |
| The values of flags and scope must be the same for |
| all work-items in the work-group. |

## Table 6.16 Built-in Synchronization Functions

### 6.13.9 Address Space Qualifier Functions

The OpenCL C programming language implements the following address space qualifier functions. We use the generic type name gentype to indicate any of the built-in data types supported by OpenCL C or a user defined type.

| Function | Description |
| :--- | :--- |
| global gentype * <br> to_global (gentype *ptr) <br> const global gentype * <br> to_global (const gentype *ptr) | Returns a pointer that points to a region in the global <br> address space if to_global can cast ptr to the global <br> address space. Otherwise it returns NULL. |
| local gentype * <br> to_local (gentype *ptr) <br> const local gentype * <br> to_local (const gentype *ptr) | Returns a pointer that points to a region in the local <br> address space if to_local can cast ptr to the local <br> address space. Otherwise it returns NULL. |
| private gentype * <br> to_private (gentype *ptr) | Returns a pointer that points to a region in the private <br> address space if to_private can cast ptr to the private <br> address space. Otherwise it returns NULL. |
| const private gentype * <br> to_private (const gentype *ptr) | Returns a valid memory fence value for ptr. |
| cl_mem_fence_flags <br> get_fence (gentype *ptr) |  |
| cl_mem_fence_flags <br> get_fence (const gentype *ptr) |  |

Table 6.17 Built-in Address Space Qualifier Functions

### 6.13.10 Async Copies from Global to Local Memory, Local to Global Memory, and Prefetch

The OpenCL C programming language implements the following functions that provide asynchronous copies between global and local memory and a prefetch from global memory.

We use the generic type name gentype to indicate the built-in data types char, char $\left\{2\left|3^{41}\right| 4|8| 16\right\}$, uchar, uchar $\{2|3| 4|8| 16\}$, short, short $\{2|3| 4|8| 16\}$, ushort, $u s h o r t\{2|3| 4|8| 16\}$, int, int $\{2|3| 4|8| 16\}$, uint, $u i n t\{2|3| 4|8| 16\}$, long, long\{2|3|4|8|16\}, ulong, ulong\{2|3|4|8|16\}, float, float\{2|3|4|8|16\}, or double, double $\{2|3| 4|8| 16\}$ as the type for the arguments unless otherwise stated.

| Function | Description |
| :---: | :---: |
| ```event_t async_work_group_copy ( __local gentype *dst, const __global gentype *src, size_t num_gentypes, event_t event) event_t async_work_group_copy ( __global gentype *dst, const __local gentype *src, size_t num_gentypes, event_t event)``` | Perform an async copy of num_gentypes gentype elements from $s r c$ to $d s t$. The async copy is performed by all work-items in a work-group and this built-in function must therefore be encountered by all workitems in a work-group executing the kernel with the same argument values; otherwise the results are undefined. This rule applies to ND-ranges implemented with uniform and non-uniform work-groups. <br> Returns an event object that can be used by wait_group_events to wait for the async copy to finish. The event argument can also be used to associate the async_work_group_copy with a previous async copy allowing an event to be shared by multiple async copies; otherwise event should be zero. <br> 0 can be implicitly and explicitly casted to event_t type. <br> If event argument is non-zero, the event object supplied in event argument will be returned. |

[^29]|  | This function does not perform any implicit synchronization of source data such as using a barrier before performing the copy. |
| :---: | :---: |
|  | Perform an async gather of num_gentypes gentype elements from $s r c$ to $d s t$. The src_stride is the stride in elements for each gentype element read from $s r c$. The $d s t \_s t r i d e ~ i s ~ t h e ~ s t r i d e ~ i n ~ e l e m e n t s ~ f o r ~ e a c h ~$ gentype element written to $d s t$. The async gather is performed by all work-items in a work-group. This built-in function must therefore be encountered by all work-items in a work-group executing the kernel with the same argument values; otherwise the results are undefined. This rule applies to ND-ranges implemented with uniform and non-uniform work-groups <br> Returns an event object that can be used by wait_group_events to wait for the async copy to finish. The event argument can also be used to associate the async_work_group_strided_copy with a previous async copy allowing an event to be shared by multiple async copies; otherwise event should be zero. <br> 0 can be implicitly and explicitly casted to event_t type. <br> If event argument is non-zero, the event object supplied in event argument will be returned. <br> This function does not perform any implicit synchronization of source data such as using a barrier before performing the copy. <br> The behavior of async_work_group_strided_copy is undefined if src_stride or $d s t_{-}$stride is 0 , or if the src_stride or $d s t_{-}$stride values cause |

$\left.\left.\begin{array}{|l|l|}\hline & \begin{array}{l}\text { the src or dst pointers to exceed the upper } \\ \text { bounds of the address space during the } \\ \text { copy. }\end{array} \\ \hline \text { void wait_group_events (int num_events, } \\ \text { event_t *event_list) }\end{array} \quad \begin{array}{l}\text { Wait for events that identify the } \\ \text { async_work_group_copy operations to } \\ \text { complete. The event objects specified in } \\ \text { event_list will be released after the wait is } \\ \text { performed. }\end{array}\right\} \begin{array}{l}\text { This function must be encountered by all } \\ \text { work-items in a work-group executing the } \\ \text { kernel with the same num_events and event } \\ \text { objects specified in event_list; otherwise } \\ \text { the results are undefined. This rule applies } \\ \text { to ND-ranges implemented with uniform } \\ \text { and non-uniform work-groups }\end{array}\right\}$

## Table 6.18 Built-in Async Copy and Prefetch Functions

NOTE: The kernel must wait for the completion of all async copies using the wait_group_events built-in function before exiting; otherwise the behavior is undefined.

### 6.13.11 Atomic Functions

The OpenCL C programming language implements a subset of the C11 atomics (refer to section 7.17 of the C 11 specification) and synchronization operations. These operations play a special role in making assignments in one work-item visible to another. A synchronization operation on one or more memory locations is either an acquire operation, a release operation, or both an acquire and release operation ${ }^{42}$. A synchronization operation without an associated memory location is a fence and can be either an acquire fence, a release fence or both an acquire and release fence. In addition, there are relaxed atomic operations, which are not synchronization operations, and atomic read-modify-write operations which have special characteristics.

The types include
memory_order
which is an enumerated type whose enumerators identify memory ordering constraints;

```
memory_scope
```

which is an enumerated type whose enumerators identify scope of memory ordering constraints;
atomic_flag
which is a 32-bit integer type representing a lock-free, primitive atomic flag; and several atomic analogs of integer types.

In the following operation definitions:

* An A refers to one of the atomic types.
\# A C refers to its corresponding non-atomic type.
An M refers to the type of the other argument for arithmetic operations. For atomic integer types, M is C .
* The functions not ending in explicit have the same semantics as the corresponding explicit function with memory_order_seq_cst for the memory_order argument.
\# The functions that do not have memory_scope argument have the same semantics as the corresponding functions with the memory_scope argument set to memory_scope_device.

[^30]NOTE: With fine-grained system SVM, sharing happens at the granularity of individual loads and stores anywhere in host memory. Memory consistency is always guaranteed at synchronization points, but to obtain finer control over consistency, the OpenCL atomics functions may be used to ensure that the updates to individual data values made by one unit of execution are visible to other execution units. In particular, when a host thread needs fine control over the consistency of memory that is shared with one or more OpenCL devices, it must use atomic and fence operations that are compatible with the C 11 atomic operations ${ }^{43}$.

### 6.13.11.1 The ATOMIC_VAR_INIT macro

The ATOMIC_VAR_INIT macro expands to a token sequence suitable for initializing an atomic object of a type that is initialization-compatible with value. An atomic object with automatic storage duration that is not explicitly initialized using ATOMIC_VAR_INIT is initially in an indeterminate state; however, the default (zero) initialization for objects with static storage duration is guaranteed to produce a valid state.

```
#define ATOMIC_VAR_INIT(C value)
```

This macro can only be used to initialize atomic objects that are declared in program scope in the global address space.

Examples:

```
global atomic_int guide = ATOMIC_VAR_INIT(42);
```

Concurrent access to the variable being initialized, even via an atomic operation, constitutes a data-race.

### 6.13.11.2 The atomic_init function

The atomic_init function non-atomically initializes the atomic object pointed to by obj to the value value.

```
void atomic_init(volatile A *obj, C value)
```


## Examples:

```
local atomic_int local_guide;
if (get_local_id(0) == 0)
    atomic_init(&guide, 42);
```

[^31]```
work_group_barrier(CLK_LOCAL_MEM_FENCE);
```


### 6.13.11.3 Order and Consistency

The enumerated type memory_order specifies the detailed regular (non-atomic) memory synchronization operations as defined in section 5.1.2.4 of the C11 specification and may provide for operation ordering. Its enumeration constants are as follows:

```
memory_order_relaxed
memory_order_acquire
memory_order_release
memory_order_acq_rel
memory_order_seq_cst
```

The memory_order can be used when performing atomic operations to global or local memory.

### 6.13.11.4 Memory Scope

The enumerated type memory_scope specifies whether the memory ordering constraints given by memory_order apply to work-items in a work-group or work-items of a kernel(s) executing on the device or across devices (in the case of shared virtual memory). Its enumeration constants are as follows:

```
memory_scope_work_item44
memory_scope_work_group
memory scope device
memory_scope_all_svm_devices
```

The memory scope should only be used when performing atomic operations to global memory. Atomic operations to local memory only guarantee memory ordering in the work-group not across work-groups and therefore ignore the memory_scope value.

### 6.13.11.5 Fences

The following fence operations are supported.

### 6.13.11.5.1 The atomic_work_item_fence function

void atomic_work_item_fence(cl_mem_fence_flags flags,

[^32]```
memory_order order,
memory_scope scope)
```

flags must be set to CLK_GLOBAL_MEM_FENCE, CLK_LOCAL_MEM_FENCE, CLK_IMAGE_MEM_FENCE or a combination of these values ORed together; otherwise the behavior is undefined. The behavior of calling atomic_work_item_fence with CLK_IMAGE_MEM_FENCE ORed together with either CLK_GLOBAL_MEM_FENCE or CLK_LOCAL_MEM_FENCE is equivalent to calling atomic_work_item_fence individually for CLK_IMAGE_MEM_FENCE and the other flags. Passing both CLK_GLOBAL_MEM_FENCE and CLK_LOCAL_MEM_FENCE to atomic_work_item_fence will synchronize memory operations to both local and global memory through some shared atomic action, as described in section 3.3.6.2 of the OpenCL API specficiation.

Depending on the value of order, this operation:
has no effects, if order == memory_order_relaxed.

* is an acquire fence, if order == memory_order_acquire.
\# is a release fence, if order == memory_order_release.
* is both an acquire fence and a release fence, if order == memory_order_acq_rel.
* is a sequentially consistent acquire and release fence, if order == memory_order_seq_cst.

For images declared with the read_write qualifier, the atomic_work_item_fence must be called to make sure that writes to the image by a work-item become visible to that workitem on subsequent reads to that image by that work-item.

### 6.13.11.6 Atomic integer and floating-point types

The list of supported atomic type names are:

```
atomic_int
atomic_uint
atomic_long45
atomic_ulong
atomic_float
atomic_double 46
atomic_intptr_t47
```

[^33]```
atomic_uintptr_t
atomic_size_t
atomic_ptrdiff_t
```

Arguments to a kernel can be declared to be a pointer to the above atomic types or the atomic_flag type.

The representation of atomic integer, floating-point and pointer types have the same size as their corresponding regular types. The atomic_flag type must be implemented as a 32-bit integer.

### 6.13.11.7 Operations on atomic types

There are only a few kinds of operations on atomic types, though there are many instances of those kinds. This section specifies each general kind.

### 6.13.11.7.1 The atomic_store functions

```
void atomic_store(volatile A *object, C desired)
void atomic_store_explicit(volatile A *object,
    C desired,
    memory_order order)
void atomic_store_explicit(volatile A *object,
    C desired,
    memory_order order,
    memory_scope scope)
```

The order argument shall not be memory_order_acquire, nor memory_order_acq_rel. Atomically replace the value pointed to by object with the value of desired. Memory is affected according to the value of order.

### 6.13.11.7.2 The atomic_load functions

```
C atomic_load(volatile A *object)
```

C atomic_load_explicit(volatile A *object,
memory_order order)
C atomic_load_explicit(volatile A *object,

[^34]```
memory_order order,
memory_scope scope)
```

The order argument shall not be memory_order_release nor memory_order_acq_rel. Memory is affected according to the value of order. Atomically returns the value pointed to by object.

### 6.13.11.7.3 The atomic_exchange functions

```
C atomic_exchange(volatile A *object, C desired)
C atomic_exchange_explicit(volatile A *object,
    C desired,
    memory_order order)
C atomic_exchange_explicit(volatile A *object,
    C desired,
    memory_order order,
    memory_scope scope)
```

Atomically replace the value pointed to by object with desired. Memory is affected according to the value of order. These operations are read-modify-write operations (as defined by section 5.1.2.4 of the C11 specification). Atomically returns the value pointed to by object immediately before the effects.

### 6.13.11.7.4 The atomic_compare_exchange functions

```
bool atomic_compare_exchange_strong(volatile A *object,
    C *expected, C desired)
bool atomic_compare_exchange_strong_explicit(
    volatile A *object,
    C *expected,
    C desired,
    memory_order success,
    memory_order failure)
bool atomic_compare_exchange_strong_explicit(
    volatile A *object,
    C *expected,
    C desired,
    memory_order success,
    memory_order failure,
    memory_scope scope)
```

```
bool atomic_compare_exchange_weak(volatile A *object,
    C *expected, C desired)
bool atomic_compare_exchange_weak_explicit(
    volatile A `object,
    C *expected,
    C desired,
    memory_order success,
    memory_order failure)
bool atomic_compare_exchange_weak_explicit(
    volatile A `object,
    C *expected,
    C desired,
    memory_order success,
    memory_order failure,
    memory_scope scope)
```

The failure argument shall not be memory_order_release nor memory_order_acq_rel. The failure argument shall be no stronger than the success argument. Atomically, compares the value pointed to by object for equality with that in expected, and if true, replaces the value pointed to by object with desired, and if false, updates the value in expected with the value pointed to by object. Further, if the comparison is true, memory is affected according to the value of success, and if the comparison is false, memory is affected according to the value of failure. These operations are atomic read-modify-write operations (as defined by section 5.1.2.4 of the C11 specification).

NOTE: The effect of the compare-and-exchange operations is

```
if (memcmp(object, expected, sizeof(*object) == 0)
    memcpy(object, &desired, sizeof(*object));
else
    memcpy(expected, object, sizeof(*object));
```

The weak compare-and-exchange operations may fail spuriously ${ }^{48}$. That is, even when the contents of memory referred to by expected and object are equal, it may return zero and store back to expected the same memory contents that were originally there.

These generic functions return the result of the comparison.

### 6.13.11.7.5 The atomic_fetch and modify functions

[^35]The following operations perform arithmetic and bitwise computations. All of these operations are applicable to an object of any atomic integer type. The key, operator, and computation correspondence is given in table below:

| key | $\boldsymbol{o p}$ | computation |
| :---: | :---: | :--- |
| add | + | addition |
| sub | - | subtraction |
| or | $\perp$ | bitwise inclusive or |
| xor | $\wedge$ | bitwise exclusive or |
| and | $\&$ | bitwise and |
| $\min$ | $\min$ | compute min |
| $\max$ | $\max$ | compute max |

NOTE: For atomic_fetch and modify functions with key = add or sub on atomic types atomic_intptr_t and atomic_uintptr_t, Mis ptrdiff_t. For atomic_fetch and modify functions with key $=$ or, xor, and, min and max on atomic types atomic_intptr_t and atomic_uintptr_t, $M$ is intptr_t and uintptr_t.

```
C atomic_fetch_key(volatile A *object, M operand)
C atomic_fetch_key_explicit(volatile A *object,
    M operand,
    memory_order order)
C atomic_fetch_key_explicit(volatile A *object,
        M operand,
        memory_order order,
        memory_scope scope)
```

Atomically replaces the value pointed to by object with the result of the computation applied to the value pointed to by object and the given operand. Memory is affected according to the value of order. These operations are atomic read-modify-write operations (as defined by section 5.1.2.4 of the C11 specification). For signed integer types, arithmetic is defined to use two's complement representation with silent wrap-around on overflow; there are no undefined results. For address types, the result may be an undefined address, but the operations otherwise have no undefined behavior. Returns atomically, the value pointed to by object immediately before the effects.

### 6.13.11.7.6 Atomic flag type and operations

The atomic_flag type provides the classic test-and-set functionality. It has two states, set (value is non-zero) and clear (value is 0 ). Operations on an object of type atomic_flag shall be lock free.

The macro ATOMIC_FLAG_INIT may be used to initialize an atomic_flag to the clear
state. An atomic_flag that is not explicitly initialized with ATOMIC_FLAG_INIT is initially in an indeterminate state.

This macro can only be used for atomic objects that are declared in program scope in the global address space with the atomic_flag type.

Example:

```
global atomic_flag guard = ATOMIC_FLAG_INIT;
```


### 6.13.11.7.7 The atomic_flag_test_and_set functions

```
bool atomic_flag_test_and_set(
    volatile atomic_flag *object)
bool atomic_flag_test_and_set_explicit(
    vōlatīle ātomic_flag *object,
    memory_order order)
bool atomic_flag_test_and_set_explicit(
    volatile atomic_flag *object,
    memory_order order,
    memory_scope scope)
```

Atomically sets the value pointed to by object to true. Memory is affected according to the value of order. These operations are atomic read-modify-write operations (as defined by section 5.1.2.4 of the C11 specification). Returns atomically, the value of the object immediately before the effects.

### 6.13.11.7.8 The atomic_flag_clear functions

```
void atomic_flag_clear(volatile atomic_flag *object)
void atomic_flag_clear_explicit(
    volatile atomic_flag *object,
    memory_order order)
void atomic_flag_clear_explicit(
        volatile atomic_flag *object,
        memory_order order,
        memory_scope scope)
```

The order argument shall not be memory_order_acquire nor memory_order_acq_rel. Atomically sets the value pointed to by object to false. Memory is affected according to the value of order.

### 6.13.11.8 Restrictions

All operations on atomic types must be performed using the built-in atomic functions. C 11 and $\mathrm{C}++11$ support operators on atomic types. OpenCL C does not support operators with atomic types. Using atomic types with operators should result in a compilation error.

* The atomic_bool, atomic_char, atomic_uchar, atomic_short, atomic_ushort, atomic_intmax_t and atomic_uintmax_t types are not supported by OpenCL C.
* OpenCL C requires that the built-in atomic functions on atomic types are lock-free.
* The _Atomic type specifier and _Atomic type qualifier are not supported by OpenCL C.

The behavior of atomic operations where pointer arguments to the atomic functions refers to an atomic type in the private address space is undefined.

### 6.13.12 Miscellaneous Vector Functions

The OpenCL C programming language implements the following additional built-in vector functions. We use the generic type name gentypen (or gentypem) to indicate the built-in datatypes char $\{2|4| 8 \mid 16\}, \operatorname{uchar}\{2|4| 8 \mid 16\}, \operatorname{short}\{2|4| 8 \mid 16\}$, ushort $\{2|4| 8 \mid 16\}, \operatorname{half}\{2|4| 8 \mid 16\}^{49}$, int $\{2|4| 8 \mid 16\}$, uint $\{2|4| 8 \mid 16\}$, long $\{2|4| 8 \mid 16\}, \operatorname{long}\{2|4| 8 \mid 16\}$, float $\{2|4| 8 \mid 16\}$ or double $\{2|4| 8 \mid 16\}^{50}$ as the type for the arguments unless otherwise stated. We use the generic name ugentypen to indicate the built-in unsigned integer data types.

| Function | Description |
| :---: | :---: |
| int vec_step (gentypen $a$ ) <br> int vec_step (char3a) int vec_step (uchar3 $a$ ) int vec_step $(\operatorname{short} 3 a)$ int vec_step (ushort3 a) int vec_step (half3 $a$ ) int vec_step (int3a) int vec_step (uint3 $a$ ) int vec_step (long3 $a$ ) int vec_step (ulong3a) int vec_step (float3 $a$ ) int vec_step (double3 a) <br> int vec_step(type) | The vec_step built-in function takes a built-in scalar or vector data type argument and returns an integer value representing the number of elements in the scalar or vector. <br> For all scalar types, vec_step returns 1 . <br> The vec_step built-in functions that take a 3component vector return 4 . <br> vec_step may also take a pure type as an argument, e.g. vec_step(float2) |
| $\begin{gathered} \text { gentypen shuffle (gentype } m x, \\ \text { ugentypen mask) } \\ \text { gentypen shuffle2 (gentypem } x, \\ \text { gentype } m, \\ \text { ugentypen mask) } \end{gathered}$ | The shuffle and shuffle2 built-in functions construct a permutation of elements from one or two input vectors respectively that are of the same type, returning a vector with the same element type as the input and length that is the same as the shuffle mask. The size of each element in the mask must match the size of each element in the result. For shuffle, only the $\operatorname{ilogb}(2 m-1)$ least significant bits of each mask element are considered. For shuffle2, only the ilogb( $2 m-1)+1$ least significant bits of each mask element are considered. Other bits in the mask shall be ignored. <br> The elements of the input vectors are numbered from left to right across one or both of the vectors. |

[^36]

Table 6.20 Built-in Miscellaneous Vector Functions

### 6.13.13 printf

The OpenCL C programming language implements the printf function.

| Function | Description |
| :--- | :--- |
| int printf(constant char * restrict format, ...) | The printf built-in function writes output to an <br> implementation-defined stream such as stdout <br> under control of the string pointed to by format <br> that specifies how subsequent arguments are <br> converted for output. If there are insufficient <br> arguments for the format, the behavior is <br> undefined. If the format is exhausted while <br> arguments remain, the excess arguments are <br> evaluated (as always) but are otherwise <br> ignored. The printf function returns when the <br> end of the format string is encountered. |
| printf returns 0 if it was executed successfully |  |
| and -1 otherwise. |  |

Table 6.21 Built-in printf Function

### 6.13.13.1 printf output synchronization

When the event that is associated with a particular kernel invocation is completed, the output of all printf() calls executed by this kernel invocation is flushed to the implementation-defined output stream. Calling clFinish on a command queue flushes all pending output by printf in previously enqueued and completed commands to the implementation-defined output stream. In the case that printf is executed from multiple work-items concurrently, there is no guarantee of ordering with respect to written data. For example, it is valid for the output of a work-item with a global id $(0,0,1)$ to appear intermixed with the output of a work-item with a global id $(0,0,4)$ and so on.

### 6.13.13.2 printf format string

The format shall be a character sequence, beginning and ending in its initial shift state. The format is composed of zero or more directives: ordinary characters (not \%), which are copied unchanged to the output stream; and conversion specifications, each of which results in fetching zero or more subsequent arguments, converting them, if applicable, according to the corresponding conversion specifier, and then writing the result to the output stream. The format is in the constant address space and must be resolvable at compile time i.e. cannot be dynamically created by the executing program, itself.

Each conversion specification is introduced by the character \%. After the \%, the following appear in sequence:
\$ Zero or more flags (in any order) that modify the meaning of the conversion specification.

* An optional minimum field width. If the converted value has fewer characters than the field width, it is padded with spaces (by default) on the left (or right, if the left adjustment flag, described later, has been given) to the field width. The field width takes the form of a nonnegative decimal integer. ${ }^{51}$ )
* An optional precision that gives the minimum number of digits to appear for the $\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}$, and $\mathbf{X}$ conversions, the number of digits to appear after the decimalpoint character for $\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}$, and $\mathbf{F}$ conversions, the maximum number of significant digits for the $\mathbf{g}$ and $\mathbf{G}$ conversions, or the maximum number of bytes to be written for $\mathbf{s}$ conversions. The precision takes the form of a period (.) followed by an optional decimal integer; if only the period is specified, the precision is taken as zero. If a precision appears with any other conversion specifier, the behavior is undefined.
* An optional vector specifier.

4 A length modifier that specifies the size of the argument. The length modifier is required with a vector specifier and together specifies the vector type. Implicit conversions between vector types are disallowed (as per section 6.2.1). If the vector specifier is not specified, the length modifier is optional.

A conversion specifier character that specifies the type of conversion to be applied.
The flag characters and their meanings are:

- The result of the conversion is left-justified within the field. (It is right-justified if this flag is not specified.)
$+\quad$ The result of a signed conversion always begins with a plus or minus sign. (It begins with a sign only when a negative value is converted if this flag is not specified. ${ }^{52}$ )
space If the first character of a signed conversion is not a sign, or if a signed conversion results in no characters, a space is prefixed to the result. If the space and + flags both appear, the space flag is ignored.
\# The result is converted to an 'alternative form''. For $\boldsymbol{o}$ conversion, it increases the precision, if and only if necessary, to force the first digit of the result to be a zero (if the value and precision are both 0 , a single 0 is printed). For $\boldsymbol{x}($ or $\boldsymbol{X})$ conversion, a nonzero

[^37]result has $\boldsymbol{0 x}$ (or $\boldsymbol{0 X}$ ) prefixed to it. For $\boldsymbol{a}, \boldsymbol{A}, \boldsymbol{e}, \boldsymbol{E}, \boldsymbol{f}, \boldsymbol{F}, \boldsymbol{g}$, and $\boldsymbol{G}$ conversions, the result of converting a floating-point number always contains a decimal-point character, even if no digits follow it. (Normally, a decimal-point character appears in the result of these conversions only if a digit follows it.) For $\boldsymbol{g}$ and $\boldsymbol{G}$ conversions, trailing zeros are not removed from the result. For other conversions, the behavior is undefined.

0 For $\boldsymbol{d}, \boldsymbol{i}, \boldsymbol{o}, \boldsymbol{u}, \boldsymbol{x}, \boldsymbol{X}, \boldsymbol{a}, \boldsymbol{A}, \boldsymbol{e}, \boldsymbol{E}, \boldsymbol{f}, \boldsymbol{F}, \boldsymbol{g}$, and $\boldsymbol{G}$ conversions, leading zeros (following any indication of sign or base) are used to pad to the field width rather than performing space padding, except when converting an infinity or NaN. If the $\mathbf{0}$ and - flags both appear, the $\boldsymbol{0}$ flag is ignored. For $\boldsymbol{d}, \boldsymbol{i}, \boldsymbol{o}, \boldsymbol{u}, \boldsymbol{x}$, and $\boldsymbol{X}$ conversions, if a precision is specified, the $\boldsymbol{0}$ flag is ignored. For other conversions, the behavior is undefined.

The vector specifier and its meaning is:
$\mathbf{v} n \quad$ Specifies that a following $\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}, \mathbf{G}, \mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}$, or $\mathbf{X}$ conversion specifier applies to a vector argument, where $n$ is the size of the vector and must be $2,3,4,8$ or 16 .

The vector value is displayed in the following general form:
valuel C value $2 \mathrm{C} \ldots . \mathrm{C}$ valuen
where C is a separator character. The value for this separator character is a comma.
If the vector specifier is not used, the length modifiers and their meanings are:
hh Specifies that a following $\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}$, or $\mathbf{X}$ conversion specifier applies to a char or uchar argument (the argument will have been promoted according to the integer promotions, but its value shall be converted to char or uchar before printing).
h Specifies that a following $\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}$, or $\mathbf{X}$ conversion specifier applies to a short or ushort argument (the argument will have been promoted according to the integer promotions, but its value shall be converted to short or unsigned short before printing).
$\mathbf{l}$ (ell) Specifies that a following $\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}$, or $\mathbf{X}$ conversion specifier applies to a long or ulong argument. The $\mathbf{l}$ modifier is supported by the full profile. For the embedded profile, the $\mathbf{l}$ modifier is supported only if 64 -bit integers are supported by the device.

If the vector specifier is used, the length modifiers and their meanings are:
hh Specifies that a following $\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}$, or $\mathbf{X}$ conversion specifier applies to a charn or ucharn argument (the argument will not be promoted).
h Specifies that a following $\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}$, or $\mathbf{X}$ conversion specifier applies to a shortn or ushortn argument (the argument will not be promoted); that
a following $\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}$, or $\mathbf{G}$ conversion specifier applies to a halfn ${ }^{53}$ argument.
hl This modifier can only be used with the vector specifier. Specifies that a following $\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}$, or $\mathbf{X}$ conversion specifier applies to a int $\boldsymbol{n}$ or uint $\boldsymbol{n}$ argument; that a following $\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}$, or $\mathbf{G}$ conversion specifier applies to a floatn argument.
$\mathbf{l}$ (ell) Specifies that a following $\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}$, or $\mathbf{X}$ conversion specifier applies to a longn or ulongn argument; that a following $\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}$, or $\mathbf{G}$ conversion specifier applies to a doublen argument. The $\mathbf{I}$ modifier is supported by the full profile. For the embedded profile, the $\mathbf{I}$ modifier is supported only if 64-bit integers or double-precision floating-point are supported by the device.

If a vector specifier appears without a length modifier, the behavior is undefined. The vector data type described by the vector specifier and length modifier must match the data type of the argument; otherwise the behavior is undefined.

If a length modifier appears with any conversion specifier other than as specified above, the behavior is undefined.

The conversion specifiers and their meanings are:
d,i The int, charn, shortn, int $\boldsymbol{n}$ or longn argument is converted to signed decimal in the style [-]dddd. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it is expanded with leading zeros. The default precision is 1 . The result of converting a zero value with a precision of zero is no characters.
$\mathbf{o , u}$,
$\mathbf{x , X}$ The unsigned int, ucharn, ushort $\boldsymbol{n}$, uint $\boldsymbol{n}$ or ulongn argument is converted to unsigned octal ( $\mathbf{0}$ ), unsigned decimal ( $\mathbf{u}$ ), or unsigned hexadecimal notation ( $\mathbf{x}$ or $\mathbf{X}$ ) in the style $d d d d$; the letters abcdef are used for $\mathbf{x}$ conversion and the
letters ABCDEF for $\mathbf{X}$ conversion. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it is expanded with leading zeros. The default precision is 1 . The result of converting a zero value with a precision of zero is no characters.
f,F A double, half $\boldsymbol{n}$, float $\boldsymbol{n}$ or doublen argument representing a floating-point number is converted to decimal notation in the style [-]ddd.ddd, where the number of digits after the decimal-point character is equal to the precision specification. If the precision is missing, it is taken as 6 ; if the precision is zero and the \# flag is not specified, no decimal-point character appears. If a decimal-point character appears, at least one digit appears before it. The value is rounded to the appropriate number of digits. A double, half $\boldsymbol{n}$, float $\boldsymbol{n}$ or doublen argument representing an infinity is converted in one of the styles [-]inf or [-]infinity - which style is implementation-defined. A double, half $\boldsymbol{n}$,

[^38]floatn or doublen argument representing a NaN is converted in one of the styles []nan or [-]nan( $n$-char-sequence) - which style, and the meaning of any $n$-charsequence, is implementation-defined. The $\mathbf{F}$ conversion specifier produces INF, INFINITY, or NAN instead of inf, infinity, or nan, respectively. ${ }^{54}$ )
$\mathbf{e}, \mathbf{E}$ A double, halfn, floatn or doublen argument representing a floating-point number is converted in the style [ -$] d . d d d \mathrm{e} \pm d d$, where there is one digit (which is nonzero if the argument is nonzero) before the decimal-point character and the number of digits after it is equal to the precision; if the precision is missing, it is taken as 6 ; if the precision is zero and the \# flag is not specified, no decimal-point character appears. The value is rounded to the appropriate number of digits. The $\mathbf{E}$ conversion specifier produces a number with $\mathbf{E}$ instead of $\mathbf{e}$ introducing the exponent. The exponent always contains at least two digits, and only as many more digits as necessary to represent the exponent. If the value is zero, the exponent is zero. A double, half $\boldsymbol{n}$, float $\boldsymbol{n}$ or doublen argument representing an infinity or NaN is converted in the style of an $\mathbf{f}$ or $\mathbf{F}$ conversion specifier.
g,G A double, halfn, floatn or doublen argument representing a floating-point number is converted in style $\mathbf{f}$ or $\mathbf{e}$ (or in style $\mathbf{F}$ or $\mathbf{E}$ in the case of a $\mathbf{G}$ conversion specifier), depending on the value converted and the precision. Let $P$ equal the precision if nonzero, 6 if the precision is omitted, or 1 if the precision is zero. Then, if a conversion with style $\mathbf{E}$ would have an exponent of $X$ : - if $P>X \geq-4$, the conversion is with style $\mathbf{f}($ or $\mathbf{F})$ and precision $P-(X+1)$. - otherwise, the conversion is with style $\mathbf{e}$ (or $\mathbf{E}$ ) and precision $P-1$. Finally, unless the \# flag is used, any trailing zeros are removed from the fractional portion of the result and the decimal-point character is removed if there is no fractional portion remaining. A double, halfn, floatn or doublen $\mathbf{e}$ argument representing an infinity or NaN is converted in the style of an $\mathbf{f}$ or $\mathbf{F}$ conversion specifier.
a,A A double, halfn, floatn or doublen argument representing a floating-point number is converted in the style $[-] 0 \mathbf{x} h . h h h h \mathbf{p} \pm d$, where there is one hexadecimal digit (which is nonzero if the argument is a normalized floating-point number and is otherwise unspecified) before the decimal-point character ${ }^{55}$ ) and the number of hexadecimal digits after it is equal to the precision; if the precision is missing, then the precision is sufficient for an exact representation of the value; if the precision is zero and the \# flag is not specified, no decimal point character appears. The letters abcdef are used for a conversion and the letters ABCDEF for A conversion. The A conversion specifier produces a number with $\mathbf{X}$ and $\mathbf{P}$ instead of $\mathbf{x}$ and $\mathbf{p}$. The exponent always contains at least one digit, and only as many more digits as necessary to represent the decimal exponent of 2 . If the value is zero, the exponent is zero. A double, halfn, floatn or doublen argument representing an infinity or NaN is converted in the style of an $\mathbf{f}$ or $\mathbf{F}$ conversion specifier.

[^39]NOTE: The conversion specifiers $\mathbf{e}, \mathbf{E}, \mathbf{g}, \mathbf{G}, \mathbf{a}, \mathbf{A}$ convert a float or half argument that is a scalar type to a double only if the double data type is supported. If the double data type is not supported, the argument will be a float instead of a double and the half type will be converted to a float.
c The int argument is converted to an unsigned char, and the resulting character is written.
s The argument shall be a literal string. ${ }^{56}$ Characters from the literal string array are written up to (but not including) the terminating null character. If the precision is specified, no more than that many bytes are written. If the precision is not specified or is greater than the size of the array, the array shall contain a null character.
p The argument shall be a pointer to void. The pointer can refer to a memory region in the global, constant, local, private or generic address space. The value of the pointer is converted to a sequence of printing characters in an implementation-defined manner.
\% A \% character is written. No argument is converted. The complete conversion specification shall be \% \%.

If a conversion specification is invalid, the behavior is undefined. If any argument is not the correct type for the corresponding conversion specification, the behavior is undefined.

In no case does a nonexistent or small field width cause truncation of a field; if the result of a conversion is wider than the field width, the field is expanded to contain the conversion result.

For $\mathbf{a}$ and $\mathbf{A}$ conversions, the value is correctly rounded to a hexadecimal floating number with the given precision.

A few examples of printf are given below:

```
float4 f = (float4)(1.0f, 2.0f, 3.0f, 4.0f);
uchar4 uc = (uchar4) (0xFA, 0xFB, 0xFC, 0xFD);
printf("f4 = %2.2v4hlf\n", f);
printf("uc = %#v4hhx\n", uc);
```

The above two printf calls print the following:

```
f4 = 1.00,2.00,3.00,4.00
uc = 0xfa,0xfb,0xfc,0xfd
```

[^40]A few examples of valid use cases of printf for the conversion specifier $\mathbf{s}$ are given below. The argument value must be a pointer to a literal string.

```
kernel void my_kernel( ... )
{
    printf("%s\n", "this is a test string\n");
}
```

A few examples of invalid use cases of printf for the conversion specifier $\mathbf{s}$ are given below:

```
kernel void my_kernel(global char *s, ... )
{
    printf("%s\n", s);
    constant char *p = "this is a test string\n";
    printf("%s\n", p);
    printf("%s\n", &p[3]);
}
```

A few examples of invalid use cases of printf where data types given by the vector specifier and length modifier do not match the argument type are given below:

```
kernel void my_kernel(global char *s, ... )
{
    uint2 ui = (uint2)(0x12345678, 0x87654321);
    printf("unsigned short value = (%#v2hx)\n", ui)
    printf("unsigned char value = (%#v2hhx)\n", ui)
}
```


### 6.13.13.3 Differences between OpenCL C and C99 printf

* The $\mathbf{l}$ modifier followed by a c conversion specifier or $\mathbf{s}$ conversion specifier is not supported by OpenCL C.
* The $\mathbf{l l}, \mathbf{j}, \mathbf{z}, \mathbf{t}$, and $\mathbf{L}$ length modifiers are not supported by OpenCL $\mathbf{C}$ but are reserved.

The $\mathbf{n}$ conversion specifier is not supported by OpenCL C but is reserved.
\# OpenCL $\mathbf{C}$ adds the optional $\mathbf{v} n$ vector specifier to support printing of vector types.

* The conversion specifiers $\mathbf{f}, \mathbf{F}, \mathbf{e}, \mathbf{E}, \mathbf{g}, \mathbf{G}, \mathbf{a}, \mathbf{A}$ convert a float argument to a double only if the double data type is supported. Refer to the description of CL_DEVICE_DOUBLE_FP_CONFIG in table 4.3. If the double data type is not supported, the argument will be a float instead of a double.

For the embedded profile, the $\mathbf{l}$ length modifier is supported only if 64-bit integers are
supported.

* In OpenCL C, printf returns 0 if it was executed successfully and -1 otherwise vs. C99 where printf returns the number of characters printed or a negative value if an output or encoding error occurred.
* In OpenCL C, the conversion specifier $\mathbf{s}$ can only be used for arguments that are literal strings.


### 6.13.14 Image Read and Write Functions

The built-in functions defined in this section can only be used with image memory objects. An image memory object can be accessed by specific function calls that read from and/or write to specific locations in the image.

Image memory objects that are being read by a kernel should be declared with the read_only qualifier. write_image calls to image memory objects declared with the read_only qualifier will generate a compilation error. Image memory objects that are being written to by a kernel should be declared with the write_only qualifier. read_image calls to image memory objects declared with the write_only qualifier will generate a compilation error.
read_image and write_image calls to the same image memory object in a kernel are supported. Image memory objects that are being read and written by a kernel should be declared with the read_write qualifier.

The read_image calls returns a four component floating-point, integer or unsigned integer color value. The color values returned by read_image are identified as $x, y, z, w$ where $x$ refers to the red component, y refers to the green component, z refers to the blue component and w refers to the alpha component.

### 6.13.14.1 Samplers

The image read functions take a sampler argument. The sampler can be passed as an argument to the kernel using clSetKernelArg, or can be declared in the outermost scope of kernel functions, or it can be a constant variable of type sampler_t declared in the program source.

Sampler variables in a program are declared to be of type sampler_t. A variable of sampler_t type declared in the program source must be initialized with a 32-bit unsigned integer constant, which is interpreted as a bit-field specifiying the following properties:

- Addressing Mode
* Filter Mode

4 Normalized Coordinates
These properties control how elements of an image object are read by read_image $\{\mathbf{f}|\mathbf{i}| \mathbf{u i}\}$.
Samplers can also be declared as global constants in the program source using the following syntax.
const sampler_t
or
constant sampler_t <sampler name> = <value>
or

```
__constant sampler_t <sampler_name> = <value>
```

Note that samplers declared using the constant qualifier are not counted towards the maximum number of arguments pointing to the constant address space or the maximum size of the constant address space allowed per device (i.e. CL_DEVICE_MAX_CONSTANT_ARGS and CL_DEVICE_MAX_CONSTANT_BUFFER_SIZE as described in table 4.3).

The sampler fields are described in table 6.22.
$\left.\begin{array}{|l|l|}\hline \text { Sampler State } & \text { Description } \\ \hline \text { <normalized coords> } & \begin{array}{l}\text { Specifies whether the } x, y \text { and } z \text { coordinates are passed in } \\ \text { as normalized or unnormalized values. This must be a } \\ \text { literal value and can be one of the following predefined } \\ \text { enums: } \\ \text { CLK_NORMALIZED_COORDS_TRUE or }\end{array} \\ \text { CLK_NORMALIZED_COORDS_FALSE. } \\ \text { The samplers used with an image in multiple calls to } \\ \text { read_image\{fi|iui\} declared in a kernel must use the } \\ \text { same value for <normalized coords>. }\end{array}\right\}$

|  | CLK_ADDRESS_CLAMP <br> coordinates will return a border color. |
| :--- | :--- |
| CLK_ADDRESS_NONE - for this addressing mode the |  |
| programmer guarantees that the image coordinates used |  |
| to sample elements of the image refer to a location inside |  |
| the image; otherwise the results are undefined. |  |
| For 1D and 2D image arrays, the addressing mode |  |
| applies only to the $x$ and $(x, y)$ coordinates. The |  |
| addressing mode for the coordinate which specifies the |  |
| array index is always |  |
| CLK_ADDRESS_CLAMP_TO_EDGE. |  |

## Table 6.22 Sampler Descriptor

## Examples:

```
const sampler_t samplerA = CLK_NORMALIZED_COORDS_TRUE |
    CLK_ADDRESS_REPEAT
    CLK_FILTER_NEAREST;
```

samplerA specifies a sampler that uses normalized coordinates, the repeat addressing mode and a nearest filter.

The maximum number of samplers that can be declared in a kernel can be queried using the CL_DEVICE_MAX_SAMPLERS token in clGetDeviceInfo.

### 6.13.14.1.1 Determining the border color or value

If <addressing mode> in sampler is CLK_ADDRESS_CLAMP, then out-of-range image coordinates return the border color. The border color selected depends on the image channel order and can be one of the following values:

[^41]4 If the image channel order is CL_A, CL_INTENSITY, CL_Rx, CL_RA, CL_RGx, CL_RGBx, CL_sRGBx, CL_ARGB, CL_BGRA, CL_ABGR, CL_RGBA, CL_sRGBA or CL_sBGRA, the border color is $(0.0 f, 0.0 f, 0.0 f, 0.0 f)$.

4 If the image channel order is CL_R, CL_RG, CL_RGB, or CL_LUMINANCE, the border color is ( $0.0 \mathrm{f}, ~ 0.0 \mathrm{f}, ~ 0.0 f, 1.0 f$ ).

If the image channel order is CL_DEPTH, the border value is $0.0 f$.

### 6.13.14.1.2 sRGB Images

The built-in image read functions will perform sRGB to linear RGB conversions if the image is an sRGB image. Writing to sRGB images from a kernel is an optional extension. The cl_khr_srgb_image_writes extension will be reported in the CL_DEVICE_EXTENSIONS string if a device supports writing to sRGB images using write_imagef.
clGetSupportedImageFormats will return the supported sRGB images if CL_MEM_READ_WRITE or CL_MEM_WRITE_ONLY is specified in flags argument and the device supports writing to an sRGB image. If cl_khr_srgb_image_writes is supported, the built-in image write functions will perform the linear to sRGB conversion.

Only the R, G and B components are converted from linear to sRGB and vice-versa. The alpha component is returned as is.

### 6.13.14.2 Built-in Image Read Functions

The following built-in function calls to read images with a sampler ${ }^{58}$ are supported.

| Function | Description |
| :---: | :---: |
| ```float4 read_imagef ( read_only image2d_t image, sampler_t sampler, int2 coord) float4 read_imagef ( read_only image2d_t image, sampler_t sampler, float2 coord)``` | Use the coordinate (coord. $x$, coord.y) to do an element lookup in the 2D image object specified by image. <br> read_imagef returns floating-point values in the range [0.0 ... 1.0] for image objects created with image_channel_data_type set to one of the predefined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16. <br> read_imagef returns floating-point values in the range $[-1.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to CL_SNORM_INT8, or CL_SNORM_INT16. <br> read_imagef returns floating-point values for image objects created with image_channel_data_type set to CL_HALF_FLOAT or CL_FLOAT. <br> The read_imagef calls that take integer coordinates must use a sampler with filter mode set to CLK_FILTER_NEAREST, normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined. <br> Values returned by read_imagef for image objects with image_channel_data_type values not specified in the description above are undefined. |
| int4 read_imagei ( read_only image2d_t image, sampler_t sampler, | Use the coordinate (coord. $x$, coord. $y$ ) to do an element lookup in the 2D image object specified by image. |

[^42]$\left.\begin{array}{|l|l|}\hline \text { int2 coord) } \\ \text { int4 read_imagei ( } \\ \text { read_only image2d_t image, } \\ \text { sampler_t sampler, } \\ \text { float2 coord) }\end{array} \quad \begin{array}{l}\text { read_imagei and read_imageui return } \\ \text { unnormalized signed integer and unsigned integer } \\ \text { values respectively. Each channel will be stored in }\end{array}\right\}$ a 32-bit integer.


| sampler_t sampler, <br> float 4 coord) | are undefined. <br> read_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imageui are undefined. <br> The read_image $\{\mathbf{i} \mid \mathbf{u i}\}$ calls support a nearest filter only. The filter_mode specified in sampler must be set to CLK_FILTER_NEAREST; otherwise the values returned are undefined. <br> Furthermore, the read_image $\{\mathbf{i} \mid \mathbf{u i}\}$ calls that take integer coordinates must use a sampler with normalized coordinates set to <br> CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined. |
| :---: | :---: |
| float4 read_imagef ( <br> read_only image2d_array_t image, sampler_t sampler, int4 coord) <br> float4 read_imagef ( read_only image2d_array_t image, sampler_t sampler, float 4 coord) | Use coord.xy to do an element lookup in the 2D image identified by coord.z in the 2D image array specified by image. <br> read_imagef returns floating-point values in the range $[0.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to one of the predefined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16. <br> read_imagef returns floating-point values in the range $[-1.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to CL_SNORM_INT8, or CL_SNORM_INT16. <br> read_imagef returns floating-point values for image objects created with image_channel_data_type set to CL_HALF_FLOAT or CL_FLOAT. |


|  | The read_imagef calls that take integer coordinates must use a sampler with filter mode set to CLK_FILTER_NEAREST, normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined. <br> Values returned by read_imagef for image objects with image_channel_data_type values not specified in the description above are undefined. |
| :---: | :---: |
| ```int4 read_imagei ( read_only image2d_array_t image, sampler_t sampler, int4 coord)``` | Use coord.xy to do an element lookup in the 2D image identified by coord. $z$ in the 2D image array specified by image. <br> read_imagei and read_imageui return |
| int4 read_imagei ( read_only image2d_array_t image, sampler_t sampler, float4 coord) | unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer. <br> read_imagei can only be used with image objects created with image_channel_data_type set to one of the following values: |
| uint4 read_imageui ( read_only image2d_array_t image, sampler_t sampler, int4 coord) <br> uint4 read_imageui ( | CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imagei are undefined. |
| sampler_t sampler, <br> float 4 coord) | read_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imageui are undefined. |
|  | The read_image $\{\mathbf{i} \mid \mathbf{u i}\}$ calls support a nearest filter only. The filter_mode specified in sampler must be set to CLK_FILTER_NEAREST; otherwise |


|  | the values returned are undefined. <br> Furthermore, the read_image $\{\mathbf{i} \mid \mathbf{u i}\}$ calls that take integer coordinates must use a sampler with normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined. |
| :---: | :---: |
| ```float4 read_imagef ( read_only image1d_t image, sampler_t sampler, int coord) float4 read_imagef ( read_only image1d_t image, sampler_t sampler, float coord)``` | Use coord to do an element lookup in the 1D image object specified by image. <br> read_imagef returns floating-point values in the range [ $0.0 \ldots 1.0$ ] for image objects created with image_channel_data_type set to one of the predefined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16. <br> read_imagef returns floating-point values in the range $[-1.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to CL_SNORM_INT8, or CL_SNORM_INT16. <br> read_imagef returns floating-point values for image objects created with image_channel_data_type set to CL_HALF_FLOAT or CL_FLOAT. <br> The read_imagef calls that take integer coordinates must use a sampler with filter mode set to CLK_FILTER_NEAREST, normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined. <br> Values returned by read_imagef for image objects with image_channel_data_type values not specified in the description above are undefined. |
| int4 read_imagei | Use coord to do an element lookup in the 1D image |



| ```float4 read_imagef ( read_only image1d_array_t image, sampler_t sampler, float2 coord)``` | read_imagef returns floating-point values in the range $[0.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to one of the predefined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16. <br> read_imagef returns floating-point values in the range $[-1.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to CL_SNORM_INT8, or CL_SNORM_INT16. <br> read_imagef returns floating-point values for image objects created with image_channel_data_type set to CL_HALF_FLOAT or CL_FLOAT. <br> The read_imagef calls that take integer coordinates must use a sampler with filter mode set to CLK_FILTER_NEAREST, normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to <br> CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined. <br> Values returned by read_imagef for image objects with image_channel_data_type values not specified in the description above are undefined. |
| :---: | :---: |
| ```int4 read_imagei ( read_only image1d_array_t image, sampler_t sampler, int2 coord)``` | Use coord. $x$ to do an element lookup in the 1D image identified by coord.y in the 1D image array specified by image. <br> read_imagei and read_imageui return |
| int4 read_imagei ( read_only image1d_array_t image, sampler_t sampler, float 2 coord) | unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer. <br> read_imagei can only be used with image objects created with image_channel_data_type set to one of the following values: |
| uint4 read_imageui ( read_only image1d_array_t image, sampler_t sampler, int 2 coord) | CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imagei |


| uint4 read_imageui ( <br> read_only image1d_array_t image, <br> sampler_t sampler, <br> float 2 coord) | are undefined. <br> read_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imageui are undefined. <br> The read_image $\{\mathbf{i} \mid \mathbf{u i}\}$ calls support a nearest filter only. The filter_mode specified in sampler must be set to CLK_FILTER_NEAREST; otherwise the values returned are undefined. <br> Furthermore, the read_image $\{\mathbf{i} \mid \mathbf{u i}\}$ calls that take integer coordinates must use a sampler with normalized coordinates set to <br> CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined. |
| :---: | :---: |
| float read_imagef ( <br> read_only image2d_depth_t image, <br> sampler_t sampler, <br> int 2 coord) <br> float read_imagef ( read_only image2d_depth_t image, sampler_t sampler, float 2 coord) | Use the coordinate (coord.x, coord.y) to do an element lookup in the 2D depth image object specified by image. <br> read_imagef returns a floating-point value in the range [ $0.0 \ldots 1.0$ ] for depth image objects created with image_channel_data_type set to CL_UNORM_INT16 or CL_UNORM_INT24. <br> read_imagef returns a floating-point value for depth image objects created with image_channel_data_type set to CL_FLOAT. <br> The read_imagef calls that take integer coordinates must use a sampler with filter mode set to CLK_FILTER_NEAREST, normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, |


|  | CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined. <br> Values returned by read_imagef for depth image objects with image_channel_data_type values not specified in the description above are undefined. |
| :---: | :---: |
| ```float read_imagef ( read_only image2d_array_depth_t image, sampler_t sampler, int4 coord) float read_imagef ( read_only image2d_array_depth_t image, sampler_t sampler, float4 coord)``` | Use coord.xy to do an element lookup in the 2D image identified by coord. $z$ in the 2D depth image array specified by image. <br> read_imagef returns a floating-point value in the range $[0.0 \ldots 1.0]$ for depth image objects created with image_channel_data_type set to CL_UNORM_INT16 or CL_UNORM_INT24. <br> read_imagef returns a floating-point value for depth image objects created with image_channel_data_type set to CL_FLOAT. <br> The read_imagef calls that take integer coordinates must use a sampler with filter mode set to CLK_FILTER_NEAREST, normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode set to CLK_ADDRESS_CLAMP_TO_EDGE, CLK_ADDRESS_CLAMP or CLK_ADDRESS_NONE; otherwise the values returned are undefined. <br> Values returned by read_imagef for image objects with image_channel_data_type values not specified in the description above are undefined. |

## Table 6.23 Built-in Image Read Functions

### 6.13.14.3 Built-in Image Sampler-less Read Functions

The sampler-less read image functions behave exactly as the corresponding read image functions described in section 6.13.14.2 that take integer coordinates and a sampler with filter mode set to CLK_FILTER_NEAREST, normalized coordinates set to CLK_NORMALIZED_COORDS_FALSE and addressing mode to CLK_ADDRESS_NONE.
aQual in Table 6.24 refers to one of the access qualifiers. For samplerless read functions this may be read_only or read_write.

| Function | Description |
| :--- | :--- |
| float4 read_imagef ( <br> aQual image2d_t image, <br> int2 coord) | Use the coordinate (coord.x, coord.y) to do an <br> element lookup in the 2D image object specified by <br> image. |
|  | read_imagef returns floating-point values in the <br> range [0.0... 1.0] for image objects created with <br> image_channel_data_type set to one of the pre- <br> defined packed formats or CL_UNORM_INT8, or <br> CL_UNORM_INT16. |
|  | read_imagef returns floating-point values in the |
|  | range [-1.0 ... 1.0] for image objects created with <br> image_channel_data_type set to CL_SNORM_INT8, <br> or CL_SNORM_INT16. |
|  | read_imagef returns floating-point values for image <br> objects created with image_channel_data_type set to <br> CL_HALF_FLOAT or CL_FLOAT. |
|  | Values returned by read_imagef for image objects <br> with image_channel_data_type values not specified |
| in the description above are undefined. |  |


|  | read_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imageui are undefined. |
| :---: | :---: |
| float4 read_imagef ( aQual image3d_t image, int4 coord) | Use the coordinate (coord.x, coord.y, coord.z) to do an element lookup in the 3D image object specified by image. coord.w is ignored. <br> read_imagef returns floating-point values in the range $[0.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to one of the predefined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16. <br> read_imagef returns floating-point values in the range $[-1.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to CL_SNORM_INT8, or CL_SNORM_INT16. <br> read_imagef returns floating-point values for image objects created with image_channel_data_type set to CL_HALF_FLOAT or CL_FLOAT. <br> Values returned by read_imagef for image objects with image_channel_data_type values not specified in the description are undefined. |
| ```int4 read_imagei ( aQual image3d_t image, int4 coord) uint4 read_imageui ( aQual image3d_t image, int4 coord)``` | Use the coordinate (coord.x, coord.y, coord.z) to do an element lookup in the 3D image object specified by image. coord. $w$ is ignored. <br> read_imagei and read_imageui return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer. <br> read_imagei can only be used with image objects created with image_channel_data_type set to one of the following values: |


|  | CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imagei are undefined. <br> read_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imageui are undefined. |
| :---: | :---: |
| float4 read_imagef ( aQual image2d_array_t image, int4 coord) | Use coord.xy to do an element lookup in the 2D image identified by coord. $z$ in the 2D image array specified by image. <br> read_imagef returns floating-point values in the range $[0.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to one of the predefined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16. <br> read_imagef returns floating-point values in the range $[-1.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to CL_SNORM_INT8, or CL_SNORM_INT16. <br> read_imagef returns floating-point values for image objects created with image_channel_data_type set to CL_HALF_FLOAT or CL_FLOAT. <br> Values returned by read_imagef for image objects with image_channel_data_type values not specified in the description above are undefined. |
| int4 read_imagei ( aQual image2d_array_t image, int4 coord) <br> uint4 read_imageui ( | Use coord.xy to do an element lookup in the 2D image identified by coord.z in the 2D image array specified by image. <br> read_imagei and read_imageui return unnormalized signed integer and unsigned integer |


| aQual image2d_array_t image, <br> int4 coord) | values respectively. Each channel will be stored in a <br> 32-bit integer. <br> read_imagei can only be used with image objects <br> created with image_channel_data_type set to one of <br> the following values: <br> CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. |
| :--- | :--- |
|  | If the image_channel_data_type is not one of the <br> above values, the values returned by read_imagei <br> are undefined. <br> read_imageui can only be used with image objects |
|  | created with image_channel_data_type set to one of <br> the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. |
| If the image_channel_data_type is not one of the |  |
| above values, the values returned by read_imageui |  |
| are undefined. |  |


| aQual image1d_t image, int coord) | or 1D image buffer object specified by image. |
| :---: | :---: |
| uint4 read_imageui ( aQual image1d_t image, int coord) | read_imagei and read_imageui return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer. |
| ```int4 read_imagei ( aQual imageld_buffer_t image, int coord)``` | read_imagei can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. |
| uint4 read_imageui ( <br> aQual image1d_buffer_t image, int coord) | If the image_channel_data_type is not one of the above values, the values returned by read_imagei are undefined. |
|  | read_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imageui are undefined. |
| float4 read_imagef ( aQual image1d_array_t image, int2 coord) | Use coord. $x$ to do an element lookup in the 1D image identified by coord.y in the 1D image array specified by image. |
|  | read_imagef returns floating-point values in the range [ $0.0 \ldots 1.0$ ] for image objects created with image_channel_data_type set to one of the predefined packed formats or CL_UNORM_INT8, or CL_UNORM_INT16. |
|  | read_imagef returns floating-point values in the range $[-1.0 \ldots 1.0]$ for image objects created with image_channel_data_type set to CL_SNORM_INT8, or CL_SNORM_INT16. |
|  | read_imagef returns floating-point values for image objects created with image_channel_data_type set to CL_HALF_FLOAT or CL_FLOAT. |


|  | Values returned by read_imagef for image objects with image_channel_data_type values not specified in the description above are undefined. |
| :---: | :---: |
| ```int4 read_imagei ( aQual image1d_array_t image, int2 coord) uint4 read_imageui ( aQual imageld_array_t image, int2 coord)``` | Use coord. $x$ to do an element lookup in the 1D image identified by coord.y in the 1D image array specified by image. <br> read_imagei and read_imageui return unnormalized signed integer and unsigned integer values respectively. Each channel will be stored in a 32-bit integer. <br> read_imagei can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imagei are undefined. <br> read_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> If the image_channel_data_type is not one of the above values, the values returned by read_imageui are undefined. |
| float read_imagef ( aQual image2d_depth_t image, int 2 coord) | Use the coordinate (coord. $x$, coord. $y$ ) to do an element lookup in the 2D depth image object specified by image. <br> read_imagef returns a floating-point value in the range [ $0.0 \ldots 1.0$ ] for depth image objects created with image_channel_data_type set to CL_UNORM_INT16 or CL_UNORM_INT24. <br> read_imagef returns a floating-point value for depth image objects created with image_channel_data_type set to CL_FLOAT. |

\(\left.\left.$$
\begin{array}{|l|l|}\hline & \begin{array}{l}\text { Values returned by read_imagef for image objects } \\
\text { with image_channel_data_type values not specified } \\
\text { in the description above are undefined. }\end{array} \\
\hline \begin{array}{l}\text { float read_imagef ( } \\
\text { aQual image2d_array_depth_t image, } \\
\text { int } \text { coord })\end{array} & \begin{array}{l}\text { Use coord.xy to do an element lookup in the 2D } \\
\text { image identified by coord.z in the 2D depth image } \\
\text { array specified by image. }\end{array} \\
\text { read_imagef returns a floating-point value in the } \\
\text { range }[0.0 \ldots .1 .0] \text { for depth image objects created } \\
\text { with image_channel_data_type set to } \\
\text { CL_UNORM_INT16 or CL_UNORM_INT24. }\end{array}
$$\right\} \begin{array}{l}read_imagef returns a floating-point value for depth <br>
image objects created with <br>

image_channel_data_type set to CL_FLOAT.\end{array}\right\}\)| Values returned by read_imagef for image objects |
| :--- |
| with image_channel_data_type values not specified |
| in the description above are undefined. |

Table 6.24 Built-in Image Sampler-less Read Functions

### 6.13.14.4 Built-in Image Write Functions

The following built-in function calls to write images are supported. Note that image writes to sRGB images are only supported if the cl_khr_srgb_image_writes extension is supported; otherwise the behavior of writing to a sRGB image is undefined.
aQual in Table 6.25 refers to one of the access qualifiers. For write functions this may be write_only or read_write.

| Function | Description |
| :--- | :--- |
| void write_imagef ( | Write color value to location specified by coord.xy |
| aQual image2d_t image, | in the 2D image object specified by image. |
| int2 coord, | Appropriate data format conversion to the specified |
| float 4 color) | image format is done before writing the color value. <br> coord.x and coord.y are considered to be <br> void write_imagei ( <br> aQual image2d_t image, <br> int2 coord, <br> int 4 color $)$ |
| unnormalized coordinates and must be in the range 0 |  |


| void write_imageui ( aQual image2d_t image, int 2 coord, uint 4 color) | created with image_channel_data_type set to one of the pre-defined packed formats or set to CL_SNORM_INT8, CL_UNORM_INT8, CL_SNORM_INT16, CL_UNORM_INT16, CL_HALF_FLOAT or CL_FLOAT. Appropriate data format conversion will be done to convert channel data from a floating-point value to actual data format in which the channels are stored. <br> write_imagei can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. <br> write_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> The behavior of write_imagef, write_imagei and write_imageui for image objects created with image_channel_data_type values not specified in the description above or with $(x, y)$ coordinate values that are not in the range ( $0 \ldots$ image width $1,0 \ldots$ image height -1 ), respectively, is undefined. |
| :---: | :---: |
| ```void write_imagef ( aQual image2d_array_t image, int 4 coord, float 4 color) void write_imagei ( aQual image2d_array_t image, int 4 coord, int4 color) void write_imageui ( aQual image2d_array_t image, int 4 coord, uint4 color)``` | Write color value to location specified by coord.xy in the 2D image identified by coord.z in the 2D image array specified by image. Appropriate data format conversion to the specified image format is done before writing the color value. coord. $x$, coord. $y$ and coord. $z$ are considered to be unnormalized coordinates and must be in the range 0 $\ldots$ image width $-1,0 \ldots$ image height -1 and $0 \ldots$ image number of layers - 1 . <br> write_imagef can only be used with image objects created with image_channel_data_type set to one of the pre-defined packed formats or set to CL_SNORM_INT8, CL_UNORM_INT8, CL_SNORM_INT16, CL_UNORM_INT16, CL_HALF_FLOAT or CL_FLOAT. Appropriate data |


|  | format conversion will be done to convert channel data from a floating-point value to actual data format in which the channels are stored. <br> write_imagei can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. <br> write_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> The behavior of write_imagef, write_imagei and write_imageui for image objects created with image_channel_data_type values not specified in the description above or with $(x, y, z)$ coordinate values that are not in the range ( $0 \ldots$ image width $1,0 \ldots$ image height $-1,0 \ldots$ image number of layers -1 ), respectively, is undefined. |
| :---: | :---: |
| ```void write_imagef ( aQual image1d_t image, int coord, float4 color) void write_imagei ( aQual image1d_t image, int coord, int 4 color) void write_imageui ( aQual image1d_t image, int coord, uint4 color) void write_imagef ( aQual image1d_buffer_t image, int coord, float4 color)``` | Write color value to location specified by coord in the 1D image or 1D image buffer object specified by image. Appropriate data format conversion to the specified image format is done before writing the color value. coord is considered to be unnormalized coordinates and must be in the range $0 \ldots$ image width -1 . <br> write_imagef can only be used with image objects created with image_channel_data_type set to one of the pre-defined packed formats or set to CL_SNORM_INT8, CL_UNORM_INT8, CL_SNORM_INT16, CL_UNORM_INT16, CL_HALF_FLOAT or CL_FLOAT. Appropriate data format conversion will be done to convert channel data from a floating-point value to actual data format in which the channels are stored. <br> write_imagei can only be used with image objects created with image_channel_data_type set to one of |


| void write_imagei ( <br> aQual image1d_buffer_t image, int coord, int 4 color) <br> void write_imageui ( aQual image1d_buffer_t image, int coord, uint4 color) | the following values: <br> CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. <br> write_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> The behavior of write_imagef, write_imagei and write_imageui for image objects created with image_channel_data_type values not specified in the description above or with coordinate values that is not in the range ( $0 \ldots$ image width -1 ), is undefined. |
| :---: | :---: |
| ```void write_imagef ( aQual image1d_array_t image, int2 coord, float4 color) void write_imagei ( aQual image1d_array_t image, int2 coord, int4 color) void write_imageui ( aQual image1d_array_t image, int2 coord, uint4 color)``` | Write color value to location specified by coord. $x$ in the 1D image identified by coord.y in the 1D image array specified by image. Appropriate data format conversion to the specified image format is done before writing the color value. coord. $x$ and coord.y are considered to be unnormalized coordinates and must be in the range $0 \ldots$ image width -1 and $0 \ldots$ image number of layers - 1 . <br> write_imagef can only be used with image objects created with image_channel_data_type set to one of the pre-defined packed formats or set to CL_SNORM_INT8, CL_UNORM_INT8, CL_SNORM_INT16, CL_UNORM_INT16, CL_HALF_FLOAT or CL_FLOAT. Appropriate data format conversion will be done to convert channel data from a floating-point value to actual data format in which the channels are stored. <br> write_imagei can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. <br> write_imageui can only be used with image objects |


|  | created with image_channel_data_type set to one of the following values: <br> CL_UNSIGNED_INT8, <br> CL_UNSIGNED_INT16 and <br> CL_UNSIGNED_INT32. <br> The behavior of write_imagef, write_imagei and write_imageui for image objects created with image_channel_data_type values not specified in the description above or with $(x, y)$ coordinate values that are not in the range ( $0 \ldots$ image width $1,0 \ldots$ image number of layers -1 ), respectively, is undefined. |
| :---: | :---: |
| void write_imagef ( <br> aQual image2d_depth_t image, int 2 coord, float depth) | Write depth value to location specified by coord.xy in the 2D depth image object specified by image. Appropriate data format conversion to the specified image format is done before writing the depth value. coord. $x$ and coord.y are considered to be unnormalized coordinates and must be in the range 0 $\ldots$ image width -1 , and $0 \ldots$ image height -1 . <br> write_imagef can only be used with image objects created with image_channel_data_type set to CL_UNORM_INT16, CL_UNORM_INT24 or CL_FLOAT. Appropriate data format conversion will be done to convert depth valye from a floatingpoint value to actual data format associated with the image. <br> The behavior of write_imagef, write_imagei and write_imageui for image objects created with image_channel_data_type values not specified in the description above or with $(x, y)$ coordinate values that are not in the range ( $0 \ldots$ image width $1,0 \ldots$ image height -1 ), respectively, is undefined. |
| ```void write_imagef ( aQual image2d_array_depth_t image, int4 coord, float depth)``` | Write depth value to location specified by coord.xy in the 2D image identified by coord.z in the 2D depth image array specified by image. Appropriate data format conversion to the specified image format is done before writing the depth value. coord.x, coord. $y$ and coord.z are considered to be unnormalized coordinates and must be in the range 0 $\ldots$ image width $-1,0 \ldots$ image height -1 and $0 \ldots$ image number of layers - 1 . |


|  | write_imagef can only be used with image objects created with image_channel_data_type set to CL_UNORM_INT16, CL_UNORM_INT24 or CL_FLOAT. Appropriate data format conversion will be done to convert depth valye from a floatingpoint value to actual data format associated with the image. <br> The behavior of write_imagef, write_imagei and write_imageui for image objects created with image_channel_data_type values not specified in the description above or with $(x, y, z)$ coordinate values that are not in the range ( $0 \ldots$ image width $1,0 \ldots$ image height $-1,0 \ldots$ image number of layers -1 ), respectively, is undefined. |
| :---: | :---: |
| void write_imagef ( <br> aQual image3d_t image, <br> int4 coord, <br> float 4 color) <br> void write_imagei ( <br> aQual image3d_t image, <br> int4 coord, <br> int 4 color) <br> void write_imageui ( <br> aQual image3d_t image, <br> int4 coord, <br> uint4 color) | Write color value to location specified by coord.xyz in the 3D image object specified by image. <br> Appropriate data format conversion to the specified image format is done before writing the color value. coord.x, coord.y and coord.z are considered to be unnormalized coordinates and must be in the range 0 $\ldots$ image width $-1,0 \ldots$ image height -1 and $0 \ldots$ image depth -1 . <br> write_imagef can only be used with image objects created with image_channel_data_type set to one of the pre-defined packed formats or set to CL_SNORM_INT8, CL_UNORM_INT8, CL_SNORM_INT16, CL_UNORM_INT16, CL_HALF_FLOAT or CL_FLOAT. Appropriate data format conversion will be done to convert channel data from a floating-point value to actual data format in which the channels are stored. <br> write_imagei can only be used with image objects created with image_channel_data_type set to one of the following values: <br> CL_SIGNED_INT8, <br> CL_SIGNED_INT16 and <br> CL_SIGNED_INT32. <br> write_imageui can only be used with image objects created with image_channel_data_type set to one of the following values: |


|  | CL_UNSIGNED_INT8, |
| :--- | :--- |
| CL_UNSIGNED_INT16 and |  |
| CL_UNSIGNED_INT32. |  |
|  | The behavior of write_imagef, write_imagei and <br> write_imageui for image objects with <br> image_channel_data_type values not specified in <br> the description above or with $(x, y, z)$ coordinate <br> values that are not in the range $(0 \ldots$ image width - <br> $1,0 \ldots$ image height $-1,0 \ldots$ image depth -1$)$ <br> respectively is undefined. |
|  |  |

## Table 6.25 Built-in Image Write Functions

### 6.13.14.5 Built-in Image Query Functions

The following built-in function calls to query image information are supported.
aQual in Table 6.26 refers to one of the access qualifiers. For query functions this may be read_only, write_only or read_write.

| Function | Description |
| :--- | :--- |
| int get_image_width ( <br> aQual image1d_t image) <br> int get_image_width ( <br> aQual image1d_buffer_t image) <br> int get_image_width ( <br> aQual image2d_t image) <br> int get_image_width ( <br> aQual image3d_t image) <br> int get_image_width ( <br> aQual image1d_array_t image) <br> int get_image_width ( <br> aQual image2d_array_t image) <br> int get_image_width ( <br> aQual image2d_depth_t image) <br> int get_image_width ( <br> aQual image2d_array_depth_t image) | Return the image width in pixels. |
|  |  |
| int get_image_height ( |  |


| aQual image2d_t image) <br> int get_image_height ( <br> aQual image3d_t image) <br> int get_image_height ( <br> aQual image2d_array_t image) <br> int get_image_height ( <br> aQual image2d_depth_t image) <br> int get_image_height ( <br> aQual image2d_array_depth_t image) |  |
| :---: | :---: |
| int get_image_depth (image3d_t image) | Return the image depth in pixels. |
| int get_image_channel_data_type ( <br> aQual image1d_t image) int get_image_channel_data_type ( <br> aQual image1d_buffer_t image) int get_image_channel_data_type ( <br> aQual image2d_t image) int get_image_channel_data_type ( <br> aQual image3d_t image) int get_image_channel_data_type ( <br> aQual image1d_array_t image) int get_image_channel_data_type ( <br> aQual image2d_array_t image) int get_image_channel_data_type ( <br> aQual image2d_depth_t image) int get_image_channel_data_type ( aQual image2d_array_depth_t image) | Return the channel data type. Valid values are: <br> CLK_SNORM_INT8 <br> CLK_SNORM_INT16 <br> CLK_UNORM_INT8 <br> CLK_UNORM_INT16 <br> CLK_UNORM_SHORT_565 <br> CLK_UNORM_SHORT_555 <br> CLK_UNORM_SHORT_101010 <br> CLK_SIGNED_INT8 <br> CLK_SIGNED_INT16 <br> CLK_SIGNED_INT32 <br> CLK_UNSIGNED_INT8 <br> CLK_UNSIGNED_INT16 <br> CLK_UNSIGNED_INT32 <br> CLK_HALF_FLOAT <br> CLK_FLOAT |
| int get_image_channel_order ( <br> aQual image1d_t image) int get_image_channel_order ( <br> aQual image1d_buffer_t image) int get_image_channel_order ( <br> aQual image2d_t image) int get_image_channel_order ( <br> aQual image3d_t image) int get_image_channel_order ( <br> aQual image1d_array_t image) int get_image_channel_order ( <br> aQual image2d_array_t image) int get_image_channel_order ( <br> aQual image2d_depth_t image) int get_image_channel_order ( aQual image2d_array_depth_t image) | Return the image channel order. Valid values are: <br> CLK_A <br> CLK_R <br> CLK_Rx <br> CLK_RG <br> CLK_RGx <br> CLK_RA <br> CLK_RGB <br> CLK_RGBx <br> CLK_RGBA <br> CLK_ARGB <br> CLK_BGRA <br> CLK_INTENSITY <br> CLK_LUMINANCE <br> CLK_ABGR |


|  | $\begin{aligned} & \hline \text { CLK_DEPTH } \\ & \text { CLK_sRGB } \\ & \text { CLK_sRGBx } \\ & \text { CLK_sRGBA } \\ & \text { CLK_sBGRA } \\ & \hline \end{aligned}$ |
| :---: | :---: |
| ```int2 get_image_dim ( aQual image2d_t image) int2 get_image_dim ( aQual image2d_array_t image) int2 get_image_dim ( aQual image2d_depth_t image) int2 get_image_dim ( aQual image2d_array_depth_t image)``` | Return the 2D image width and height as an int2 type. The width is returned in the $x$ component, and the height in the $y$ component. |
| int4 get_image_dim ( aQual image3d_t image) | Return the 3D image width, height, and depth as an int 4 type. The width is returned in the $x$ component, height in the $y$ component, depth in the $z$ component and the $w$ component is 0 . |
| size_t get_image_array_size( aQual image2d_array_t image) size_t get_image_array_size( aQual image2d_array_depth_t image) | Return the number of images in the 2D image array. |
| size_t get_image_array_size( aQual image1d_array_t image) | Return the number of images in the 1D image array. |

## Table 6.26 Built-in Image Query Functions

The values returned by get_image_channel_data_type and get_image_channel_order as specified in table 6.26 with the CLK_ prefixes correspond to the CL_ prefixes used to describe the image channel order and data type in tables 5.4 and 5.5. For example, both CL_UNORM_INT8 and CLK_UNORM_INT8 refer to an image channel data type that is an unnormalized unsigned 8-bit integer.

### 6.13.14. 6 Reading and writing to the same image in a kernel

The atomic_work_item_fence (CLK_IMAGE_MEM_FENCE) built-in function can be used to make sure that sampler-less writes are visible to later reads by the same work-item. Only a scope of memory_scope_work_item and an order of memory_order_acq_rel is valid for atomic_work_item_fence when passed the CLK_IMAGE_MEM_FENCE flag. If multiple work-items are writing to and reading from multiple locations in an image, the work_group_barrier (CLK_IMAGE_MEM_FENCE) should be used.

Consider the following example:

```
kernel void
foo(read_write image2d_t img, ... )
{
    int2 coord;
    coord.x = (int)get_global_id(0);
    coord.y = (int)get_global_id(1);
    float4 clr = read_imagef(img, coord);
    write_imagef(img, coord, clr);
    // required to ensure that following read from image at
    // location coord returns the latest color value.
    atomic_work_item_fence(
    CLK IMAGE MEM FENCE,
    memory_order_acq_rel,
    memory_scope_work_item);
    float4 clr_new = read_imagef(img, coord);
```


### 6.13.14.7 Mapping image channels to color values returned by read_image and color values passed to write_image to image channels

The following table describes the mapping of the number of channels of an image element to the appropriate components in the float4, int4 or uint4 vector data type for the color values returned by read_image $\{\mathbf{f}|\mathbf{i}| \mathbf{u i}\}$ or supplied to write_image $\{\mathbf{f}|\mathbf{i}| \mathbf{u i}\}$. The unmapped components will be set to 0.0 for red, green and blue channels and will be set to 1.0 for the alpha channel.
$\left.\begin{array}{|c|c|}\hline \text { Channel Order } & \begin{array}{c}\text { float4, int4 or uint4 } \\ \text { components of channel data }\end{array} \\ \hline \text { CL_R, CL_Rx } & (r, 0.0,0.0,1.0) \\ \hline \text { CL_A } & (0.0,0.0,0.0, a) \\ \hline \text { CL_RG, CL_RGx } & (r, 9,0.0,1.0) \\ \hline \text { CL_RA } & (r, 0.0,0.0, a) \\ \hline \text { CL_RGB, CL_RGBx, } & (r, ~ g, b, 1.0) \\ \hline \text { CL_SRGB, CL_sRGBx }\end{array}\right)$

For CL_DEPTH images, a scalar value is returned by read_imagef or supplied to write_imagef.

NOTE: A kernel that uses a sampler with the CL_ADDRESS_CLAMP addressing mode with multiple images may result in additional samplers being used internally by an implementation. If the same sampler is used with multiple images called via read_image $\{\mathbf{f}|\mathbf{i}| \mathbf{u i}\}$, then it is possible that an implementation may need to allocate an additional sampler to handle the different border color values that may be needed depending on the image formats being used. These implementation allocated samplers will count against the maximum sampler values supported by the device and given by CL_DEVICE_MAX_SAMPLERS. Enqueuing a kernel that requires more samplers than the implementation can support will result in a CL_OUT_OF_RESOURCES error being returned.

### 6.13.15 Work-group Functions

The OpenCL C programming language implements the following built-in functions that operate on a work-group level. These built-in functions must be encountered by all work-items in a work-group executing the kernel. We use the generic type name gentype to indicate the builtin data types half ${ }^{59}$, int, uint, long, ulong, float or double ${ }^{60}$ as the type for the arguments.

| Function | Description |
| :---: | :---: |
| int work_group_all (int predicate) | Evaluates predicate for all work-items in the work-group and returns a non-zero value if predicate evaluates to non-zero for all workitems in the work-group. |
| int work_group_any (int predicate) | Evaluates predicate for all work-items in the work-group and returns a non-zero value if predicate evaluates to non-zero for any workitems in the work-group. |
| gentype work_group_broadcast (gentype $a$,size_t local_id)gentype work_group_broadcast (gentype $a$,size_t local_id_x,size_t local_id_y)gentype work_group_broadcast (gentype $a$, <br> size_t local_id_x, <br> size_t $l o c a l \_i d \_y$, <br> size_t $\left.l o c a l \_i d \_z\right)$ | Broadcast the value of $x$ for work-item identified by local_id to all work-items in the work-group. <br> local_id must be the same value for all workitems in the work-group. |
| gentype work_group_reduce_<op> ( gentype $x$ ) | Return result of reduction operation specified by 〈op> for all values of $x$ specified by workitems in a work-group. |
| gentype <br> work_group_scan_exclusive_<op> ( gentype $x$ ) | Do an exclusive scan operation specified by <op> of all values specified by work-items in the work-group. The scan results are returned for each work-item. <br> The scan order is defined by increasing 1D |

[^43]|  | linear global ID within the work-group. |
| :--- | :--- |
| gentype <br> work_group_scan_inclusive_<op> ( <br> gentype $x)$ | Do an inclusive scan operation specified by <br> <op> of all values specified by work-items in <br> the work-group. The scan results are returned <br> for each work-item. |
|  | The scan order is defined by increasing 1D <br> linear global ID within the work-group. |

## Table 6.27 Built-in Work-group Functions

The <op> in work_group_reduce_<op>, work_group_scan_exclusive_<op> and work_group_scan_inclusive_<op> defines the operator and can be add, min or max.

The inclusive scan operation takes a binary operator $o p$ with an identity I and n (where n is the size of the work-group) elements [ $\left.a_{0}, a_{1}, \ldots a_{n-1}\right]$ and returns [ $a_{0},\left(a_{0} o p a_{1}\right), \ldots$ ( $\mathrm{a}_{0}$ op $\mathrm{a}_{1}$ op ... op $\mathrm{a}_{\mathrm{n}-1}$ )]. If <op> = add, the identity I is 0 . If <op> $=$ min, the identity I is INT_MAX, UINT_MAX, LONG_MAX, ULONG_MAX, for int, uint, long, ulong types and is +INF for floating-point types. Similarly if <op> = max, the identity I is INT_MIN, 0 , LONG_MIN, 0 and -INF.

Consider the following example:

```
void foo(int *p)
{
    int prefix_sum_val = work_group_scan_inclusive_add(
    p[get_local_id(0)]);
}
```

For the example above, let's assume that the work-group size is 8 and $p$ points to the following elements [ $\left.\begin{array}{llllllll}3 & 1 & 7 & 0 & 4 & 1 & 6 & 3\end{array}\right]$. Work-item 0 calls work_group_scan_inclusive_add with 3 and returns 3. Work-item 1 calls work_group_scan_inclusive_add with 1 and returns 4. The full set of values returned by work_group_scan_inclusive_add for work-items 0... 7 are [ 3 4 1111151622 25].

The exclusive scan operation takes a binary associative operator $o p$ with an identity I and n (where n is the size of the work-group) elements [ $\mathrm{a}_{0}, \mathrm{a}_{1}, \ldots a_{\mathrm{n}-1}$ ] and returns [I, $a_{0}$, $\left(\mathrm{a}_{0}\right.$ op $\left.\mathrm{a}_{1}\right)$, ... ( $\mathrm{a}_{0}$ op $\mathrm{a}_{1}$ op ... op $\mathrm{a}_{\mathrm{n}-2}$ ) ]. For the example above, the exclusive scan add operation on the ordered set $\left[\begin{array}{llllllll}3 & 1 & 7 & 0 & 4 & 1 & 6 & 3\end{array}\right]$ would return $\left[\begin{array}{llllll}0 & 3 & 4 & 11 & 11 & 15\end{array}\right.$ 16 22].

NOTE: The order of floating-point operations is not guaranteed for the work_group_reduce_<op>, work_group_scan_inclusive_<op> and work_group_scan_exclusive_<op> built-in functions that operate on half, float and double data
types. The order of these floating-point operations is also non-deterministic for a given workgroup.

### 6.13.16 Pipe Functions

A pipe is identified by specifying the pipe keyword with a type. The data type specifies the size of each packet in the pipe. The pipe keyword is a type modifier. When it is applied to another type $T$, the result is a pipe type whose elements (or packets) are of type $T$. The packet type $\mathbf{T}$ may be any supported OpenCL $C$ scalar and vector integer or floating-point data types, or a user-defined type built from these scalar and vector data types.

Examples:

```
pipe int4 pipeA; // a pipe with int4 packets
pipe user_type_t pipeB; // a pipe with user_type_t packets
```

The read_only (or __read_only) and write_only (or __write_only) qualifiers must be used with the pipe qualifier when a pipe is a parameter of a kernel or of a user-defined function to identify if a pipe can be read from or written to by a kernel and its callees and enqueued child kernels. If no qualifier is specified, read_only is assumed.

A kernel cannot read from and write to the same pipe object. Using the read_write (or read_write) qualifier with the pipe qualifier is a compilation error.

In the following example

```
kernel void
foo (read_only pipe fooA_t pipeA,
    write_ooly pipe foob
{
}
```

pipeA is a read-only pipe object, and pipeB is a write-only pipe object.
The macro CLK_NULL_RESERVE_ID refers to an invalid reservation ID.

### 6.13.16.1 Restrictions

* Pipes can only be passed as arguments to a function (including kernel functions). The C operators (refer to section 6.3 of the OpenCL 2.0 specification) cannot be used with variables declared with the pipe qualifier.
* The pipe qualifier cannot be used with variables declared inside a kernel, a structure or union field, a pointer type, an array, global variables declared in program scope or the return type of a function.


### 6.13.16.2 Built-in Pipe Read and Write Functions

The OpenCL C programming language implements the following built-in functions that read from or write to a pipe. We use the generic type name gentype to indicate the built-in OpenCL C scalar or vector integer or floating-point data types ${ }^{61}$ or any user defined type built from these scalar and vector data types can be used as the type for the arguments to the pipe functions listed in table 6.28.

| Function | Description |
| :---: | :---: |
| $\begin{aligned} & \text { int read_pipe ( } \\ & \text { read_only pipe gentype } p, \\ & \text { gentype *ptr) } \\ & \hline \end{aligned}$ | Read packet from pipe $p$ into $p t r$. Returns 0 if read_pipe is successful and a negative value if the pipe is empty. |
| int write_pipe ( write_only pipe gentype $p$, const gentype *ptr) | Write packet specified by $p t r$ to pipe $p$. Returns 0 if write_pipe is successful and a negative value if the pipe is full. |
| int read_pipe ( <br> read_only pipe gentype $p$, reserve_id_t reserve_id, uint index, gentype *ptr) | Read packet from the reserved area of the pipe referred to by reserve_id and index into ptr. <br> The reserved pipe entries are referred to by indices that go from $0 \ldots$ num_packets -1 . <br> Returns 0 if read_pipe is successful and a negative value otherwise. |
| int write_pipe ( write_only pipe gentype $p$, reserve_id_t reserve_id, uint index, const gentype *ptr) | Write packet specified by ptr to the reserved area of the pipe referred to by reserve_id and index. <br> The reserved pipe entries are referred to by indices that go from $0 \ldots$ num_packets -1 . <br> Returns 0 if write_pipe is successful and a negative value otherwise. |
| ```reserve_id_t reserve_read_pipe ( read_only pipe gentype p uint num_packets) reserve_id_t reserve_write_pipe ( write_only pipe gentype p, uint num_packets)``` | Reserve num_packets entries for reading from or writing to pipe $p$. Returns a valid reservation ID if the reservation is successful. |

[^44]| void commit_read_pipe ( <br> read_only pipe gentype $p$, <br> reserve_id_t reserve_id $)$ | Indicates that all reads and writes to num_packets <br> associated with reservation reserve_id are <br> completed. |
| :---: | :--- |
| void commit_write_pipe $($ <br> write_only pipe gentype $p$, <br> reserve_id_t reserve_id $)$ | Return true if reserve_id is a valid reservation ID <br> and false otherwise. |
| bool is_valid_reserve_id $($ <br> reserve_id_t reserve_id $)$ |  |

Table 6.28 Built-in Pipe Functions

### 6.13.16.3 Built-in Work-group Pipe Read and Write Functions

The OpenCL C programming language implements the following built-in pipe functions that operate at a work-group level. These built-in functions must be encountered by all work-items in a work-group executing the kernel with the same argument values; otherwise the behavior is undefined. We use the generic type name gentype to indicate the built-in OpenCL C scalar or vector integer or floating-point data types ${ }^{62}$ or any user defined type built from these scalar and vector data types can be used as the type for the arguments to the pipe functions listed in table 6.29 .

| Function | Description |
| :---: | :---: |
| reserve_id_t <br> work_group_reserve_read_pipe ( read_only pipe gentype $p$, uint num_packets) <br> reserve_id_t <br> work_group_reserve_write_pipe ( write_only pipe gentype $p$, uint num_packets) | Reserve num_packets entries for reading from or writing to pipe $p$. Returns a valid reservation ID if the reservation is successful. <br> The reserved pipe entries are referred to by indices that go from $0 \ldots$ num_packets -1 . |
| void work_group_commit_read_pipe ( read_only pipe gentype $p$, reserve_id_t reserve_id) <br> void work_group_commit_write_pipe ( write_only pipe gentype $p$, reserve_id_t reserve_id) | Indicates that all reads and writes to num_packets associated with reservation reserve_id are completed. |

Table 6.29 Built-in Pipe Work-group Functions

[^45]NOTE: The read_pipe and write_pipe functions that take a reservation ID as an argument can be used to read from or write to a packet index. These built-ins can be used to read from or write to a packet index one or multiple times. If a packet index that is reserved for writing is not written to using the write_pipe function, the contents of that packet in the pipe are undefined. commit_read_pipe and work_group_commit_read_pipe remove the entries reserved for reading from the pipe. commit_write_pipe and work_group_commit_write_pipe ensures that the entries reserved for writing are all added in-order as one contiguous set of packets to the pipe.

There can only be CL_DEVICE_PIPE_MAX_ACTIVE_RESERVATIONS (refer to table 4.3) reservations active (i.e. reservation IDs that have been reserved but not committed) per workitem or work-group for a pipe in a kernel executing on a device.

Work-item based reservations made by a work-item are ordered in the pipe as they are ordered in the program. Reservations made by different work-items that belong to the same work-group can be ordered using the work-group barrier function. The order of work-item based reservations that belong to different work-groups is implementation defined.

Work-group based reservations made by a work-group are ordered in the pipe as they are ordered in the program. The order of work-group based reservations by different work-groups is implementation defined.

### 6.13.16.4 Built-in Pipe Query Functions

The OpenCL C programming language implements the following built-in query functions for a pipe. We use the generic type name gent ype to indicate the built-in OpenCL C scalar or vector integer or floating-point data types ${ }^{63}$ or any user defined type built from these scalar and vector data types can be used as the type for the arguments to the pipe functions listed in table 6.30 .
aQual in Table 6.30 refers to one of the access qualifiers. For pipe query functions this may be read_only or write_only.

| Function | Description |
| :--- | :--- |
| uint get_pipe_num_packets ( <br> aQual pipe gentype $p$ ) | Returns the number of available entries in the pipe. <br> The number of available entries in a pipe is a <br> dynamic value. The value returned should be <br> considered immediately stale. |
| uint get_pipe_max_packets ( |  |
| aQual pipe gentype $p$ ) |  |$\quad$| Returns the maximum number of packets specified |
| :--- |
| when pipe was created. |

[^46]
## Table 6.30 Built-in Pipe Query Functions

### 6.13.16.5 Restrictions

The following behavior is undefined:
A kernel fails to call reserve_pipe before calling read_pipe or write_pipe that take a reservation ID.

* A kernel calls read_pipe, write_pipe, commit_read_pipe or commit_write_pipe with an invalid reservation ID.
* A kernel calls read_pipe or write_pipe with an valid reservation ID but with an index that is not a value from $0 \ldots$ num_packets -1 specified to the corresponding call to reserve_pipe.
* A kernel calls read_pipe or write_pipe with a reservation ID that has already been committed (i.e. a commit_read_pipe or commit_write_pipe with this reservation ID has already been called).
* A kernel fails to call commit_read_pipe for any reservation ID obtained by a prior call to reserve_read_pipe.
\# A kernel fails to call commit_write_pipe for any reservation ID obtained by a prior call to reserve_write_pipe.
* The contents of the reserved data packets in the pipe are undefined if the kernel does not call write_pipe for all entries that were reserved by the corresponding call to reserve_pipe.
* Calls to read_pipe that takes a reservation ID and commit_read_pipe or write_pipe that takes a reservation ID and commit_write_pipe for a given reservation ID must be called by the same kernel that made the reservation using reserve_read_pipe or reserve_write_pipe. The reservation ID cannot be passed to another kernel including child kernels.


### 6.13.17 Enqueuing Kernels

OpenCL 2.0 allows a kernel to independently enqueue to the same device, without host interaction. A kernel may enqueue code represented by Block syntax, and control execution order with event dependencies including user events and markers. There are several advantages to using the Block syntax: it is more compact; it does not require a cl_kernel object; and enqueuing can be done as a single semantic step.

The following table describes the list of built-in functions that can be used to enqueue a kernel(s).

The macro CLK_NULL_EVENT refers to an invalid device event. The macro CLK_NULL_QUEUE refers to an invalid device queue.

### 6.13.17.1 Built-in Functions - Enqueuing a kernel

## Built-in Function Description

int enqueue_kernel (
queue_t queue,
kernel_enqueue_flags_t flags, const ndrange_t ndrange, void (^block)(void))
int enqueue_kernel (
queue_t queue,
kernel_enqueue_flags_t flags, const ndrange_t ndrange, uint num_events_in_wait_list, const clk_event_t *event_wait_list, clk_event_t *event_ret, void (^block)(void))
int enqueue_kernel (
queue_t queue,
kernel_enqueue_flags_t flags, const ndrange_t ndrange, void (^block)(local void *, ...), uint size $0, \ldots$ )
int enqueue_kernel (
queue_t queue,
kernel_enqueue_flags_t flags, const ndrange_t ndrange, uint num_events_in_wait_list,

Enqueue the block for execution to queue.
If an event is returned, enqueue_kernel performs an implicit retain on the returned event.

```
const clk_event_t *event_wait_list,
clk_event_t *event_ret,
void (^block)(local void *, ..),
uint size0, ...)
```


## Table 6.31 Built-in Kernel Enqueue Functions

The enqueue_kernel built-in function allows a work-item to enqueue a block. Work-items can enqueue multiple blocks to a device queue(s).

The enqueue_kernel built-in function returns CLK_SUCCESS if the block is enqueued successfully and returns CLK_ENQUEUE_FAILURE otherwise. If the -g compile option is specified in compiler options passed to clCompileProgram or clBuildProgram when compiling or building the parent program, the following errors may be returned instead of CLK_ENQUEUE_FAILURE to indicate why enqueue_kernel failed to enqueue the block:

* CLK_INVALID_QUEUE if queue is not a valid device queue.
* CLK_INVALID_NDRANGE if ndrange is not a valid ND-range descriptor or if the program was compiled with -cl-uniform-work-group-size and the local_work_size is specified in ndrange but the global_work_size specified in ndrange is not a multiple of the local_work_size.
* CLK_INVALID_EVENT_WAIT_LIST if event_wait_list is NULL and num_events_in_wait_list > 0, or if event_wait_list is not NULL and num_events_in_wait_list is 0 , or if event objects in event_wait_list are not valid events.
* CLK_DEVICE_QUEUE_FULL if queue is full.

CLK_INVALID_ARG_SIZE if size of local memory arguments is 0 .
CLK_EVENT_ALLOCATION_FAILURE if event_ret is not NULL and an event could not be allocated.

CLK_OUT_OF_RESOURCES if there is a failure to queue the block in queue because of insufficient resources needed to execute the kernel.

Below are some examples of how to enqueue a block.

```
kernel void
my_func_A(global int *a, global int *b, global int *c)
{
}
```

```
kernel void
my_func_B(global int *a, global int *b, global int *c)
{
    ndrange t ndrange;
    // buil\overline{d ndrange information}
    ...
    // example - enqueue a kernel as a block
    enqueue_kernel(get_default_queue(), ndrange,
                                    ^{my_func_A(a, b, c);});
}
kernel void
my_func_C(global int *a, global int *b, global int *c)
{
    ndrange_t ndrange;
    // build ndrange information
    // note that a, b and c are variables in scope of
    // the block
    void (^my_block_A)(void) = ^{my_func_A(a, b, c);};
    // enqueue the block variable
    enqueue_kernel(get_default_queue(),
                                    CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
                                    ndrange,
                                    my_block_A);
    ...
}
```

The example below shows how to declare a block literal and enqueue it.

```
kernel void
my_func(global int *a, global int *b)
{
    ndrange_t ndrange;
    // buil\overline{d} ndrange information
    ...
    // note that a, b and c are variables in scope of
    // the block
    void (^my_block_A)(void) =
        ^{ size_t id = get_global_id(0);
            b[id] += a[id];
```

```
        };
    // enqueue the block variable
    enqueue_kernel(get_default_queue(),
        CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
        ndrange,
        my_block_A);
    // or we could have done the following
    enqueue_kernel(get_default_queue(),
        CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
        ndrange,
        ^{
        size_t id = get_global_id(0);
        b[id] += a[id];
        };
}
```

NOTE: Blocks passed to enqueue_kernel cannot use global variables or stack variables local to the enclosing lexical scope that are a pointer type in the local or private address space.

Example:

```
kernel void
foo(global int *a, local int *lptr, ...)
{
    enqueue_kernel(get_default_queue(),
        CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
        ndrange,
        ^{
            size_t id = get_global_id(0);
            local int *p = lptr; // undefined behavior
        } );
}
```


### 6.13.17.2 Arguments that are a pointer type to local address space

A block passed to enqueue_kernel can have arguments declared to be a pointer to local memory. The enqueue_kernel built-in function variants allow blocks to be enqueued with a variable number of arguments. Each argument must be declared to be a void pointer to local memory. These enqueue_kernel built-in function variants also have a corresponding number of arguments each of type uint that follow the block argument. These arguments specify the size of each local memory pointer argument of the enqueued block.

Some examples follow:

```
kernel void
my_func_A_local_arg1(global int *a, local int *lptr, ...)
{
    ...
}
kernel void
my_func_A_local_arg2(global int *a,
    local int *lptr1, local float4 *lptr2, ...)
{
}
kernel void
my_func_B(global int *a, ...)
{
        ndrange_t ndrange = ndrange_1d(...);
        uint local_mem_size = compute_local_mem_size();
        enqueue_kernel(get_default_queue(),
            CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
                ndrange,
                ^(local void *p) {
                    my_func_A_local_arg1(a, (local int *)p, ...);},
                local_mem_size);
}
kernel void
my_func_C(global int *a, ...)
{
    ...
    ndrange_t ndrange = ndrange_1d(...);
        void (^my_blk_A)(local void *, local void *) =
        ^(local void *lptr1, local void *lptr2){
                my_func_A_local_arg2(
                            a,
                            (local int *)lptr1,
                            (local float4 *)lptr2, ...);};
        // calculate local memory size for lptr
        // argument in local address space for my_blk_A
        uint local_mem_size = compute_local_mem_size();
```

```
    enqueue_kernel(get_default_queue(),
        CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
    ndrange,
    my_blk_A,
    local_mem_size, local_mem_size*4);
}
```


### 6.13.17.3 A Complete Example

The example below shows how to implement an iterative algorithm where the host enqueues the first instance of the nd-range kernel (dp_func_A). The kernel dp_func_A will launch a kernel (evaluate_dp_work_A) that will determine if new nd-range work needs to be performed. If new nd-range work does need to be performed, then evaluate_dp_work_A will enqueue a new instance of $d p \_f u n c \_A$. This process is repeated until all the work is completed.

```
kernel void
dp_func_A(queue_t q, ...)
{
    // queue a single instance of evaluate_dp_work_A to
    // device queue q. queued kernel begins execution after
    // kernel dp_func_A finishes
    if (get_global_id(0) == 0)
    {
        enqueue_kernel(q,
                            CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
                                    ndrange_1d(1),
                                    ^{evaluate_dp_work_A(q, ...);});
    }
}
kernel void
evaluate_dp_work_A(queue_t q,...)
{
    // check if more work needs to be performed
    bool more_work = check_new_work(...);
    if (more_work)
    {
        size_t global_work_size = compute_global_size(...);
        void (^dp_func_A_blk)(void) =
```

```
                    ^{dp_func_A(q, ...});
        // get local WG-size for kernel dp_func_A
        size_t local_work_size =
            get_kernel_work_group_size(dp_func_A_blk);
        // build nd-range descriptor
        ndrange_t ndrange = ndrange_1D(global_work_size,
                        local_work_size);
        // enqueue dp_func_A
        enqueue_kerne\overline{l}
                            CLK_ENQUEUE_FLAGS_WAIT_KERNEL,
    ndrange,
    dp_func_A_blk);
    }
    ...
}
```


### 6.13.17.4 Determining when a child kernel begins execution

The kernel_enqueue_flags_t ${ }^{64}$ argument to enqueue_kernel built-in functions can be used to specify when the child kernel begins execution. Supported values are described in the table below:

| kernel_enqueue_flags_t enum | Description |
| :--- | :--- |
| CLK_ENQUEUE_FLAGS_NO_WAIT | Indicates that the enqueued kernels <br> do not need to wait for the parent <br> kernel to finish execution before <br> they begin execution. |
| CLK_ENQUEUE_FLAGS_WAIT_KERNEL | Indicates that all work-items of the <br> parent kernel must finish executing <br> and all immediate 65 <br> side effects <br> committed before the enqueued <br> child kernel may begin execution. |
| CLK_ENQUEUE_FLAGS_WAIT_WORK_GROUP ${ }^{66}$ | Indicates that the enqueued kernels <br> wait only for the workgroup that <br> enqueued the kernels to finish <br> before they begin execution. |

[^47]
## Table 6.32 Kernel Enqueue Flags

NOTE: The kernel_enqueue_flags_t flags are useful when a kernel enqueued from the host and executing on a device enqueues kernels on the device. The kernel enqueued from the host may not have an event associated with it. The kernel_enqueue_flags_t flags allow the developer to indicate when the child kernels can begin execution.

### 6.13.17.5 Determining when a parent kernel has finished execution

A parent kernel's execution status is considered to be complete when it and all its child kernels have finished execution. The execution status of a parent kernel will be CL_COMPLETE if this kernel and all its child kernels finish execution successfully. The execution status of the kernel will be an error code (given by a negative integer value) if it or any of its child kernels encounter an error, or are abnormally terminated.

For example, assume that the host enqueues a kernel $k$ for execution on a device. Kernel $k$ when executing on the device enqueues kernels $A$ and $B$ to a device queue(s). The enqueue_kernel call to enqueue kernel B specifies the event associated with kernel A in the event_wait_list argument i.e. wait for kernel A to finish execution before kernel B can begin execution. Let's assume kernel $A$ enqueues kernels $X, Y$ and $Z$. kernel $A$ is considered to have finished execution i.e. its execution status is CL_COMPLETE only after A and the kernels A enqueued (and any kernels these enqueued kernels enqueue and so on) have finished execution.

### 6.13.17.6 Built-in Functions - Kernel Query Functions

| Built-in Function | Description |
| :---: | :---: |
| ```uint get_kernel_work_group_size ( void (^block)(void)); uint get_kernel_work_group_size ( void (^block)(local void *, ...));``` | This provides a mechanism to query the maximum work-group size that can be used to execute a block on a specific device given by device. <br> block specifies the block to be enqueued. |
| ```uint get_kernel_preferred_ work_group_size_multiple ( void (^block)(void)); uint get_kernel_preferred_ work_group_size_multiple ( void (^block)(local void *, ...));``` | Returns the preferred multiple of work-group size for launch. This is a performance hint. Specifying a work-group size that is not a multiple of the value returned by this query as the value of the local work size argument to enqueue_kernel will not fail to enqueue the block for execution unless the work-group size specified is larger than the device maximum. |

## Table 6.33 Built-in Kernel Query Functions

### 6.13.17.7 Built-in Functions - Queuing other commands

The following table describes the list of built-in functions that can be used to enqueue commands such as a marker.

| Built-in Function | Description |
| :--- | :--- |
| int enqueue_marker ( <br> queue_t queue, <br> uint num_events_in_wait_list, <br> const clk_event_t *event_wait_list, <br> clk_event_t *event_ret) | Enqueue a marker command to queue. <br> The marker command waits for a list of <br> events specified by event_wait_list to <br> complete before the marker completes. |
| event_ret must not be NULL as otherwise <br> this is a no-op. |  |
| If an event is returned, enqueue_marker <br> performs an implicit retain on the returned <br> event. |  |

Table 6.34 Built-in Other Enqueue Functions
The enqueue_marker built-in function returns CLK_SUCCESS if the marked command is enqueued successfully and returns CLK_ENQUEUE_FAILURE otherwise. If the -g compile option is specified in compiler options passed to clCompileProgram or clBuildProgram, the following errors may be returned instead of CLK_ENQUEUE_FAILURE to indicate why enqueue_marker failed to enqueue the marker command:

CLK_INVALID_QUEUE if queue is not a valid device queue.

* CLK_INVALID_EVENT_WAIT_LIST if event_wait_list is NULL, or if event_wait_list is not NULL and num_events_in_wait_list is 0 , or if event objects in event_wait_list are not valid events.
* CLK_DEVICE_QUEUE_FULL if queue is full.
* CLK_EVENT_ALLOCATION_FAILURE if event_ret is not NULL and an event could not be allocated.
* CLK_OUT_OF_RESOURCES if there is a failure to queue the block in queue because of insufficient resources needed to execute the kernel.


### 6.13.17.8 Built-in Functions - Event Functions

The following table describes the list of built-in functions that work on events.
$\left.\begin{array}{|l|l|}\hline \begin{array}{l}\text { Built-in Function }\end{array} & \text { Description } \\ \hline \begin{array}{c}\text { void retain_event ( } \\ \text { clk_event_t event) }\end{array} & \begin{array}{l}\text { Increments the event refernce count. event } \\ \text { must be an event returned by } \\ \text { enqueue_kernel or enqueue_marker or } \\ \text { a user event. }\end{array} \\ \hline \begin{array}{c}\text { void release_event ( } \\ \text { clk_event_t event) }\end{array} & \begin{array}{l}\text { Decrements the event reference count. The } \\ \text { event object is deleted once the event reference } \\ \text { count is zero, the specific command identified by } \\ \text { this event has completed (or terminated) and } \\ \text { there are no commands in any device command } \\ \text { queue that require a wait for this event to } \\ \text { complete. }\end{array} \\ \hline \begin{array}{l}\text { event must be an event returned by }\end{array} \\ \hline \text { enqueue_kernel, enqueue_marker or a } \\ \text { user event. }\end{array}\right\}$

|  | CLK_PROFILING_COMMAND_EXEC_TIME |
| :--- | :--- |
|  | value is a pointer to two 64-bit values. |
|  | The first 64-bit value describes the elapsed time |
| CL_PROFILING_COMMAND_END - |  |
| CL_PROFLING_COMMAND_START for the |  |
| command identified by event in nanoseconds. |  |
|  | The second 64-bit value describes the elapsed <br> time CL_PROFILING_COMMAND_COMPLETE <br> - CL_PROFILING_COMAMND_START for the <br> command identified by event in nanoseconds. |
|  | NOTE: The behavior of <br> capture_event_profiling_info when <br> called multiple times for the same event is <br> undefined. |

## Table 6.35 Built-in Event Functions

Events can be used to identify commands enqueued to a command-queue from the host. These events created by the OpenCL runtime can only be used on the host i.e. as events passed in event_wait_list argument to various clEnqueue APIs or runtime APIs that take events as arguments such as clRetainEvent, clReleaseEvent, clGetEventProfilingInfo.

Similarly, events can be used to identify commands enqueued to a device queue (from a kernel). These event objects cannot be passed to the host or used by OpenCL runtime APIs such as the clEnqueueAPIs or runtime APIs that take event arguments.
clRetainEvent and clReleaseEvent will return CL_INVALID_OPERATION if event specified is an event that refers to any kernel enqueued to a device queue using enqueue_kernel or enqueue_marker or is a user event created by create_user_event.

Similarly, clSetUserEventStatus can only be used to set the execution status of events created using clCreateUserEvent. User events created on the device can be set using set_user_event_status built-in function.

The example below shows how events can be used with kernels enqueued to multiple device queues.

```
extern void barA_kernel(...);
extern void barB_kernel(...);
kernel void
```

```
foo(queue_t q0, queue q1, ...)
{
    clk_event_t evt0;
    // enqueue kernel to queue q0
    enqueue_kernel(q0,
        CLK_ENQUEUE_FLAGS_NO_WAIT,
        ndrange_A,
        0, NULL, &evt0,
        ^{barA_kernel(...);} );
    // enqueue kernel to queue q1
    enqueue_kernel(q1,
        CLK_ENQUEUE_FLAGS_NO_WAIT,
        ndrange_B,
        1, &evt0, NULL,
        ^{barB_kernel(...);} );
    // release event evt0. This will get released
    // after barA_kernel enqueued in queue q0 has finished
    // execution and barB_kernel enqueued in queue q1 and
    // waits for evt0 is submitted for execution i.e. wait
    // for evtO is satisfied.
    release_event(evt0);
}
```

The example below shows how the marker command can be used with kernels enqueued to a device queue.

```
kernel void
foo(queue_t q, ...)
{
    ...
    clk_event_t marker_event;
    clk_event_t events[2];
    enqueue_kernel(q,
                            CLK_ENQUEUE_FLAGS_NO_WAIT,
        ndrange,
        0, NULL, &events[0],
        ^{barA_kernel(...);} );
    enqueue_kernel(q,
    CLK_ENQUEUE_FLAGS_NO_WAIT,
        ndrange,
        0, NULL, &events[1],
```

```
        ^{barB_kernel(...);} );
    // barA kernel and barB kernel can be executed
    // out of order. we need to wait for both these
    // kernels to finish execution before barC_kernel
    // starts execution so we enqueue a marker command and
    // then enqueue barC_kernel that waits on the event
    // associated with the marker.
    enqueue_marker(q, 2, events, &marker_event);
    enqueue_kernel(q,
        CLK_ENQUEUE_FLAGS_NO_WAIT,
        1, &marker_event, NULL,
        ^{barC_kernel(...);} );
    release event(events[0];
    release_event(events[1]);
    release_event(marker_event);
}
```


### 6.13.17.9 Built-in Functions - Helper Functions

| Built-in Function | Description |
| :--- | :--- |
| queue_t get_default_queue (void) | Returns the default device queue. If a <br> default device queue has not been created, <br> CLK_NULL_QUEUE is returned. |
| ndrange_t ndrange_1D ( <br> size_t global_work_size) | Builds a 1D, 2D or 3D ND-range <br> descriptor. |
| ndrange_t ndrange_1D ( |  |
| size_t global_work_size, |  |
| size_t local_work_size) |  |
| ndrange_t ndrange_1D ( |  |
| size_t global_work_offset, |  |
| size_t global_work_size, |  |
| size_t local_work_size) |  |$\quad$| ndrange_t ndrange_2D ( |
| :--- |
| const size_t global_work_size[2]) |

```
    const size_t global_work_size[2],
    const size_t local_work_size[2])
ndrange_t ndrange_2D (
    const size_t global_work_offset[2],
    const size_t global_work_size[2],
    const size_t local_work_size[2])
ndrange_t ndrange_3D (
    const size_t global_work_size[3])
ndrange_t ndrange_3D (
    const size_t global_work_size[3],
    const size_t local_work_size[3])
ndrange_t ndrange_3D (
    const size_t global_work_offset[3],
    const size_t global_work_size[3],
    const size_t local_work_size[3])
```

Table 6.36 Built-in Helper Functions

## 7. OpenCL Numerical Compliance

This section describes features of the C99 and IEEE 754 standards that must be supported by all OpenCL compliant devices.

This section describes the functionality that must be supported by all OpenCL devices for single precision floating-point numbers. Currently, only single precision floating-point is a requirement. Double precision floating-point is an optional feature.

### 7.1 Rounding Modes

Floating-point calculations may be carried out internally with extra precision and then rounded to fit into the destination type. IEEE 754 defines four possible rounding modes:

Round to nearest even

* Round toward $+\infty$

4 Round toward $-\infty$

* Round toward zero

Round to nearest even is currently the only rounding mode required ${ }^{67}$ by the OpenCL specification for single precision and double precision operations and is therefore the default rounding mode. In addition, only static selection of rounding mode is supported. Dynamically reconfiguring the rounding modes as specified by the IEEE 754 spec is unsupported.

### 7.2 INF, NaN and Denormalized Numbers

INF and NaNs must be supported. Support for signaling NaNs is not required.
Support for denormalized numbers with single precision floating-point is optional.
Denormalized single precision floating-point numbers passed as input or produced as the output of single precision floating-point operations such as add, sub, mul, divide, and the functions defined in sections 6.11 .2 (math functions), 6.11.4 (common functions) and 6.11.5 (geometric functions) may be flushed to zero.

[^48]
### 7.3 Floating-Point Exceptions

Floating-point exceptions are disabled in OpenCL. The result of a floating-point exception must match the IEEE 754 spec for the exceptions not enabled case. Whether and when the implementation sets floating-point flags or raises floating-point exceptions is implementationdefined. This standard provides no method for querying, clearing or setting floating-point flags or trapping raised exceptions. Due to non-performance, non-portability of trap mechanisms and the impracticality of servicing precise exceptions in a vector context (especially on heterogeneous hardware), such features are discouraged.

Implementations that nevertheless support such operations through an extension to the standard shall initialize with all exception flags cleared and the exception masks set so that exceptions raised by arithmetic operations do not trigger a trap to be taken. If the underlying work is reused by the implementation, the implementation is however not responsible for reclearing the flags or resetting exception masks to default values before entering the kernel. That is to say that kernels that do not inspect flags or enable traps are licensed to expect that their arithmetic will not trigger a trap. Those kernels that do examine flags or enable traps are responsible for clearing flag state and disabling all traps before returning control to the implementation. Whether or when the underlying work-item (and accompanying global floating-point state if any) is reused is implementation-defined.

The expressions math_errorhandling and MATH_ERREXCEPT are reserved for use by this standard, but not defined. Implementations that extend this specification with support for floating-point exceptions shall define math_errorhandling and MATH_ERREXCEPT per ISO / IEC 9899 : TC2.

### 7.4 Relative Error as ULPs

In this section we discuss the maximum relative error defined as ulp (units in the last place). Addition, subtraction, multiplication, fused multiply-add and conversion between integer and a single precision floating-point format are IEEE 754 compliant and are therefore correctly rounded. Conversion between floating-point formats and explicit conversions specified in section 6.2.3 must be correctly rounded.

The ULP is defined as follows:

```
If x is a real number that lies between two finite
consecutive floating-point numbers a and b, without being
equal to one of them, then ulp(x) = |b - a|, otherwise
ulp(x) is the distance between the two non-equal finite
floating-point numbers nearest x. Moreover, ulp(NaN) is
NaN.
```

Attribution: This definition was taken with consent from Jean-Michel Muller with slight clarification for behavior at zero. Refer to ftp://ftp.inria.fr/INRIA/publication/publi-pdf/RR/RR5504.pdf.

Table $7.1^{68}$ describes the minimum accuracy of single precision floating-point arithmetic operations given as ULP values. The reference value used to compute the ULP value of an arithmetic operation is the infinitely precise result.

| Function | Min Accuracy - ULP values ${ }^{\mathbf{6 9}}$ |
| ---: | :--- |
| $\boldsymbol{x}+\boldsymbol{y}$ | Correctly rounded |
| $\boldsymbol{x}-\boldsymbol{y}$ | Correctly rounded |
| $\boldsymbol{x} * \boldsymbol{y}$ | Correctly rounded |
| $\mathbf{1 . 0} / \boldsymbol{x}$ | $<=2.5$ ulp |
| $\boldsymbol{x} / \boldsymbol{y}$ | $<=2.5 \mathrm{ulp}$ |
|  |  |
| acos | $<=4 \mathrm{ulp}$ |
| acospi | $<=5 \mathrm{ulp}$ |
| asin | $<=4$ ulp |
| asinpi | $<=5 \mathrm{ulp}$ |
| atan | $<=5$ ulp |
| atan2 | $<=6$ ulp |
| atanpi | $<=5$ ulp |
| atan2pi | $<=6$ ulp |
| acosh | $<=4$ ulp |
| asinh | $<=4$ ulp |
| atanh | $<=5 \mathrm{ulp}$ |
| cbrt | $<=2$ ulp |
| ceil | Correctly rounded |
| copysign | 0 ulp |
| cos | $<=4$ ulp |
| cosh | $<=4$ ulp |
| cospi | $<=4$ ulp |
| erfc | $<=16$ ulp |
| erf | $<=16$ ulp |
| exp | $<=3$ ulp |
| exp2 | $<=3$ ulp |
| exp10 | $<=3$ ulp |
| expm1 | $<=3$ ulp |
| fabs | 0 ulp |
| fdim | Correctly rounded |
| floor | Correctly rounded |
|  |  |

[^49]| fma | Correctly rounded |
| :---: | :---: |
| fmax | 0 ulp |
| fmin | 0 ulp |
| fmod | 0 ulp |
| fract | Correctly rounded |
| frexp | 0 ulp |
| hypot | <= 4 ulp |
| ilogb | 0 ulp |
| Idexp | Correctly rounded |
| $\log$ | <= 3 ulp |
| $\log 2$ | <= 3 ulp |
| $\log 10$ | <=3 ulp |
| $\log 1 \mathrm{p}$ | <= 2 ulp |
| logb | 0 ulp |
| mad | Implemented either as a correctly rounded fma or as a multiply followed by an add both of which are correctly rounded |
| maxmag | 0 ulp |
| minmag | 0 ulp |
| modf | 0 ulp |
| nan | 0 ulp |
| nextafter | 0 ulp |
| pow(x, y) | <= 16 ulp |
| pown(x, y) | <= 16 ulp |
| powr(x, y) | <= 16 ulp |
| remainder | 0 ulp |
| remquo | 0 ulp |
| rint | Correctly rounded |
| rootn | <= 16 ulp |
| round | Correctly rounded |
| rsqrt | <=2 ulp |
| sin | <= 4 ulp |
| sincos | <=4 ulp for sine and cosine values |
| sinh | <=4 ulp |
| sinpi | <=4 ulp |
| sqrt | <= 3 ulp |
| tan | <= 5 ulp |
| tanh | <= 5 ulp |
| tanpi | <=6 ulp |
| tgamma | <= 16 ulp |
| trunc | Correctly rounded |
|  |  |
| half_cos | <= 8192 ulp |
| half_divide | <= 8192 ulp |


| half_exp | $<=8192$ ulp |
| ---: | :--- |
| half_exp2 | $<=8192 \mathrm{ulp}$ |
| half_exp10 | $<=8192 \mathrm{ulp}$ |
| half_log | $<=8192 \mathrm{ulp}$ |
| half_log2 | $<=8192 \mathrm{ulp}$ |
| half_log10 | $<=8192 \mathrm{ulp}$ |
| half_powr | $<=8192 \mathrm{ulp}$ |
| half_recip | $<=8192 \mathrm{ulp}$ |
| half_rsqrt | $<=8192 \mathrm{ulp}$ |
| half_sin | $<=8192$ ulp |
| half_sqrt | $<=8192$ ulp |
| half_tan | $<=8192$ ulp |
| native_cos | Implementation-defined |
| native_divide | Implementation-defined |
| native_exp | Implementation-defined |
| native_exp2 | Implementation-defined |
| native_exp10 | Implementation-defined |
| native_log | Implementation-defined |
| native_log2 | Implementation-defined |
| native_log10 | Implementation-defined |
| native_powr | Implementation-defined |
| native_recip | Implementation-defined |
| native_rsqrt | Implementation-defined |
| native_sin | Implementation-defined |
| native_sqrt | Implementation-defined |
| native_tan | Implementation-defined |

Table 7.1 ULP values for single precision built-in math functions
Table 7.2 describes the minimum accuracy of commonly used single precision floating-point arithmetic operations given as ULP values if the -cl-fast-relaxed-math compiler option is specified when compiling or building an OpenCL program. The minimum accuracy of math functions not defined in table 7.2 when the -cl-fast-relaxed-math compiler option is specified is as defined in table 7.1. The reference value used to compute the ULP value of an arithmetic operation is the infinitely precise result.

| Function | Min Accuracy - ULP values ${ }^{70}$ |
| ---: | :--- |
| $\mathbf{1 . 0} / \boldsymbol{x}$ | $<=2.5$ ulp for $x$ in the domain of $2^{-126}$ to $2^{126}$ |
| $\boldsymbol{x} / \boldsymbol{y}$ | $<=2.5$ ulp for $x$ in the domain of $2^{-62}$ to $2^{62}$ and $y$ in <br> the domain of $2^{-62}$ to $2^{62}$. |
| $\operatorname{acos}(\boldsymbol{x})$ | $<=4096$ ulp |

[^50]| $\operatorname{acospi}(x)$ | Implemented as $\operatorname{acos}(\mathbf{x})$ * M_PI_F. For nonderived implementations, the error is $<=8192$ ulp. |
| :---: | :---: |
| $\operatorname{asin}(x)$ | <= 4096 ulp |
| $\operatorname{asinpi}(x)$ | Implemented as $\operatorname{asin}(\mathbf{x}) *$ M_PI_F. For nonderived implementations, the error is $<=8192$ ulp. |
| $\operatorname{atan}(x)$ | <= 4096 ulp |
| $\operatorname{atan} 2(y, x)$ | Implemented as $\operatorname{atan}(\mathbf{y} / \mathbf{x})$ for $x>0, \operatorname{atan}(\mathbf{y} / \mathbf{x})+$ M_PI_F for $x<0$ and $y>0$ and $\operatorname{atan}(\mathbf{y} / \mathbf{x})$ M_PI_F for $\mathrm{x}<0$ and $\mathrm{y}<0$. |
| $\operatorname{atanpi}(x)$ | Implemented as atan(x) * M_PI_F. For nonderived implementations, the error is $<=8192$ ulp. |
| $\operatorname{atan} 2 \mathrm{pi}(y, x)$ | Implemented as $\operatorname{atan} 2(\mathbf{y}, \mathbf{x}) *$ M_PI_F. For nonderived implementations, the error is $<=8192$ ulp. |
| $\operatorname{acosh}(x)$ | Implemented as $\boldsymbol{\operatorname { l o g }}\left(\mathrm{x}+\operatorname{sqrt}\left(\mathbf{x}^{*} \mathrm{x}-\mathbf{1}\right)\right.$ ). |
| $\operatorname{asinh}(x)$ | Implemented as $\log \left(\mathbf{x}+\operatorname{sqrt}\left(\mathbf{x}^{*} \mathbf{x}+\mathbf{1}\right)\right.$ ). |
| $\operatorname{cbrt}(\boldsymbol{x})$ | Implemented as rootn( $\mathbf{x}, \mathbf{3}$ ). For non-derived implementations, the error is $<=8192 \mathrm{ulp}$. |
| $\cos (x)$ | For $x$ in the domain $[-\pi, \pi]$, the maximum absolute error is $<=2^{-11}$ and larger otherwise. |
| $\cosh (x)$ | Implemented as $0.5 * \exp (\mathbf{x})+\exp (-\mathbf{x})$. For nonderived implementations, the error is $<=8192$ ulp. |
| $\operatorname{cospi}(x)$ | For $x$ in the domain $[-1,1]$, the maximum absolute error is $<=2^{-11}$ and larger otherwise. |
| $\exp (x)$ | $<=3+$ floor(fabs( $2 * x$ ) ) ulp |
| $\exp 2(x)$ | <= $3+$ floor(fabs( $2 * x$ ) ) ulp |
| $\exp 10(x)$ | Derived implementations implement this as $\exp 2(x$ * $\log 2(10)$ ). For non-derived implementations, the error is $<=8192 \mathrm{ulp}$. |
| $\operatorname{expm} 1(x)$ | Derived implementations implement this as $\exp (x)$ -1 . For non-derived implementations, the error is <= 8192 ulp . |
| $\log (x)$ | For $x$ in the domain [0.5,2] the maximum absolute error is $\left\langle=2^{-21}\right.$; otherwise the maximum error is $<=3$ ulp for the full profile and $<=4$ ulp for the embedded profile |
| $\log 2(x)$ | For $x$ in the domain $[0.5,2]$ the maximum absolute error is $<=2^{-21}$; otherwise the maximum error is $<=3$ ulp for the full profile and $<=4$ ulp for the embedded profile |
| $\log 10(x)$ | For $x$ in the domain [0.5,2] the maximum absolute error is $<=2^{-21}$; otherwise the maximum error is $<=3$ ulp for the full profile and $<=4$ ulp for the embedded profile |
| $\log 1 \mathrm{p}(x)$ | Derived implementations implement this as $\log (x+$ <br> 1). For non-derived implementations, the error is <= 8192 ulp. |


| pow(x, y) | Undefined for $\mathrm{x}=0$ and $\mathrm{y}=0$ or for $\mathrm{x}<0$ and noninteger $y$. For $x>=0$ or $x<0$ and even $y$, derived implementations implement this as $\exp 2(y *$ $\log 2(x))$. For $\mathrm{x}<0$ and odd y , derived implementations implement this as -exp2 $(y *$ $\log 2(\mathbf{f a b s}(x))$. For non-derived implementations, the error is $<=8192$ ulp. |
| :---: | :---: |
| $\operatorname{pown}(\mathrm{x}, \mathrm{y})$ | Defined only for integer values of $y$. Undefined for $\mathrm{x}=0$ and $\mathrm{y}=0$. For $\mathrm{x}>=0$ or $\mathrm{x}<0$ and even y , derived implementations implement this as $\exp 2(y$ $* \log 2(x)$ ). For $\mathrm{x}<0$ and odd y , derived implementations implement this as $-\exp 2(y *$ $\log 2(\mathbf{f a b s}(x))$. For non-derived implementations, the error is $<=8192$ ulp. |
| powr(x, y) | Defined only for $\mathrm{x}>=0$. Undefined for $\mathrm{x}=0$ and y $=0$. Derived implementations implement this as $\exp 2(y * \log 2(x))$. For non-derived implementations, the error is <= 8192 ulp. |
| $\operatorname{rootn}(x, y)$ | Defined for $\mathrm{x}>0$ and for $\mathrm{x}<0$ when y is odd. Undefined for $\mathrm{x}=0$ and $\mathrm{y}=0$. Derived implementations implement this as $\exp 2(\log 2(x) /$ $y)$ for $\mathrm{x}>0$. Derived implementations implement this as $-\exp 2(\log 2(-x) / y)$ for $\mathrm{x}<0$. For nonderived implementations, the error is $<=8192$ ulp. |
| $\sin (x)$ | For $x$ in the domain $[-\pi, \pi]$, the maximum absolute error is $<=2^{-11}$ and larger otherwise. |
| $\sin \cos (x)$ | ulp values as defined for $\sin (x)$ and $\cos (x)$ |
| $\sinh (x)$ | Implemented as $0.5 * \exp (\mathbf{x})-\exp (-\mathbf{x})$. For nonderived implementations, the error is $<=8192$ ulp. |
| $\operatorname{sinpi}(x)$ | For $x$ in the domain $[-1,1]$, the maximum absolute error is $<=2^{-11}$ and larger otherwise. |
| $\boldsymbol{\operatorname { t a n }}(x)$ | Derived implementations implement this as $\sin (\boldsymbol{x})$ * (1.0f $/ \cos (\boldsymbol{x})$ ). For non-derived implementations, the error is <= 8192 ulp. |
| $\operatorname{tanpi}(x)$ | Derived implementations implement this as $\boldsymbol{\operatorname { t a n }}(\boldsymbol{x} *$ M_PI_F). For non-derived implementations, the error is $<=8192$ ulp for $x$ in the domain $[-1,1]$. |
| $x * y+z$ | Implemented either as a correctly rounded fma or as a multiply and an add both of which are correctly rounded. |

Table 7.2 ULP values for single precision built-in math functions with fast relaxed math

Table 7.3 describes the minimum accuracy of double precision floating-point arithmetic operations given as ULP values. The reference value used to compute the ULP value of an
arithmetic operation is the infinitely precise result.

| Function | Min Accuracy - ULP values ${ }^{\mathbf{7 1}}$ |
| :---: | :---: |
| $x+y$ | Correctly rounded |
| $x-y$ | Correctly rounded |
| $x * y$ | Correctly rounded |
| $1.0 / x$ | Correctly rounded |
| $x / y$ | Correctly rounded |
|  |  |
| acos | <= 4 ulp |
| acospi | <= 5 ulp |
| asin | <= 4 ulp |
| asinpi | <= 5 ulp |
| atan | <= 5 ulp |
| $\operatorname{atan} 2$ | <= 6 ulp |
| atanpi | <= 5 ulp |
| atan2pi | <=6 ulp |
| acosh | <= 4 ulp |
| asinh | <= 4 ulp |
| atanh | <= 5 ulp |
| cbrt | <=2 ulp |
| ceil | Correctly rounded |
| copysign | 0 ulp |
| cos | <= 4 ulp |
| cosh | <=4 ulp |
| cospi | <= 4 ulp |
| erfc | <= 16 ulp |
| erf | <= 16 ulp |
| exp | <= 3 ulp |
| exp2 | <= 3 ulp |
| $\exp 10$ | <= 3 ulp |
| expm1 | <= 3 ulp |
| fabs | 0 ulp |
| fdim | Correctly rounded |
| floor | Correctly rounded |
| fma | Correctly rounded |
| fmax | 0 ulp |
| fmin | 0 ulp |
| fmod | 0 ulp |
| fract | Correctly rounded |
| frexp | 0 ulp |
| hypot | <= 4 ulp |
| ilogb | 0 ulp |

[^51]| $\operatorname{ldexp}$ | Correctly rounded |
| ---: | :--- |
| $\mathbf{\operatorname { l o g }}$ | $<=3 \mathrm{ulp}$ |
| $\mathbf{l o g} 2$ | $<=3 \mathrm{ulp}$ |
| $\mathbf{l o g} 10$ | $<=3 \mathrm{ulp}$ |
| $\mathbf{l o g 1 p}$ | $<=2$ ulp |
| $\mathbf{l o g b}$ | 0 ulp |
| mad | Any value allowed (infinite ulp) |
| maxmag | 0 ulp |
| minmag | 0 ulp |
| modf | 0 ulp |
| nan | 0 ulp |
| nextafter | 0 ulp |
| pow(x, $\mathbf{y})$ | $<=16$ ulp |
| pown(x, y) | $<=16$ ulp |
| powr(x, y) | $<=16$ ulp |
| remainder | 0 ulp |
| remquo | 0 ulp |
| rint | Correctly rounded |
| rootn | $<=16$ ulp |
| round | Correctly rounded |
| rsqrt | $<=2$ ulp |
| sin | $<=4$ ulp |
| sincos | $<=4$ ulp for sine and cosine values |
| sinh | $<=4$ ulp |
| sinpi | $<=4$ ulp |
| sqrt | Correctly rounded |
| tan | $<=5$ ulp |
| tanh | $<=5$ ulp |
| tanpi | $<=6$ ulp |
| tgamma | $<=16$ ulp |
| trunc | Correctly rounded |
|  |  |
|  |  |
|  |  |

Table 7.3 ULP values for double precision built-in math functions

### 7.5 Edge Case Behavior

The edge case behavior of the math functions (section 6.13.2) shall conform to sections F. 9 and G. 6 of ISO/IEC 9899:TC 2 (commonly known as C99, TC2), except where noted below in section 7.5.1.

### 7.5.1 Additional Requirements Beyond C99 TC2

Functions that return a NaN with more than one NaN operand shall return one of the NaN operands. Functions that return a NaN operand may silence the NaN if it is a signaling NaN . A non-signaling NaN shall be converted to a non-signaling NaN. A signaling NaN shall be converted to a NaN, and should be converted to a non-signaling NaN. How the rest of the NaN payload bits or the sign of NaN is converted is undefined.
half_<funcname> functions behave identically to the function of the same name without the half_ prefix. They must conform to the same edge case requirements (see sections F. 9 and G. 6 of C99, TC2). For other cases, except where otherwise noted, these single precision functions are permitted to have up to 8192 ulps of error (as measured in the single precision result), although better accuracy is encouraged.

The usual allowances for rounding error (section 7.4) or flushing behavior (section 7.5.3) shall not apply for those values for which section F. 9 of C99, TC2, or sections 7.5.1 and 7.5.3 below (and similar sections for other floating-point precisions) prescribe a result (e.g. ceil ( $-1<$ $\mathrm{x}<0$ ) returns -0 ). Those values shall produce exactly the prescribed answers, and no other. Where the $\pm$ symbol is used, the sign shall be preserved. For example, $\sin ( \pm 0)= \pm 0$ shall be interpreted to mean $\sin (+0)$ is +0 and $\sin (-0)$ is -0 .

```
\(\operatorname{acospi}(1)=+0\).
\(\operatorname{acospi}(x)\) returns a NaN for \(|x|>1\).
asinpi \(( \pm 0)= \pm 0\).
asinpi \((x)\) returns a NaN for \(|x|>1\).
atanpi \(( \pm 0)= \pm 0\).
atanpi \(( \pm \infty)= \pm 0.5\).
atan2pi \(( \pm 0,-0)= \pm 1\).
\(\operatorname{atan} 2 \mathbf{p i}( \pm 0,+0)= \pm 0\).
atan2pi \(( \pm 0, x)\) returns \(\pm 1\) for \(x<0\).
\(\operatorname{atan} 2 \mathbf{p i}( \pm 0, x)\) returns \(\pm 0\) for \(x>0\).
atan2pi \((y, \pm 0)\) returns -0.5 for \(y<0\).
atan2pi \((y, \pm 0)\) returns 0.5 for \(y>0\).
\(\boldsymbol{\operatorname { a t a n }} 2 \mathbf{p i}( \pm y,-\infty)\) returns \(\pm 1\) for finite \(y>0\).
atan2pi \(( \pm y,+\infty)\) returns \(\pm 0\) for finite \(y>0\).
atan2pi \(( \pm \infty, x)\) returns \(\pm 0.5\) for finite \(x\).
atan2pi \(( \pm \infty,-\infty)\) returns \(\pm 0.75\).
atan2pi \(( \pm \infty,+\infty)\) returns \(\pm 0.25\).
ceil \((-1<x<0)\) returns -0 .
cospi ( \(\pm 0\) ) returns 1
cospi \((n+0.5)\) is +0 for any integer \(n\) where \(n+0.5\) is representable.
cospi \(( \pm \infty)\) returns a NaN .
```

$\exp 10( \pm 0)$ returns 1.
$\exp 10(-\infty)$ returns +0.
$\exp 10(+\infty)$ returns $+\infty$.
distance $(x, y)$ calculates the distance from $x$ to $y$ without overflow or extraordinary precision loss due to underflow.
fdim ( any, NaN ) returns NaN .
fdim ( NaN , any) returns NaN .
fmod ( $\pm 0, \mathrm{NaN})$ returns NaN .
frexp ( $\pm \infty, \exp$ ) returns $\pm \infty$ and stores 0 in $\exp$.
frexp ( $\mathrm{NaN}, \exp$ ) returns the NaN and stores 0 in exp.
fract ( $x$, iptr) shall not return a value greater than or equal to 1.0 , and shall not return a value less than 0 .
fract $(+0, i p t r)$ returns +0 and +0 in iptr.
fract ( -0 , iptr ) returns -0 and -0 in iptr.
fract ( + inf, iptr ) returns +0 and + inf in iptr.
fract ( -inf , iptr ) returns -0 and -inf in iptr.
fract ( NaN, iptr ) returns the NaN and NaN in iptr.
length calculates the length of a vector without overflow or extraordinary precision loss due to underflow.
lgamma_r $(x$, signp $)$ returns 0 in signp if $x$ is zero or a negative integer.
nextafter $(-0, y>0)$ returns smallest positive denormal value.
nextafter $(+0, y<0)$ returns smallest negative denormal value.
normalize shall reduce the vector to unit length, pointing in the same direction without overflow or extraordinary precision loss due to underflow.
normalize $(v)$ returns $v$ if all elements of $v$ are zero.
normalize $(v)$ returns a vector full of NaNs if any element is a NaN .
normalize $(v)$ for which any element in $v$ is infinite shall proceed as if the elements in $v$ were replaced as follows:

$$
\begin{aligned}
\text { for }(\mathrm{i}= & 0 ; \mathrm{i}<\operatorname{sizeof}(v) / \operatorname{sizeof}(v[0]) ; \mathrm{i}++) \\
& v[\mathrm{i}]=\operatorname{isinf}(v[\mathrm{i}]) ? \operatorname{copysign}(1.0, v[\mathrm{i}]): 0.0 * v[\mathrm{i}] ;
\end{aligned}
$$

pow ( $\pm 0,-\infty$ ) returns $+\infty$
pown $(x, 0)$ is 1 for any $x$, even zero, NaN or infinity.
pown $( \pm 0, n)$ is $\pm \infty$ for odd $n<0$.
pown $( \pm 0, n)$ is $+\infty$ for even $n<0$.
pown $( \pm 0, n)$ is +0 for even $n>0$.
pown $( \pm 0, n)$ is $\pm 0$ for odd $n>0$.
powr $(x, \pm 0)$ is 1 for finite $x>0$.
$\operatorname{powr}( \pm 0, y)$ is $+\infty$ for finite $y<0$.
powr $( \pm 0,-\infty)$ is $+\infty$.
powr $( \pm 0, y)$ is +0 for $y>0$.
powr $(+1, y)$ is 1 for finite $y$.
powr $(x, y)$ returns NaN for $x<0$.
powr ( $\pm 0, \pm 0$ ) returns NaN .
powr $(+\infty, \pm 0)$ returns NaN .
powr $(+1, \pm \infty)$ returns NaN .
powr $(x, \mathrm{NaN})$ returns the NaN for $x>=0$.
powr ( $\mathrm{NaN}, y$ ) returns the NaN .
$\boldsymbol{\operatorname { r i n t }}(-0.5<=x<0)$ returns -0 .
remquo ( $x, y, \& q u o$ ) returns a NaN and 0 in quo if $x$ is $\pm \infty$, or if $y$ is 0 and the other argument is non- NaN or if either argument is a NaN .
$\operatorname{rootn}( \pm 0, n)$ is $\pm \infty$ for odd $n<0$.
$\operatorname{rootn}( \pm 0, n)$ is $+\infty$ for even $n<0$.
$\operatorname{rootn}( \pm 0, n)$ is +0 for even $n>0$.
$\operatorname{rootn}( \pm 0, n)$ is $\pm 0$ for odd $n>0$.
$\operatorname{rootn}(x, n)$ returns a NaN for $x<0$ and $n$ is even.
$\operatorname{rootn}(x, 0)$ returns a NaN .
round $(-0.5<x<0)$ returns -0 .
$\operatorname{sinpi}( \pm 0)$ returns $\pm 0$.
$\operatorname{sinpi}(+n)$ returns +0 for positive integers $n$.
$\operatorname{sinpi}(-n)$ returns -0 for negative integers $n$.
$\operatorname{sinpi}( \pm \infty)$ returns a NaN .
tanpi ( $\pm 0$ ) returns $\pm 0$.
tanpi $( \pm \infty)$ returns a NaN .
tanpi ( $n$ ) is copysign $(0.0, n$ ) for even integers $n$.
tanpi $(n)$ is copysign $(0.0,-n)$ for odd integers $n$.
tanpi $(n+0.5)$ for even integer $n$ is $+\infty$ where $n+0.5$ is representable.
tanpi $(n+0.5)$ for odd integer $n$ is $-\infty$ where $n+0.5$ is representable.
trunc $(-1<x<0)$ returns -0.

### 7.5.2 Changes to C99 TC2 Behavior

modf behaves as though implemented by:

```
gentype modf ( gentype value, gentype *iptr )
{
    *iptr = trunc( value );
    return copysign( isinf( value ) ? 0.0 : value - *iptr, value );
}
```

rint always rounds according to round to nearest even rounding mode even if the caller is in some other rounding mode.

### 7.5.3 Edge Case Behavior in Flush To Zero Mode

If denormals are flushed to zero, then a function may return one of four results:

1. Any conforming result for non-flush-to-zero mode
2. If the result given by 1 . is a sub-normal before rounding, it may be flushed to zero
3. Any non-flushed conforming result for the function if one or more of its sub-normal operands are flushed to zero.
4. If the result of 3 . is a sub-normal before rounding, the result may be flushed to zero.

In each of the above cases, if an operand or result is flushed to zero, the sign of the zero is undefined.

If subnormals are flushed to zero, a device may choose to conform to the following edge cases for nextafter instead of those listed in section 7.5.1:
nextafter $(+$ smallest normal, $y<+$ smallest normal $)=+0$.
nextafter $(-$ smallest normal, $y>-$ smallest normal $)=-0$.
nextafter $(-0, y>0)$ returns smallest positive normal value.
nextafter $(+0, y<0)$ returns smallest negative normal value.
For clarity, subnormals or denormals are defined to be the set of representable numbers in the range $0<\mathrm{x}<$ TYPE_MIN and -TYPE_MIN $<\mathrm{x}<-0$. They do not include $\pm 0$. A non-zero number is said to be sub-normal before rounding if after normalization, its radix-2 exponent is less than (TYPE_MIN_EXP - 1). ${ }^{72}$

[^52]
## 8. Image Addressing and Filtering

Let $\mathrm{w}_{\mathrm{t}}, \mathrm{h}_{\mathrm{t}}$ and $\mathrm{d}_{\mathrm{t}}$ be the width, height (or image array size for a 1D image array) and depth (or image array size for a 2D image array) of the image in pixels. Let coord. xy also referred to as $(s, t)$ or coord. xyz also referred to as $(s, t, r)$ be the coordinates specified to read_image $\{\mathbf{f}|\mathbf{i}| \mathbf{u i}\}$. The sampler specified in read_image $\{\mathbf{f} \mathbf{i} \mathbf{i} \mid \mathbf{u i}\}$ is used to determine how to sample the image and return an appropriate color.

### 8.1 Image Coordinates

This affects the interpretation of image coordinates. If image coordinates specified to read_image $\{\mathbf{f}|\mathbf{i}| \mathbf{u i}\}$ are normalized (as specified in the sampler), the $s, t$, and $r$ coordinate values are multiplied by $\omega_{t}, h_{t}$, and $d_{t}$ respectively to generate the unnormalized coordinate values. For image arrays, the image array coordinate (i.e. $t$ if it is a 1D image array or $r$ if it is a 2D image array) specified to read_image\{f|i|ui\} must always be the un-normalized image coordinate value.

Let ( $u, v, w$ ) represent the unnormalized image coordinate values.

### 8.2 Addressing and Filter Modes

We first describe how the addressing and filter modes are applied to generate the appropriate sample locations to read from the image if the addressing mode is not CLK_ADDRESS_REPEAT nor CLK_ADDRESS_MIRRORED_REPEAT.

After generating the image coordinate ( $\mathrm{u}, \mathrm{v}, \mathrm{w}$ ) we apply the appropriate addressing and filter mode to generate the appropriate sample locations to read from the image.

If values in $(u, v, w)$ are INF or NaN, the behavior of read_image $\{\mathbf{f}|\mathbf{i}| \mathbf{u i}\}$ is undefined.

## Filter Mode $=$ CLK_FILTER_NEAREST

When filter mode is CLK_FILTER_NEAREST, the image element in the image that is nearest (in Manhattan distance) to that specified by ( $u, v, w$ ) is obtained. This means the image element at location ( $i, j, k$ ) becomes the image element value, where

```
i = address_mode((int)floor(u))
j = address_mode((int)floor(v))
k = address_mode((int)floor(w))
```

For a 3D image, the image element at location (i, j, k) becomes the color value. For a 2D image, the image element at location $(i, j)$ becomes the color value.

Table 8.1 describes the address_mode function.

| Addressing Mode | Result of address_mode(coord) |  |
| :--- | :--- | :---: |
| CLK_ADDRESS_CLAMP_TO_EDGE | clamp (coord, 0, size - 1) |  |
| CLK_ADDRESS_CLAMP | clamp (coord, -1, size) |  |
| CLK_ADDRESS_NONE | coord |  |

Table 8.1 Addressing modes to generate texel location.
The size term in table 8.1 is $w_{t}$ for $u, h_{t}$ for $v$ and $d_{t}$ for $w$.

The clamp function used in table 8.1 is defined as:

$$
\text { clamp }(a, b, c)=\text { return }(a<b) \text { ? b : }((a>c) \text { ? c : a) }
$$

If the selected texel location ( $i, j, k$ ) refers to a location outside the image, the border color is used as the color value for this texel.

## Filter Mode = CLK_FILTER_LINEAR

When filter mode is CLK_FILTER_LINEAR, a $2 \times 2$ square of image elements for a 2 D image or a $2 \times 2 \times 2$ cube of image elements for a 3D image is selected. This $2 \times 2$ square or 2 $x 2 \times 2$ cube is obtained as follows.

Let

```
i0 = address_mode((int)floor(u - 0.5))
j0 = address_mode((int)floor(v - 0.5))
k0 = address mode((int)floor(w - 0.5))
i1 = address_mode((int)floor(u - 0.5) + 1)
j1 = address_mode((int)floor(v - 0.5) + 1)
k1 = address_mode((int)floor(w - 0.5) + 1)
a = frac(u - 0.5)
b = frac(v - 0.5)
c = frac(w - 0.5)
```

where $\mathrm{frac}(\mathrm{x}$ ) denotes the fractional part of x and is computed as $\mathrm{x}-\mathrm{floor}(\mathrm{x})$.
For a 3D image, the image element value is found as

$$
\begin{aligned}
\mathrm{T}= & (1-\mathrm{a}) \star(1-\mathrm{b}) \star(1-c) \star \mathrm{T}_{\mathrm{iOj0k0}} \\
& +\mathrm{a} *(1-\mathrm{b}) \star(1-c) \star \mathrm{T}_{\mathrm{i} 1 \mathrm{j} 0 \mathrm{k} 0}
\end{aligned}
$$

$$
\begin{aligned}
& +(1-a) * b *(1-c) * T_{i 0 j 1 k 0} \\
& +a * b *(1-c) * T_{i 1 j 1 k 0} \\
& +(1-a) \star(1-b) * c * T_{i 0 j 0 k 1} \\
& +a \star(1-b) \star c * T_{i 1 j 0 k 1} \\
& +(1-a) * b * c * T_{i 0 j 1 k 1} \\
& +a \neq b * c * T_{i 1 j 1 k 1}
\end{aligned}
$$

where $T_{i j k}$ is the image element at location (i,j,k) in the 3D image.
For a 2D image, the image element value is found as

$$
\begin{aligned}
\mathrm{T}= & (1-a) \star(1-b) \star \mathrm{T}_{i 0 j 0} \\
& +a *(1-b) \star \mathrm{T}_{i 1 j 0} \\
& +(1-a) \star b * \mathrm{~T}_{i 0 j 1} \\
& +a * \mathrm{~b} * \mathrm{~T}_{i 1 j 1}
\end{aligned}
$$

where $T_{i j}$ is the image element at location $(i, j)$ in the 2D image.
If any of the selected $\mathrm{T}_{\mathrm{ijk}}$ or $\mathrm{T}_{i j}$ in the above equations refers to a location outside the image, the border color is used as the color value for $\mathrm{T}_{i j k}$ or $\mathrm{T}_{i j}$.

If the image channel type is CL_FLOAT or CL_HALF_FLOAT and any of the image elements $\mathrm{T}_{i j k}$ or $\mathrm{T}_{i j}$ is INF or NaN, the behavior of the built-in image read function is undefined.

We now discuss how the addressing and filter modes are applied to generate the appropriate sample locations to read from the image if the addressing mode is CLK_ADDRESS_REPEAT.

If values in $(s, t, r)$ are INF or NaN, the behavior of the built-in image read functions is undefined.

## Filter Mode $=$ CLK_FILTER_NEAREST

When filter mode is CLK_FILTER_NEAREST, the image element at location (i,j,k) becomes the image element value, with $i, j$ and $k$ computed as

```
u = (s - floor(s)) * wt
i = (int)floor(u)
if (i > wt - 1)
        i = i - wt
v = (t - floor(t)) * ht
j = (int)floor(v)
if (j > ht - 1)
    j = j - ht
```

```
w = (r - floor(r)) * dt
k = (int)floor(w)
if (k > dt - 1)
    k}=\textrm{k}-\mp@subsup{\textrm{d}}{\textrm{t}}{
```

For a 3D image, the image element at location (i, $j, k$ ) becomes the color value. For a 2D image, the image element at location ( $i, j$ ) becomes the color value.

## Filter Mode = CLK_FILTER_LINEAR

When filter mode is CLK_FILTER_LINEAR, a $2 \times 2$ square of image elements for a 2D image or a $2 \times 2 \times 2$ cube of image elements for a 3D image is selected. This $2 \times 2$ square or 2 x 2 x 2 cube is obtained as follows.

Let

```
u = (s - floor(s)) * wt
i0 = (int)floor(u - 0.5)
i1 = i0 + 1
if (i0 < 0)
    iO = wt + i0
if (il > wt - 1)
        il = il - Wt
v = (t - floor(t)) * ht
j0 = (int)floor(v - 0.5)
j1 = j0 + 1
if (j0< < )
        j0 = ht + j0
if (j1 > ht - 1)
        j1 = j1 - ht
w = (r - floor(r)) * dt
k0 = (int)floor(w - 0.5)
k1 = k0 + 1
if (k0 < 0)
        k0 = dt + k0
if (k1 > dt - 1)
        k1 = k1 - dt
a = frac(u - 0.5)
b = frac(v - 0.5)
c = frac(w - 0.5)
```

where $f$ rac ( $x$ ) denotes the fractional part of $x$ and is computed as $x-f l o o r(x)$.

For a 3D image, the image element value is found as

$$
\begin{aligned}
\mathrm{T}= & (1-a) \star(1-b) *(1-c) \star \mathrm{T}_{i 0 j 0 k 0} \\
& +a *(1-b) *(1-c) * T_{i 1 j 0 k 0} \\
& +(1-a) * b *(1-c) * T_{i 0 j 1 k 0} \\
& +a * b *(1-c) * T_{i 1 j 1 k 0} \\
& +(1-a) *(1-b) * c * T_{i 0 j 0 k 1} \\
& +a *(1-b) * c * T_{i 1 j 0 k 1} \\
& +(1-a) * b * c * T_{i 0 j 1 k 1} \\
& +a * b * c * T_{i 1 j 1 k 1}
\end{aligned}
$$

where $T_{i j k}$ is the image element at location ( $i, j, k$ ) in the 3D image.
For a 2D image, the image element value is found as

$$
\begin{aligned}
\mathrm{T}= & (1-a) \star(1-b) * \mathrm{~T}_{i 0 j 0} \\
& +a *(1-b) * \mathrm{~T}_{i 1 j 0} \\
& +(1-a) \star b * \mathrm{~T}_{i 0 j 1} \\
& +a * b * \mathrm{~T}_{i 1 j 1}
\end{aligned}
$$

where $T_{i j}$ is the image element at location $(i, j)$ in the 2D image.
If the image channel type is CL_FLOAT or CL_HALF_FLOAT and any of the image elements $\mathrm{T}_{\mathrm{ijk}}$ or $\mathrm{T}_{\mathrm{ij}}$ is INF or NaN, the behavior of the built-in image read function is undefined.

We now discuss how the addressing and filter modes are applied to generate the appropriate sample locations to read from the image if the addressing mode is
CLK_ADDRESS_MIRRORED_REPEAT. The CLK_ADDRESS_MIRRORED_REPEAT addressing mode causes the image to be read as if it is tiled at every integer seam with the interpretation of the image data flipped at each integer crossing. For example, the ( $s, t, r$ ) coordinates between 2 and 3 are addressed into the image as coordinates from 1 down to 0 . If values in $(s, t, r)$ are INF or NaN , the behavior of the built-in image read functions is undefined.

## Filter Mode $=$ CLK_FILTER_NEAREST

When filter mode is CLK_FILTER_NEAREST, the image element at location (i,j,k) becomes the image element value, with $i, j$ and $k$ computed as

```
s' = 2.0f * rint(0.5f * s)
s' = fabs(s - s')
u = s' * wt
i = (int)floor(u)
i = min(i, wt - 1)
```

```
t'=2.0f * rint(0.5f * t)
t'}=\textrm{fabs}(t-\mp@subsup{t}{}{\prime}
v = t' * ht
j = (int)floor(v)
j = min(j, ht - 1)
r'=2.0f * rint(0.5f * r)
r' = fabs(r - r')
w = r'}**\mp@subsup{d}{t}{
k = (int)floor(w)
k}=\operatorname{min}(k,\mp@subsup{d}{t}{}-1
```

For a 3D image, the image element at location ( $i, j, k$ ) becomes the color value. For a 2D image, the image element at location $(i, j)$ becomes the color value.

## Filter Mode = CLK_FILTER_LINEAR

When filter mode is CLK_FILTER_LINEAR, a $2 \times 2$ square of image elements for a 2D image or a $2 \times 2 \times 2$ cube of image elements for a 3D image is selected. This $2 \times 2$ square or 2 $x 2 \times 2$ cube is obtained as follows.

Let

```
s' = 2.0f * rint(0.5f * s)
s' = fabs(s - s')
u = s' * wt
i0 = (int) floor(u - 0.5f)
i1 = i0 + 1
i0 = max(i0, 0)
i1 = min(i1, wt - 1)
t'=2.0f * rint(0.5f * t)
t' = fabs(t - t')
v = t' * ht
j0 = (int)floor(v - 0.5f)
j1 = j0 + 1
j0 = max(j0, 0)
j1 = min(j1, ht - 1)
r' = 2.0f * rint(0.5f * r)
r' = fabs(r - r')
w = r' * dt
k0 = (int)floor(w - 0.5f)
k1 = k0 + 1
k0 = max(k0, 0)
k1 = min(k1, dt - 1)
```

$$
\begin{aligned}
& \mathrm{a}=\operatorname{frac}(\mathrm{u}-0.5) \\
& \mathrm{b}=\operatorname{frac}(\mathrm{v}-0.5) \\
& \mathrm{c}=\operatorname{frac}(\mathrm{w}-0.5)
\end{aligned}
$$

where $\mathrm{frac}(\mathrm{x})$ denotes the fractional part of x and is computed as $\mathrm{x}-\mathrm{floor}(\mathrm{x})$.
For a 3D image, the image element value is found as

$$
\begin{aligned}
& T=(1-a) *(1-b) *(1-c) * T_{i 0 j 0 k 0} \\
& +a *(1-b) *(1-c) * T_{i 1 j 0 k 0} \\
& +(1-a) * b *(1-c) * T_{i 0 j 1 k 0} \\
& +\mathrm{a} * \mathrm{~b} *(1-\mathrm{c}) * \mathrm{~T}_{\mathrm{i} 1 \mathrm{j1k0}} \\
& +(1-a) *(1-b) * c * T_{i 0 j 0 k 1} \\
& +\mathrm{a} *(1-\mathrm{b}) \text { * c * Tilj0k1 } \\
& +(1-a) * b * c * T_{i 0 j 1 k 1} \\
& +\mathrm{a} * \mathrm{~b} \text { * c * } \mathrm{T}_{\mathrm{i} 1 \mathrm{j} 1 \mathrm{k} 1}
\end{aligned}
$$

where $T_{i j k}$ is the image element at location (i,j,k) in the 3D image.
For a 2D image, the image element value is found as

$$
\begin{aligned}
\mathrm{T}= & (1-a) \star(1-b) * \mathrm{~T}_{i 0 j 0} \\
& +a *(1-b) * \mathrm{~T}_{i 1 j 0} \\
& +(1-a) \star b * \mathrm{~T}_{i 0 j 1} \\
& +a * b * \mathrm{~T}_{i 1 j 1}
\end{aligned}
$$

where $T_{i j}$ is the image element at location $(i, j)$ in the 2D image.
For a 1D image, the image element value is found as

$$
\begin{aligned}
\mathrm{T}= & (1-a) \star \mathrm{T}_{i 0} \\
& +\mathrm{a}^{\star} \mathrm{T}_{\mathrm{i} 1}
\end{aligned}
$$

where $T_{i}$ is the image element at location (i) in the 1D image.
If the image channel type is CL_FLOAT or CL_HALF_FLOAT and any of the image elements $\mathrm{T}_{i j k}$ or $\mathrm{T}_{i j}$ is INF or NaN, the behavior of the built-in image read function is undefined.

## NOTE

If the sampler is specified as using unnormalized coordinates (floating-point or integer coordinates), filter mode set to CLK_FILTER_NEAREST and addressing mode set to one of the following modes - CLK_ADDRESS_NONE, CLK_ADDRESS_CLAMP_TO_EDGE or CLK_ADDRESS_CLAMP, the location of the image element in the image given by $(i, j, k)$ in section 8.2 will be computed without any loss of precision.

For all other sampler combinations of normalized or unnormalized coordinates, filter and addressing modes, the relative error or precision of the addressing mode calculations and the image filter operation are not defined by this revision of the OpenCL specification. To ensure a minimum precision of image addressing and filter calculations across any OpenCL device, for these sampler combinations, developers should unnormalize the image coordinate in the kernel and implement the linear filter in the kernel with appropriate calls to read_image\{f|i|ui\} with a sampler that uses unnormalized coordinates, filter mode set to CLK_FILTER_NEAREST, addressing mode set to CLK_ADDRESS_NONE, CLK_ADDRESS_CLAMP_TO_EDGE or CLK_ADDRESS_CLAMP and finally performing the interpolation of color values read from the image to generate the filtered color value.

### 8.3 Conversion Rules

In this section we discuss conversion rules that are applied when reading and writing images in a kernel.

### 8.3.1 Conversion rules for normalized integer channel data types

In this section we discuss converting normalized integer channel data types to floating-point values and vice-versa.

### 8.3.1.1 Converting normalized integer channel data types to floating-point values

For images created with image channel data type of CL_UNORM_INT8 and CL_UNORM_INT16, read_imagef will convert the channel values from an 8-bit or 16-bit unsigned integer to normalized floating-point values in the range $[0.0 f$... 1.0].

For images created with image channel data type of CL_SNORM_INT8 and CL_SNORM_INT16, read_imagef will convert the channel values from an 8-bit or 16-bit signed integer to normalized floating-point values in the range $[-1.0$... 1.0].

These conversions are performed as follows:
CL_UNORM_INT8 (8-bit unsigned integer) $\rightarrow$ float
normalized float value $=($ float $) \mathrm{c} / 255.0 \mathrm{f}$
CL_UNORM_INT_101010 (10-bit unsigned integer) $\rightarrow$ float

```
    normalized float value = (float)c / 1023.0f
```

CL_UNORM_INT16 (16-bit unsigned integer) $\rightarrow$ float
normalized float value $=($ float $) \mathrm{c} /$ 65535.0f
CL_SNORM_INT8 (8-bit signed integer) $\rightarrow$ float
normalized float value $=\max (-1.0 f, \quad(f l o a t) c / 127.0 f)$
CL_SNORM_INT16 (16-bit signed integer) $\rightarrow$ float
normalized float value $=\max (-1.0 f,(f l o a t) c / 32767.0 f)$
The precision of the above conversions is $<=1.5$ ulp except for the following cases.
For CL_UNORM_INT8

```
0 must convert to \(0.0 f\) and 255 must convert to \(1.0 f\)
```

For CL_UNORM_INT_101010

```
0 must convert to \(0.0 f\) and 1023 must convert to \(1.0 f\)
```

For CL_UNORM_INT16

```
0 must convert to \(0.0 f\) and 65535 must convert to \(1.0 f\)
```

For CL_SNORM_INT8

```
-128 and -127 must convert to -1.0f,
    0 must convert to 0.0f and
    127 must convert to 1.0f
```

For CL_SNORM_INT16

$$
\begin{aligned}
& -32768 \text { and }-32767 \text { must convert to }-1.0 f \text {, } \\
& 0 \text { must convert to } 0.0 f \text { and } \\
& 32767 \text { must convert to } 1.0 f
\end{aligned}
$$

### 8.3.1.2 Converting floating-point values to normalized integer channel data types

For images created with image channel data type of CL_UNORM_INT8 and CL_UNORM_INT16, write_imagef will convert the floating-point color value to an 8-bit or 16-bit unsigned integer.

For images created with image channel data type of CL_SNORM_INT8 and CL_SNORM_INT16, write_imagef will convert the floating-point color value to an 8-bit or 16-bit signed integer.

The preferred method for how conversions from floating-point values to normalized integer values are performed is as follows:

```
float \(\rightarrow\) CL_UNORM_INT8 (8-bit unsigned integer)
        convert_uchar_sat_rte(f * 255.0f)
float \(\rightarrow\) CL_UNORM_INT_101010 (10-bit unsigned integer)
        min(convert_ushort_sat_rte(f * 1023.0f), 0x3ff)
float \(\rightarrow\) CL_UNORM_INT16 (16-bit unsigned integer)
        convert_ushort_sat_rte(f * 65535.0f)
float \(\rightarrow\) CL_SNORM_INT8 (8-bit signed integer)
        convert_char_sat_rte(f * 127.0f)
float \(\rightarrow\) CL_SNORM_INT16 (16-bit signed integer)
convert_short_sat_rte(f * 32767.0f)
```

Please refer to section 6.2.3.3 for out-of-range behavior and saturated conversions rules.
OpenCL implementations may choose to approximate the rounding mode used in the conversions described above. If a rounding mode other than round to nearest even (_rte) is used, the absolute error of the implementation dependant rounding mode vs. the result produced by the round to nearest even rounding mode must be $<=0.6$.
float $\rightarrow$ CL_UNORM_INT8 (8-bit unsigned integer)

$$
\text { Let } f_{\text {preferred }}=\text { convert_uchar_sat_rte(f * 255.0f) }
$$

```
Let fapprox =
    convert_uchar_sat_<impl-rounding-mode>(f * 255.0f)
fabs(flpreferred - fapprox) must be <= 0.6
```

float $\rightarrow$ CL_UNORM_INT_101010 (10-bit unsigned integer)

```
Let f}\mp@subsup{f}{preferred = convert_ushort_sat_rte(f * 1023.0f)}{(f)
Let fapprox =
    convert_ushort_sat_<impl-rounding-mode>(f * 1023.0f)
fabs(fpreferred - fapprox) must be <= 0.6
```

float $\rightarrow$ CL_UNORM_INT16 (16-bit unsigned integer)

```
Let f}\mp@subsup{f}{\mathrm{ preferred = convert_ushort_sat_rte(f * 65535.0f)}}{\mathrm{ (f }
Let f}\mp@subsup{f}{\mathrm{ approx }}{=
    convert_ushort_sat_<impl-rounding-mode>(f * 65535.0f)
fabs(f}\mp@subsup{f}{\mathrm{ preferred }}{}-\mp@subsup{f}{\mathrm{ approx }}{})\mathrm{ must be }<=0.
```

float $\rightarrow$ CL_SNORM_INT8 (8-bit signed integer)

```
Let fpreferred = convert_char_sat_rte(f * 127.0f)
Let fapprox =
        convert_char_sat_<impl_rounding_mode>(f * 127.0f)
fabs(fpreferred - fapprox) must be <= 0.6
```

float $\rightarrow$ CL_SNORM_INT16 (16-bit signed integer)

```
Let fpreferred = convert_short_sat_rte(f * 32767.0f)
Let fapprox =
    convert_short_sat_<impl-rounding-mode>(f * 32767.0f)
fabs(fpreferred - fapprox) must be <= 0.6
```


### 8.3.2 Conversion rules for half precision floating-point channel data type

For images created with a channel data type of CL_HALF_FLOAT, the conversions from half to float are lossless (as described in section 6.1.1.1). Conversions from float to half round the mantissa using the round to nearest even or round to zero rounding mode. Denormalized numbers for the half data type which may be generated when converting a float to a half
may be flushed to zero. A float NaN must be converted to an appropriate NaN in the half type. A float INF must be converted to an appropriate INF in the half type.

### 8.3.3 Conversion rules for floating-point channel data type

The following rules apply for reading and writing images created with channel data type of CL_FLOAT.

NaNs may be converted to a NaN value(s) supported by the device.
Denorms can be flushed to zero.

All other values must be preserved.

### 8.3.4 Conversion rules for signed and unsigned 8-bit, 16-bit and 32-bit integer channel data types

Calls to read_imagei with channel data type values of CL_SIGNED_INT8, CL_SIGNED_INT16 and CL_SIGNED_INT32 return the unmodified integer values stored in the image at specified location.

Calls to read_imageui with channel data type values of CL_UNSIGNED_INT8, CL_UNSIGNED_INT16 and CL_UNSIGNED_INT32 return the unmodified integer values stored in the image at specified location.

Calls to write_imagei will perform one of the following conversions:
32 bit signed integer $\rightarrow 8$-bit signed integer

```
convert_char_sat(i)
```

32 bit signed integer $\rightarrow$ 16-bit signed integer

```
convert_short_sat(i)
```

32 bit signed integer $\rightarrow$ 32-bit signed integer

```
no conversion is performed
```

Calls to write_imageui will perform one of the following conversions:
32 bit unsigned integer $\rightarrow 8$-bit unsigned integer

```
convert_uchar_sat(i)
```

32 bit unsigned integer $\rightarrow$ 16-bit unsigned integer

```
convert_ushort_sat(i)
```

32 bit unsigned integer $\rightarrow$ 32-bit unsigned integer

```
no conversion is performed
```

The conversions described in this section must be correctly saturated.

### 8.3.5 Conversion rules for sRGBA and sBGRA images

Standard RGB data, which roughly displays colors in a linear ramp of luminosity levels such that an average observer, under average viewing conditions, can view them as perceptually equal steps on an average display. All 0 's maps to 0.0 f , and all 1 's maps to 1.0 f . The sequence of unsigned integer encodings between all 0 's and all 1's represent a nonlinear progression in the floating-point interpretation of the numbers between 0.0 f to 1.0 f . For more detail, see the SRGB color standard, IEC 61996-2-1, at IEC (International Electrotechnical Commission).

Conversion from sRGB space is automatically done by read_imagef built-in functions if the image channel order is one of the sRGB values described above. When reading from an sRGB image, the conversion from sRGB to linear RGB is performed before the filter specified in the sampler specified to read_imagef is applied. If the format has an alpha channel, the alpha data is stored in linear color space. Conversion to sRGB space is automatically done by write_imagef built-in functions if the image channel order is one of the sRGB values described above and the device supports writing to sRGB images.

If the format has an alpha channel, the alpha data is stored in linear color space.
The following is the conversion rule for converting a normalized 8-bit unsigned integer sRGB color value to a floating-point linear RGB color value using read imagef.

```
Convert the normalized 8-bit unsigned integer R, G and B
channel values to a floating-point value (call it c) as per
rules described in section 8.3.1.1.
if (c <= 0.04045),
    result = c / 12.92;
else
    result = powr((c + 0.055) / 1.055, 2.4);
```

The resulting floating point value, if converted back to an sRGB value without rounding to a 8-
bit unsigned integer value, must be within 0.5 ulp of the original sRGB value.
The following are the conversion rules for converting a linear RGB floating-point color value (call it c) to a normalized 8-bit unsigned integer sRGB value using write imagef.

```
if (c is NaN) c = 0.0;
if (c > 1.0)
    c = 1.0;
else if (c < 0.0)
    c = 0.0;
else if (c < 0.0031308)
    c = 12.92 * c;
else
    c}=1.055*\mp@subsup{c}{}{1.0/2.4 - 0.055;
scaled_reference_result = c * 255
channe\overline{l}_componen\overline{t}= L scaled_reference_result + 1/2 \rfloor
```

The precision of the above conversion should be such that :
| generated_channel_component - scaled_reference_result | <= 0.6 where generated_channel_component is the actual value that the implementation produces and being checked for conformance.

### 8.4 Selecting an Image from an Image Array

Let ( $u, v, w$ ) represent the unnormalized image coordinate values for reading from and/or writing to a 2D image in a 2D image array.

When read using a sampler, the 2 D image layer selected is computed as:

```
layer = clamp(rint(w), 0, dt - 1)
```

otherwise the layer selected is computed as:
layer = w
(since $w$ is already an integer) and the result is undefined if $w$ is not one of the integers $0,1, \ldots \mathrm{~d}_{\mathrm{t}}$ - 1 .

Let $(u, v)$ represent the unnormalized image coordinate values for reading from and/or writing to a 1D image in a 1D image array.

When read using a sampler, the 1 D image layer selected is computed as:

```
layer = clamp(rint(v), 0, ht - 1)
```

otherwise the layer selected is computed as:
layer = v
(since $v$ is already an integer) and the result is undefined if $v$ is not one of the integers $0,1, \ldots \mathrm{~h}_{\mathrm{t}}$ 1.


[^0]:    ${ }^{1}$ When any scalar value is converted to bool, the result is 0 if the value compares equal to 0 ; otherwise, the result is 1.
    ${ }^{2}$ The double scalar type is an optional type that is supported if CL_DEVICE_DOUBLE_FP_CONFIG in table 4.3 for a device is not zero.

[^1]:    ${ }^{3}$ Built-in vector data types are supported by the OpenCL implementation even if the underlying compute device does not support any or all of the vector data types. These are to be converted by the device compiler to appropriate instructions that use underlying built-in types supported natively by the compute device. Refer to Appendix B for a description of the order of the components of a vector type in memory.
    ${ }^{4}$ The double vector type is an optional type that is supported if CL_DEVICE_DOUBLE_FP_CONFIG in table 4.3 for a device is not zero.

[^2]:    ${ }^{5}$ That is, for the purpose of applying type-based aliasing rules, a built-in vector data type will be considered equivalent to the corresponding array type.

[^3]:    ${ }^{6}$ Unless the cl_khr_fp16 extension is supported and has been enabled.
    ${ }^{7}$ Unless the cl_khr_fp16 extension is supported and has been enabled.

[^4]:    ${ }^{8}$ For conversions to floating-point format, when a finite source value exceeds the maximum representable finite floating-point destination value, the rounding mode will affect whether the result is the maximum finite floatingpoint value or infinity of same sign as the source value, per IEEE-754 rules for rounding.

[^5]:    ${ }^{9}$ In addition, some other extensions to the C language designed to support particular vector ISA (e.g. AltiVec ${ }^{\mathrm{TM}}$, CELL Broadband Engine ${ }^{\text {TM }}$ Architecture) use such conversions in conjunction with swizzle operators to achieve type unconversion. So as to support legacy code of this type, as_typen() allows conversions between vectors of the same size but different numbers of elements, even though the behavior of this sort of conversion is not likely to be portable except to other OpenCL implementations for the same hardware architecture. AltiVec ${ }^{\text {TM }}$ is a trademark of Motorola Inc. Cell Broadband Engine is a trademark of Sony Computer Entertainment, Inc.

[^6]:    ${ }^{10}$ Only if double precision is supported.
    ${ }^{11}$ Unless the cl_khr_fp16 extension is supported and has been enabled.
    ${ }^{12}$ While the union is intended to reflect the organization of data in memory, the as_type () and as_typen () constructs are intended to reflect the organization of data in register. The as_type() and as_typen() constructs are intended to compile to no instructions on devices that use a shared register file designed to operate on both the operand and result types. Note that while differences in memory organization are expected to largely be limited to those arising from endianness, the register based representation may also differ due to size of the element in register. (For example, an architecture may load a char into a 32 -bit register, or a char vector into a SIMD vector register with fixed 32 -bit element size.) If the element count does not match, then the implementation should pick a data representation that most closely matches what would happen if an appropriate result type operator was applied to a register containing data of the source type. If the number of elements matches, then the as_typen() should faithfully reproduce the behavior expected from a similar data type reinterpretation using memory/unions. So, for example if an implementation stores all single precision data as double in register, it should implement as_int( float ) by first downconverting the double to single precision and then (if necessary) moving the single precision bits to a register suitable for operating on integer data. If data stored in different address spaces do not have the same endianness, then the "dominant endianness" of the device should prevail.

[^7]:    ${ }^{13}$ Only if double precision is supported.

[^8]:    ${ }^{14}$ This is different from the standard integer conversion rank described in C 99 TC 2 , section 6.3.1.1.

[^9]:    ${ }^{15}$ The pre- and post- increment operators may have unexpected behavior on floating-point values and are therefore not supported for floating-point scalar and vector built-in types. For example, if variable a has type float and holds the value $0 x 1.0 \mathrm{p} 25 \mathrm{f}$, then a++ returns $0 x 1.0 \mathrm{p} 25 \mathrm{f}$. Also, (a++)-- is not guaranteed to return a, if a has fractional value. In non-default rounding modes, (a++)-- may produce the same result as $a++$ or $a--$ for large $a$.

[^10]:    ${ }^{16}$ To test whether any or all elements in the result of a vector relational operator test true, for example to use in the context in an if () statement, please see the any and all builtins in section 6.11.6.
    ${ }^{17}$ To test whether any or all elements in the result of a vector equality operator test true, for example to use in the context in an if () statement, please see the any and all builtins in section 6.11.6.

[^11]:    ${ }^{18}$ Integer promotion is described in ISO/IEC 9899:1999 in section 6.3.1.1.

[^12]:    ${ }^{19}$ Variable length arrays are not supported in OpenCL C. Refer to section 6.9.d.
    ${ }^{20}$ Except for 3-component vectors whose size is defined as $4 *$ size of each scalar component.
    ${ }^{21}$ Bit-field struct members are not supported in OpenCL C. Refer to section 6.9.c.

[^13]:    ${ }^{22}$ Among the invalid values for dereferencing a pointer by the unary $*$ operator are a null pointer, an address inappropriately aligned for the type of object pointed to, and the address of an object after the end of its lifetime. If $* \mathbf{P}$ is an lvalue and $\mathbf{T}$ is the name of an object pointer type, $*(\mathbf{T}) \mathbf{P}$ is an lvalue that has a type compatible with that to which $\mathbf{T}$ points.
    ${ }^{23}$ Thus, $\boldsymbol{\&} * \mathbf{E}$ is equivalent to $\mathbf{E}$ (even if $\mathbf{E}$ is a null pointer), and $\boldsymbol{\&}(\mathbf{E} 1[\mathbf{E} 2])$ to $((\mathbf{E} 1)+(\mathbf{E} 2))$. It is always true that if $\mathbf{E}$ is an lvalue that is a valid operand of the unary $\boldsymbol{\&}$ operator, $* \boldsymbol{\&} \mathbf{E}$ is an lvalue equal to $\mathbf{E}$.

[^14]:    ${ }^{24}$ Implicit in autovectorization is the assumption that any libraries called from the __kernel must be recompilable at run time to handle cases where the compiler decides to merge or separate workitems. This probably means that such libraries can never be hard coded binaries or that hard coded binaries must be accompanied either by source or some retargetable intermediate representation. This may be a code security question for some.

[^15]:    ${ }^{25}$ Items struckthrough are restrictions in a previous version of OpenCL C that are no longer present in OpenCL C 2.0.

[^16]:    ${ }^{26}$ Unless the cl_khr_fp16 extension is supported and has been enabled.

[^17]:    ${ }^{27}$ This syntax is already part of the clang source tree on which most vendors have based their OpenCL implementations. Additionally, blocks based closures are supported by the clang open source C compiler as well as Mac OS X's C and Objective C compilers. Specifically, Mac OS X's Grand Central Dispatch allows applications to queue tasks as a block.

[^18]:    ${ }^{28}$ OpenCL C does not allow function pointers (see section 6.9) primarily because it is difficult or expensive to implement generic indirections to executable code in many hardware architectures that OpenCL targets. OpenCL C's design of Blocks is intended to respect that same condition, yielding the restrictions listed here. As such, Blocks allow a form of dynamically enqueued function scheduling without providing a form of runtime synchronous dynamic dispatch analogous to function pointers.

[^19]:    ${ }^{29}$ i.e. the global_work_size values specified to clEnqueueNDRangeKernel are not evenly divisable by the local_work_size values for each dimension.

[^20]:    ${ }^{30} \mathbf{f m i n}$ and fmax behave as defined by C99 and may not match the IEEE 754-2008 definition for minNum and maxNum with regard to signaling NaNs. Specifically, signaling NaNs may behave as quiet NaNs.
    ${ }^{31}$ The $\min ()$ operator is there to prevent fract(-small ) from returning 1.0. It returns the largest positive floatingpoint number less than 1.0.

[^21]:    ${ }^{32}$ The user is cautioned that for some usages, e.g. $\operatorname{mad}(a, b,-a * b)$, the definition of $\operatorname{mad}()$ is loose enough that almost any result is allowed from $\boldsymbol{m a d}()$ for some values of $a$ and $b$.

[^22]:    ${ }^{33}$ Frequently vector operations need $n+1$ bits temporarily to calculate a result. The rhadd instruction gives you an extra bit without needing to upsample and downsample. This can be a profound performance win.

[^23]:    ${ }^{34}$ The mix and smoothstep functions can be implemented using contractions such as mad or fma.

[^24]:    ${ }^{35}$ The geometric functions can be implemented using contractions such as mad or fma.

[^25]:    ${ }^{36}$ If an implementation extends this specification to support IEEE-754 flags or exceptions, then all builtin functions defined in table 6.14 shall proceed without raising the invalid floating-point exception when one or more of the operands are NaNs.

[^26]:    ${ }^{37}$ The half $n$ type is only defined by the cl_khr_fp16 extension described in section 9.5 of the OpenCL 1.2 Extension Specification.

[^27]:    38 vload3, and vload_half3 read $x, y, z$ components from address $(p+(o f f s e t * 3))$ into a 3-component vector. vstore3, and vstore_half3 write $x, y, z$ components from a 3-component vector to address ( $p+($ offset * 3) ).

[^28]:    ${ }^{39}$ The built-in function barrier has been renamed work_group_barrier. For backward compatibility, barrier is also supported.
    ${ }^{40}$ Refer to section 6.13 .11 for description of memory_scope.

[^29]:    ${ }^{41}$ async_work_group_copy and async_work_group_strided_copy for 3-component vector types behave as async_work_group_copy and async_work_group_strided_copy respectively for 4-component vector types.

[^30]:    42 The C11 consume operation is not supported.

[^31]:    ${ }^{43}$ We can't require C11 atomics since host programs can be implemented in other programming languages and versions of C or $\mathrm{C}++$, but we do require that the host programs use atomics and that those atomics be compatible with those in C11.

[^32]:    ${ }^{44}$ This value for memory_scope can only be used with atomic_work_item_fence with flags set to CLK_IMAGE_MEM_FENCE.

[^33]:    ${ }^{45}$ The atomic_long and atomic_ulong types are supported if the cl_khr_int64_base_atomics and cl_khr_int64_extended_atomics extensions are supported and have been enabled.
    ${ }^{46}$ The atomic_double type is only supported if double precision is supported and the cl_khr_int64_base_atomics and cl_khr_int64_extended_atomics extensions are supported and have been enabled.

[^34]:    ${ }^{47}$ If the device address space is 64-bits, the data types atomic_intptr_t, atomic_uintptr_t, atomic_size_t and atomic_ptrdiff_t are supported if the cl_khr_int64_base_atomics and cl_khr_int64_extended_atomics extensions are supported and have been enabled.

[^35]:    ${ }^{48}$ This spurious failure enables implementation of compare-and-exchange on a broader class of machines, e.g. load-locked store-conditional machines.

[^36]:    ${ }^{49}$ Only if the cl_khr_fp16 extension is supported and has been enabled.
    ${ }^{50}$ Only if double precision is supported.

[^37]:    ${ }^{51}$ Note that $\mathbf{0}$ is taken as a flag, not as the beginning of a field width.
    ${ }^{52}$ The results of all floating conversions of a negative zero, and of negative values that round to zero, include a minus sign.

[^38]:    ${ }^{53}$ Only if the cl_khr_fp16 extension is supported and has been enabled.

[^39]:    ${ }^{54}$ When applied to infinite and NaN values, the,-+ , and space flag characters have their usual meaning; the \# and $\mathbf{0}$ flag characters have no effect.
    55 Binary implementations can choose the hexadecimal digit to the left of the decimal-point character so that subsequent digits align to nibble (4-bit) boundaries.

[^40]:    ${ }^{56}$ No special provisions are made for multibyte characters. The behavior of printf with the $\mathbf{s}$ conversion specifier is undefined if the argument value is not a pointer to a literal string.

[^41]:    ${ }^{57}$ This is similar to the GL_ADDRESS_CLAMP_TO_BORDER addressing mode.

[^42]:    ${ }^{58}$ The built-in function calls to read images with a sampler are not supported for image1d_buffer_t image types.

[^43]:    ${ }^{59}$ Only if the cl_khr_fp16 extension is supported and has been enabled.
    ${ }^{60}$ Only if double precision is supported.

[^44]:    ${ }^{61}$ The half scalar and vector types can only be used if the cl_khr_fp16 extension is supported and has been enabled. The double scalar and vector types can only be used if double precision is supported.

[^45]:    ${ }^{62}$ The half scalar and vector types can only be used if the cl_khr_fp16 extension is supported and has been enabled. The double scalar and vector types can only be used if double precision is supported.

[^46]:    ${ }^{63}$ The half scalar and vector types can only be used if the cl_khr_fp16 extension is supported and has been enabled. The double scalar and vector types can only be used if double precision is supported.

[^47]:    ${ }^{64}$ Implementations are not required to honor this flag. Implementations may not schedule kernel launch earlier than the point specified by this flag, however.
    ${ }^{65}$ Immediate meaning not side effects resulting from child kernels. The side effects would include stores to global memory and pipe reads and writes.
    ${ }^{66}$ This acts as a memory synchronization point between work-items in a work-group and child kernels enqueued by work-items in the work-group.

[^48]:    ${ }^{67}$ Except for the embedded profile whether either round to zero or round to nearest rounding mode may be supported for single precision floating-point.

[^49]:    ${ }^{68}$ The ULP values for built-in math functions lgamma and lgamma_r is currently undefined.
    ${ }^{69} 0 \mathrm{ulp}$ is used for math functions that do not require rounding.

[^50]:    ${ }^{70} 0$ ulp is used for math functions that do not require rounding.

[^51]:    ${ }^{71} 0$ ulp is used for math functions that do not require rounding.

[^52]:    ${ }^{72}$ Here TYPE_MIN and TYPE_MIN_EXP should be substituted by constants appropriate to the floating-point type under consideration, such as FLT_MIN and FLT_MIN_EXP for float.

